

# Commutation via Relative Termination

work in progress

**Nao Hirokawa**

JAIST

Aart Middeldorp

University of Innsbruck

## Confluence by Decreasingness

### THEOREM

Hirokawa and Middeldorp 2011

left-linear and locally confluent TRS  $\mathcal{R}$  is confluent if

$\text{CPS}(\mathcal{R})/\mathcal{R}$  is terminating

- critical peak steps  
$$\overbrace{\text{CPS}(\mathcal{R})} = \left\{ \begin{array}{l} s \rightarrow t \\ s \rightarrow u \end{array} \middle| t \xleftarrow{\mathcal{R}} s \xrightarrow{\mathcal{R}} u \text{ is critical peak} \right\}$$
- $$\xrightarrow{S/\mathcal{R}} = \xrightarrow{\mathcal{R}}^* \cdot \xrightarrow{S} \cdot \xrightarrow{\mathcal{R}}^*$$

### NOTE

generalization of

- Rosen's orthogonality (1973)
- left-linear case of Knuth and Bendix' completeness criterion

## Example I

consider locally confluent **non-terminating** TRS  $\mathcal{R}$ :

$$\begin{array}{lll} x + 0 \rightarrow x & x + s(y) \rightarrow s(x + y) & \text{ones} \rightarrow s(0) : \text{ones} \\ 0 + y \rightarrow y & s(x) + y \rightarrow s(x + y) & \end{array}$$

critical peak steps  $\text{CPS}(\mathcal{R})$ :

$$\begin{array}{lll} 0 + 0 \rightarrow 0 & & \\ 0 + s(y) \rightarrow s(0 + y) & s(x) + 0 \rightarrow s(x + 0) & s(x) + s(y) \rightarrow s(x + s(y)) \\ 0 + s(y) \rightarrow s(y) & s(x) + 0 \rightarrow s(x) & s(x) + s(y) \rightarrow s(s(x) + y) \end{array}$$

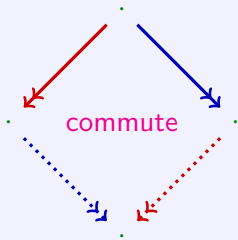
$\text{CPS}(\mathcal{R})/\mathcal{R}$  is terminating. hence  $\mathcal{R}$  is confluent

- presence of overlaps is no problem
- very weak at commutative rule  $x + y \rightarrow y + x$

# Commutativity

## DEFINITION

$\rightarrow$  and  $\rightarrow$  **commute** if  $*\leftarrow \cdot \rightarrow^* \subseteq \rightarrow^* \cdot *\leftarrow$



$\rightarrow$  and  $\rightarrow$  **locally commute** if  $\leftarrow \cdot \rightarrow \subseteq \rightarrow^* \cdot *\leftarrow$

## LEMMA

- $\mathcal{R}$  is confluent if  $\mathcal{R}$  and  $\mathcal{R}$  commute
- $\mathcal{R} \cup \mathcal{S}$  is confluent if  $\mathcal{R}$  and  $\mathcal{S}$  are confluent and commute

## Example I

is next TRS confluent?

$$x + 0 \rightarrow x \quad x + s(y) \rightarrow s(x + y) \quad x + p(y) \rightarrow p(x + y)$$

$$0 + y \rightarrow y \quad s(x) + y \rightarrow s(x + y) \quad p(x) + y \rightarrow p(x + y)$$

$$x + y \rightarrow y + x$$

**not confluent:**

$$p(s(x + y)) \xleftarrow[\mathcal{R}]{*} s(x) + p(y) \xrightarrow[\mathcal{R}]{*} s(p(x + y))$$

## Commutation by Closedness

do  $\mathcal{R}$  and  $\mathcal{S}$  commute?

$$x + 0 \rightarrow x \qquad x + s(y) \rightarrow s(x + y) \qquad x + p(y) \rightarrow p(x + y)$$

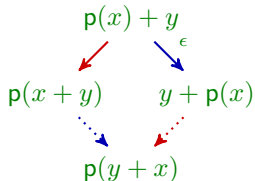
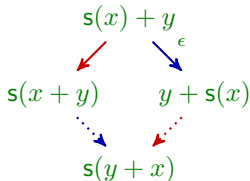
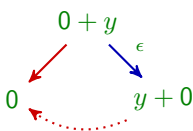
$$0 + y \rightarrow y \qquad s(x) + y \rightarrow s(x + y) \qquad p(x) + y \rightarrow p(x + y)$$

$$x + y \rightarrow y + x$$

**THEOREM** Toyama 1988; van Oostrom 1994; Aoto, Yoshida and Toyama 2009

left-linear TRSs  $\mathcal{R}$  and  $\mathcal{S}$  commute if

$$\leftarrow_{\mathcal{R}} \times \xrightarrow{\epsilon}_{\mathcal{S}} \subseteq \leftarrow_{\mathcal{S}} \cdot \xrightarrow{*}_{\mathcal{R}} \qquad \leftarrow_{\mathcal{R}} \times \xrightarrow{>\epsilon}_{\mathcal{S}} \subseteq \leftarrow_{\mathcal{R}} \circ \leftarrow_{\mathcal{S}}$$



## Aim: Unify

### THEOREM

Aoto, Yoshida and Toyama 2009

left-linear TRSs  $\mathcal{R}$  and  $\mathcal{S}$  commute if

$$\begin{array}{c} \longleftarrow \\ \mathcal{R} \end{array} \times \begin{array}{c} \xrightarrow{\epsilon} \\ \mathcal{S} \end{array} \subseteq \begin{array}{c} \longleftarrow \\ \mathcal{S} \end{array} \cdot \begin{array}{c} \xrightarrow{*} \\ \mathcal{R} \end{array} \qquad \begin{array}{c} \xrightarrow{\epsilon} \\ \mathcal{R} \end{array} \times \begin{array}{c} \xrightarrow{\epsilon} \\ \mathcal{S} \end{array} \subseteq \begin{array}{c} \longleftarrow \\ \mathcal{R} \end{array}$$

### THEOREM

Hirokawa and Middeldorp 2011

left-linear and locally confluent TRS  $\mathcal{R}$  is confluent if

$\text{CPS}(\mathcal{R})/\mathcal{R}$  is terminating

$$\text{CPS}(\mathcal{R}) = \left\{ \begin{array}{l} s \rightarrow t \\ s \rightarrow u \end{array} \middle| t \xleftarrow{\mathcal{R}} s \xrightarrow{\mathcal{R}} u \text{ is critical peak} \right\}$$

Q

find generalized criterion

# Main Result

## THEOREM

left-linear and locally commuting TRSs  $\mathcal{R}$  and  $\mathcal{S}$  commute if  
 $\text{CPS}(\mathcal{R}, \mathcal{S}) / (\mathcal{R} \cup \mathcal{S})$  is terminating

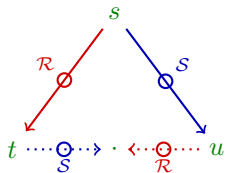
$$\text{CPS}(\mathcal{R}, \mathcal{S}) = \left\{ \begin{array}{l} s \rightarrow t \\ s \rightarrow u \end{array} \middle| \begin{array}{l} t \xleftarrow{\mathcal{R}} s \xrightarrow{\mathcal{S}} u \text{ is non-closed critical peak} \end{array} \right\}$$

- $t \xleftarrow{\mathcal{R}} s \xrightarrow{\mathcal{S}} u$  is closed if  $t \xrightarrow{\mathcal{S}} \cdot \xleftarrow{\mathcal{R}}^* u$
- $t \xleftarrow{\mathcal{R}} s \xrightarrow{\mathcal{S}} u$  is closed if  $t \xleftarrow{\mathcal{R}} u$

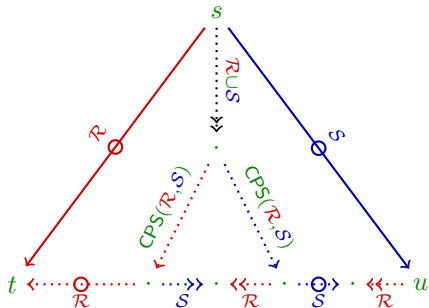


# Proof

if  $t \xrightarrow[\mathcal{R}]{\circlearrowleft} s \xrightarrow[S]{} u$  then



or



$(\rightarrow, \rightarrow)$  is decreasing with respect to  $>$ , where

- $>$  is extension of  $\rightarrow_{CPS(\mathcal{R}, S) / (\mathcal{R}US)}$  with bottom  $\perp$

- $t \rightarrow_s u$  if  $s = \perp$  and  $t \xrightarrow[\mathcal{R}]{} u$        $t \rightarrow_s u$  if  $s \xrightarrow[\mathcal{R}US]{}^* t \xrightarrow[S]{} u$

## Remark

our proof

- uses **proof terms** for multisteps:
  - let  $F(y)$  denote rule  $0 + y \rightarrow y$
  - $s(0 + 0) \dashv\dashv s(0)$  is witnessed by  $s(F(0))$
  
- separates **critical peak lemma** from main proof

## Example of Direct Usage

left-linear locally confluent TRS  $\mathcal{R}$

$$\begin{array}{lll} x + 0 \rightarrow x & x + s(y) \rightarrow s(x + y) & x + p(y) \rightarrow p(x + y) \\ 0 + y \rightarrow y & s(x) + y \rightarrow s(x + y) & p(x) + y \rightarrow p(x + y) \\ x + y \rightarrow y + x & s(p(x)) \rightarrow x & p(s(x)) \rightarrow x \end{array}$$

non-closed critical peak steps  $\text{CPS}_{\mathcal{R}}(\mathcal{R})$

$$\begin{array}{ll} p(x) + s(y) \rightarrow s(p(x) + y) & p(x) + s(y) \rightarrow p(x + s(y)) \\ s(x) + p(y) \rightarrow p(s(x) + y) & s(x) + p(y) \rightarrow s(x + p(y)) \\ s(x) + p(y) \rightarrow s(x + p(y)) & s(x) + p(y) \rightarrow p(s(x) + y) \\ p(x) + s(y) \rightarrow p(x + s(y)) & p(x) + s(y) \rightarrow s(p(x) + y) \\ x + s(p(y)) \rightarrow x + y & x + s(p(y)) \rightarrow s(x + p(y)) \\ s(p(x)) + y \rightarrow x + y & s(p(x)) + y \rightarrow s(p(x) + y) \\ x + p(s(y)) \rightarrow x + y & x + p(s(y)) \rightarrow p(x + s(y)) \\ p(s(x)) + y \rightarrow x + y & p(s(x)) + y \rightarrow p(s(x) + y) \end{array}$$

$\text{CPS}(\mathcal{R}, \mathcal{R})/\mathcal{R}$  is terminating. hence  $\mathcal{R}$  is confluent

## Experiments on 168 left-linear TRSs in Cops

	direct				decomposition	
	closed	critical	peak	new	rule label	new+rl
yes	21		28	41	50	62
timeout	0		0	0	1	<b>106</b>

- 30 sec timeout
- local confluence/commutation check by 4-steps rewriting
- (relative) termination check by  $T_1T_2$
- Bruteforce for decomposition

REMARK

$$41 - (21 \cup 28) = 2$$

# Summary

unified

**Toyama (1988)**

commutativity

**van Oostrom (1994)**

development closed

**H-M (2011)**

critical peak steps

## FUTURE WORK

- simplify proof
- efficient algorithm for decomposition
- integrate more:

**Okui (1998)**

simultaneous closed

**Oyamaguchi-Ohta (2004)**

upside parallel closed

**van Oostrom (2012)**

critical valley steps