Observation of the definitions of the classes…

Def: Class $\mathcal{P}$ (Chapter 5)
Set L is in the class $\mathcal{P}$ ⇔
There exists a poly-time computable predicate R such that
for each $x \in \Sigma^*$, $x \in L \Leftrightarrow R(x)$

Def: Class $\mathcal{NP}$ (Def 5.2)
Set L is in the class $\mathcal{NP}$ ⇔
There exists a poly q and a poly-time computable pred. R s.t.
for each $x \in \Sigma^*$, $x \in L \Leftrightarrow \exists w \in \Sigma^*: |w| \leqq q(|x|)[R(x,w)]$

Def: Class co-$\mathcal{NP}$ (Theorem 5.5)
Set L is in the class co-$\mathcal{NP}$ ⇔
There exists a poly q and a poly-time computable pred. R s.t.
for each $x \in \Sigma^*$, $x \in L \Leftrightarrow \forall w \in \Sigma^*: |w| \leqq q(|x|)[R(x,w)]$

計算量クラス間の定義を概観すると…

クラス$\mathcal{P}$の定義(5章)
集合Lがクラス$\mathcal{P}$に入る ⇔
以下を満たす多項式時間計算可能述語Rが存在:
各 $x \in \Sigma^*$で$x \in L \Leftrightarrow R(x)$

クラス$\mathcal{NP}$の定義(定義5.2)
集合Lがクラス$\mathcal{NP}$に入る ⇔
以下を満たす多項式qと多項式時間計算可能述語Rが存在:
各 $x \in \Sigma^*$で$x \in L \Leftrightarrow \exists w \in \Sigma^*: |w| \leqq q(|x|)[R(x,w)]$

クラスco-$\mathcal{NP}$の定義(定理5.5)
集合Lがクラスco-$\mathcal{NP}$に入る ⇔
以下を満たす多項式qと多項式時間計算可能述語Rが存在:
各 $x \in \Sigma^*$で$x \in L \Leftrightarrow \forall w \in \Sigma^*: |w| \leqq q(|x|)[R(x,w)]$

Chapter 6. Analysis on Polynomial-Time Computability

6.1. Polynomial-time Reducibility

Def.6.1:
Let A and B be arbitrary sets.
(1) function $h: A \rightarrow B$: polynomial-time reduction
⇔
(a) h is a total function from Σ^* onto Σ^*
(b) $x \in \Sigma^*[x \in A \leftrightarrow h(x) \in B]$
(c) h is polynomial-time computable.

(2) When there is a polynomial-time reduction from A to B, we say A is polynomial-time reducible to B.
Then, we denote by $A \leq_m^P B$

第6章 多項式時間計算可能性の分析

6.1. 多項式時間還元可能性

定義6.1:
AとBを任意の集合とする.
(1) 関数 $h: A \rightarrow B$: 多項式時間還元(polynomial-time reduction)
⇔
(a) h はΣ^*からΣ^*への全域的関数
(b) $x \in \Sigma^*[x \in A \leftrightarrow h(x) \in B]$
(c) h は多項式時間計算可能.

(2) AからBへの多項式時間還元が存在するとき,
AはBへ多項式時間還元可能という(polynomial time reducible).
このとき，次のように書く: $A \leq_m^P B$

$A \leq_m^P B$ within polynomial time, hardness of A $\leq$ that of B

定理6.1 $A \leq_m^P B$ leads to,

(1) $B \in \mathcal{P} \rightarrow A \in \mathcal{P}$.
(2) $B \in \mathcal{NP} \rightarrow A \in \mathcal{NP}$.
(3) $B \in$ co-$\mathcal{NP}$ $\rightarrow A \in$ co-$\mathcal{NP}$.
(4) $B \in \mathcal{EXP} \rightarrow A \in \mathcal{EXP}$.

Note: class $\mathcal{E}$ is exceptional. Generally, $B \in \mathcal{E} \rightarrow A \in \mathcal{E}$ is not true.

Ex.6.2: If we define ONE$\equiv \{1\}$, for each set L in $\mathcal{P}$ we have
$L \leq_m^P$ ONE

If we define $h(x) \equiv \begin{cases} 1, & \text{if } x \in L, \\ 0, & \text{otherwise} \end{cases}$

(1) h is a total function from Σ^* onto Σ^*.
(2) $x \in \Sigma^*[x \in L \leftrightarrow h(x) \in \mathrm{ONE}]$
(3) h is polynomial-time computable(so is computation $L \in \mathcal{P} \rightarrow x \in L$)

$A \leq_m^P B$ 多項式時間の範囲内では, Aの難しさ$\leq$ Bの難しさ

定理6.1. $A \leq_m^P B$ のとき,

(1) $B \in \mathcal{P} \rightarrow A \in \mathcal{P}$.
(2) $B \in \mathcal{NP} \rightarrow A \in \mathcal{NP}$.
(3) $B \in$ co-$\mathcal{NP}$ $\rightarrow A \in$ co-$\mathcal{NP}$.
(4) $B \in \mathcal{EXP} \rightarrow A \in \mathcal{EXP}$.

補注:クラス$\mathcal{E}$は例外. 一般には, $B \in \mathcal{E} \rightarrow A \in \mathcal{E}$とはならない.

例6.2: ONE$\equiv \{1\}$と定義するとき, クラス$\mathcal{P}$のすべての集合Lについて $L \leq_m^P$ ONE
が成り立つ. $h(x) \equiv \begin{cases} 1, & x \in L\text{のとき}, \\ 0, & \text{その他のとき} \end{cases}$
と定義すると, (1) hはΣ^*からΣ^*への全域的関数.
(2) $x \in \Sigma^*[x \in L \leftrightarrow h(x) \in \mathrm{ONE}]$
(3) hは多項式時間計算可能($L \in \mathcal{P} \rightarrow x \in L$の判定も多項式時間内)

Theorem 6.2: A, B, C: arbitrary sets

(1) $A \le_m^P A$

(2) $A \le_m^P B \wedge B \le_m^P C \to A \le_m^P C$

Def: $A \equiv_m^P B \leftrightarrow A \le_m^P B \wedge B \le_m^P A$

$\equiv_m^P$ is an equivalence relation.

定理6.2: A, B, C: 任意の集合

(1) $A \le_m^P A$

(2) $A \le_m^P B \wedge B \le_m^P C \to A \le_m^P C$

定義: $A \equiv_m^P B \leftrightarrow A \le_m^P B \wedge B \le_m^P A$

$\equiv_m^P$ は同値関係

Relation among satisfiability problems of propositional expressions

2SAT (propositional satisfiability problem)
3SAT
SAT
ExSAT (extended propositional satisfiability problem)

$2\text{SAT} \le_m^P 3\text{SAT}$

- at most k…trivial
- exactly k…
  - ➢easy if you can repeat the same literal.
  - ➢the other case … good exercise!

Similarly,

$3\text{SAT} \le_m^P \text{SAT} \le_m^P \text{ExSAT}$

$2\text{SAT} \le_m^P 3\text{SAT} \le_m^P \text{SAT} \le_m^P \text{ExSAT}$ (6.1)

Here, if we can show

$\text{ExSAT} \le_m^P 3\text{SAT}$

then we have

$3\text{SAT} \equiv_m^P \text{SAT} \equiv_m^P \text{ExSAT}$

命題論理式の充足可能性問題の間の関係

2SAT （命題論理式充足性問題：二和積形式）
3SAT （命題論理式充足性問題：三和積形式）
SAT （命題論理式充足性問題）
ExSAT （拡張命題論理式充足性問題）

$2\text{SAT} \le_m^P 3\text{SAT}$

- 高々k個…自明
- ちょうどk個…
  - ➢同じリテラルを使ってよいなら簡単。
  - ➢だめなら…考えてみよう！

同様に，

$3\text{SAT} \le_m^P \text{SAT} \le_m^P \text{ExSAT}$

$2\text{SAT} \le_m^P 3\text{SAT} \le_m^P \text{SAT} \le_m^P \text{ExSAT}$ (6.1)

ここで

$\text{ExSAT} \le_m^P 3\text{SAT}$

であることを示せると，

$3\text{SAT} \equiv_m^P \text{SAT} \equiv_m^P \text{ExSAT}$

となる．

Ex. 6.3: Reduction from ExSAT to 3SAT

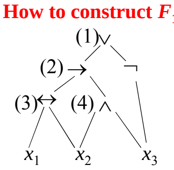
$$E_1(x_1,x_2,x_3) \equiv [[x_1 \leftrightarrow x_2] \to [x_2 \wedge x_3]] \vee \neg x_3$$

$$F_1(x_1,x_2,x_3) \equiv U_1 \wedge [U_1 \leftrightarrow [U_2 \vee \neg x_3]] \wedge [U_2 \leftrightarrow [U_3 \to U_4]] \wedge [U_3 \leftrightarrow [x_1 \leftrightarrow x_2]] \wedge [U_4 \leftrightarrow [x_2 \wedge x_3]]$$

Then, [E_1 is satisfiable] $\leftrightarrow$ [F_1 is satisfiable] (6.2)

F_1 is easier to be converted to 3SAT form.

How to construct F_1



(1) $V_1 \equiv V_2 \vee \neg x_3$

(2) $V_2 \equiv [V_3 \to V_4]$

(3) $V_3 \equiv [x_1 \leftrightarrow x_2]$

(4) $V_4 \equiv x_2 \wedge x_3$

To construct F_1 we let $V_i \to U_i$, and connect expressions of V_i by $\wedge$

例6.3: ExSATから3SATへの還元

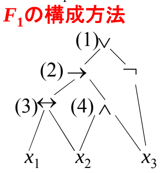
$$E_1(x_1,x_2,x_3) \equiv [[x_1 \leftrightarrow x_2] \to [x_2 \wedge x_3]] \vee \neg x_3$$

$$F_1(x_1,x_2,x_3) \equiv U_1 \wedge [U_1 \leftrightarrow [U_2 \vee \neg x_3]] \wedge [U_2 \leftrightarrow [U_3 \to U_4]] \wedge [U_3 \leftrightarrow [x_1 \leftrightarrow x_2]] \wedge [U_4 \leftrightarrow [x_2 \wedge x_3]]$$

このとき，[E_1が充足可能] $\leftrightarrow$ [F_1が充足可能] (6.2)

F_1は三和積形式に直しやすい形になっている．

F_1の構成方法



(1) $V_1 \equiv V_2 \vee \neg x_3$

(2) $V_2 \equiv [V_3 \to V_4]$

(3) $V_3 \equiv [x_1 \leftrightarrow x_2]$

(4) $V_4 \equiv x_2 \wedge x_3$

F_1を構成するために，$V_i \to U_i$とし，V_iの定義式を $\wedge$で結ぶ

From the construction of F_1
(1) F_1 is never true unless each U_i is $V_i(x_1, x_2, x_3)$.
(2) If each U_i is $V_i(x_1, x_2, x_3)$, we have $F_1 = E_1$

The above properties are proved by using induction.
proof is omitted.

Conversion to 3SAT form

$a \to b = \neg a \vee b$
$a \leftrightarrow b = (a \to b) \wedge (b \to a) = [\neg a \vee b] \wedge [\neg b \vee a]$: useful relations

$$
\begin{aligned}
U_1 \leftrightarrow [U_2 \vee \neg x_3] &= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg[U_2 \vee \neg x_2]] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee [\neg U_2 \wedge x_2]] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg U_2] \wedge [U_1 \vee x_2] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg U_2 \vee \neg U_2] \wedge [U_1 \vee x_2 \vee x_2]
\end{aligned}
$$

Others are similar.
Thus, every 3SAT form is converted.

F_1 の構成方法より,
(1)各U_i の値を$V_i(x_1, x_2, x_3)$としない限り, F_1は真にはならない.
(2)各U_iの値を$V_i(x_1, x_2, x_3)$としたとき, $F_1 = E_1$

上の性質が成り立つことは, 帰納法を用いるなどして証明可能.
証明は省略.

三和積形式への変換

$a \to b = \neg a \vee b$
$a \leftrightarrow b = (a \to b) \wedge (b \to a) = [\neg a \vee b] \wedge [\neg b \vee a]$であることを用いる.

$$
\begin{aligned}
U_1 \leftrightarrow [U_2 \vee \neg x_3] &= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg[U_2 \vee \neg x_2]] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee [\neg U_2 \wedge x_2]] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg U_2] \wedge [U_1 \vee x_2] \\
&= [\neg U_1 \vee U_2 \vee \neg x_3] \wedge [U_1 \vee \neg U_2 \vee \neg U_2] \wedge [U_1 \vee x_2 \vee x_2]
\end{aligned}
$$

他も同様.
よって, すべて三和積形式に変形できることがわかる.

6.2.Completeness based on Polynomial-time Reducibility

6.2.1. Definition of Completeness and its Basic Properties

Def.6.2: For a class $\mathcal{C}$, if a set A satisfies the following conditions, then it is called $\mathcal{C}$-complete (under $\leq_m^P$)
(a) $\forall L \in \mathcal{C}\, [L \leq_m^P A]$
(b) $A \in \mathcal{C}$
Note: Sets satisfying the condition (a) are called $\mathcal{C}$-hard.

6.2. 多項式時間還元可能性に基づく完全性

6.2.1. 完全性の定義とその基本的性質

定義6.2: 計算量クラス$\mathcal{C}$に対し, 集合Aが次の条件を満たすとき, それを($\leq_m^P$の下で)$\mathcal{C}$-完全という.
(a) $\forall L \in \mathcal{C}\, [L \leq_m^P A]$
(b) $A \in \mathcal{C}$
補注: 条件(a)を満たす集合は$\mathcal{C}$-困難.

6.2.Completeness based on Polynomial-time Reducibility

6.2.1. Definition of Completeness and its Basic Properties

Ex.6.5. Examples of $\mathcal{NP}$-complete sets
3SAT, SAT, ExSAT, DHAM, KNAP, BIN, VC, etc
$\mathcal{EXP}$-complete sets
EVAL-IN-E, HALT-IN-E, etc.

EVAL-IN-E:
Input: $< a, x, \bar{t} >$
a: the code of a program with 1 input, $x \in \Sigma^*, \bar{t} \geq 0$
Output: $\mathit{eval\text{-}in\text{-}time}(a, x, \bar{2}^t) = \mathit{accept}$?

6.2. 多項式時間還元可能性に基づく完全性

6.2.1. 完全性の定義とその基本的性質

例6.5. クラス$\mathcal{NP}$の完全集合の例
3SAT, SAT, ExSAT, DHAM, KNAP, BIN, VCなど
クラス$\mathcal{EXP}$の完全集合
EVAL-IN-E, HALT-IN-Eなど

EVAL-IN-E:
入力: $< a, x, \bar{t} >$
a: 1入力プログラムのコード, $x \in \Sigma^*, \bar{t} \geq 0$
出力: $\mathit{eval\text{-}in\text{-}time}(a, x, \bar{2}^t) = \mathit{accept}$?

9/14

Theorem 6.3. For any $\mathcal{C}$-hard (or $\mathcal{C}$-complete) set A,

(1) $A \in \mathcal{P} \rightarrow \mathcal{C} \subseteq \mathcal{P}$ CP: $\mathcal{C} \not\subset \mathcal{P} \rightarrow A \notin \mathcal{P}$
(2) $A \in \mathcal{NP} \rightarrow \mathcal{C} \subseteq \mathcal{NP}$ CP: $\mathcal{C} \not\subset \mathcal{NP} \rightarrow A \notin \mathcal{NP}$
(3) $A \in$ co-$\mathcal{NP} \rightarrow \mathcal{C} \subseteq$ co-$\mathcal{NP}$ CP: $\mathcal{C} \not\subset$ co-$\mathcal{NP} \rightarrow A \notin$ co-$\mathcal{NP}$
(4) $A \in \mathcal{EXP} \rightarrow \mathcal{C} \subseteq \mathcal{EXP}$ CP: $\mathcal{C} \not\subset \mathcal{EXP} \rightarrow A \notin \mathcal{EXP}$

Proof: CP: contraposition
(1) Let B be any $\mathcal{C}$-set. Then, since A is $\mathcal{C}$-hard,
$B \leq_m^P A$ and by the assumption $A \in \mathcal{P}$ we have $B \in \mathcal{P}$ (Th. 6.1)
(2), (3), (4) are similar.

9/14

定理6.3. 任意の$\mathcal{C}$-困難集合(含:$\mathcal{C}$-完全集合)Aに対し,

(1) $A \in \mathcal{P} \rightarrow \mathcal{C} \subseteq \mathcal{P}$ 対偶は $\mathcal{C} \not\subset \mathcal{P} \rightarrow A \notin \mathcal{P}$
(2) $A \in \mathcal{NP} \rightarrow \mathcal{C} \subseteq \mathcal{NP}$ 対偶は $\mathcal{C} \not\subset \mathcal{NP} \rightarrow A \notin \mathcal{NP}$
(3) $A \in$ co-$\mathcal{NP} \rightarrow \mathcal{C} \subseteq$ co-$\mathcal{NP}$ 対偶は $\mathcal{C} \not\subset$ co-$\mathcal{NP} \rightarrow A \notin$ co-$\mathcal{NP}$
(4) $A \in \mathcal{EXP} \rightarrow \mathcal{C} \subseteq \mathcal{EXP}$ 対偶は $\mathcal{C} \not\subset \mathcal{EXP} \rightarrow A \notin \mathcal{EXP}$

証明:
(1) Bを任意の$\mathcal{C}$集合とすると, Aは$\mathcal{C}$-困難だから,
$B \leq_m^P A$ 一方, $A \in \mathcal{P}$の仮定より, $B \in \mathcal{P}$ (定理6.1)
(2), (3), (4)も同様

10/14

Theorem 6.3. For any $\mathcal{C}$-hard (or $\mathcal{C}$-complete) set A,

(1) $A \in \mathcal{P} \rightarrow \mathcal{C} \subseteq \mathcal{P}$ CP: $\mathcal{C} \not\subset \mathcal{P} \rightarrow A \notin \mathcal{P}$
(2) $A \in \mathcal{NP} \rightarrow \mathcal{C} \subseteq \mathcal{NP}$ CP: $\mathcal{C} \not\subset \mathcal{NP} \rightarrow A \notin \mathcal{NP}$
(3) $A \in$ co-$\mathcal{NP} \rightarrow \mathcal{C} \subseteq$ co-$\mathcal{NP}$ CP: $\mathcal{C} \not\subset$ co-$\mathcal{NP} \rightarrow A \notin$ co-$\mathcal{NP}$
(4) $A \in \mathcal{EXP} \rightarrow \mathcal{C} \subseteq \mathcal{EXP}$ CP: $\mathcal{C} \not\subset \mathcal{EXP} \rightarrow A \notin \mathcal{EXP}$

Theorem 5.9.
(1) $\mathcal{NP} \subseteq$ co-$\mathcal{NP}$
$\rightarrow \mathcal{NP} =$ co-$\mathcal{NP}$

Ex.6.6: Meaning of Theorem 6.3 (class $\mathcal{NP}$)
Let A be $\mathcal{NP}$-complete set.

By the contraposition of Theorem 6.3(1) we have
$\mathcal{NP} \neq \mathcal{P} \rightarrow A \notin \mathcal{P}$

By the contraposition of Theorem 6.3(3) and that of Theorem 5.9(1),
$A \notin$ co-$\mathcal{NP}$

That is, $\mathcal{NP}$-complete sets are $\mathcal{NP}$-sets that cannot be recognized in polynomial time unless $\mathcal{P} = \mathcal{NP}$.

10/14

定理6.3. 任意の$\mathcal{C}$-困難集合(含:$\mathcal{C}$-完全集合)Aに対し,

(1) $A \in \mathcal{P} \rightarrow \mathcal{C} \subseteq \mathcal{P}$ 対偶は $\mathcal{C} \not\subset \mathcal{P} \rightarrow A \notin \mathcal{P}$
(2) $A \in \mathcal{NP} \rightarrow \mathcal{C} \subseteq \mathcal{NP}$ 対偶は $\mathcal{C} \not\subset \mathcal{NP} \rightarrow A \notin \mathcal{NP}$
(3) $A \in$ co-$\mathcal{NP} \rightarrow \mathcal{C} \subseteq$ co-$\mathcal{NP}$ 対偶は $\mathcal{C} \not\subset$ co-$\mathcal{NP} \rightarrow A \notin$ co-$\mathcal{NP}$
(4) $A \in \mathcal{EXP} \rightarrow \mathcal{C} \subseteq \mathcal{EXP}$ 対偶は $\mathcal{C} \not\subset \mathcal{EXP} \rightarrow A \notin \mathcal{EXP}$

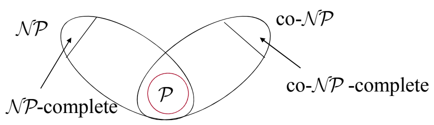
例6.6. 定理6.3の意味(クラス$\mathcal{NP}$)
Aを$\mathcal{NP}$-完全集合とする.

定理5.9.
(1) $\mathcal{NP} \subseteq$ co-$\mathcal{NP} \rightarrow \mathcal{NP} =$ co-$\mathcal{NP}$

定理6.3(1)の対偶より,
$\mathcal{NP} \neq \mathcal{P} \rightarrow A \notin \mathcal{P}$
定理6.3(3)の対偶と定理5.9(1)の対偶より,
$A \notin$ co-$\mathcal{NP}$
つまり, $\mathcal{NP}$-完全集合は$\mathcal{P} \neq \mathcal{NP}$である限り,
多項式時間では認識できない.

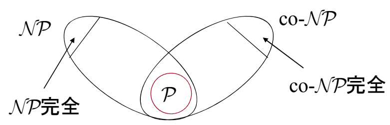
11/14

$\mathcal{NP}$-compete sets are $\mathcal{NP}$-sets that do not belong to $\mathcal{NP} \cap$ co-$\mathcal{NP}$ unless $\mathcal{P} = \mathcal{NP}$.



11/14

$\mathcal{NP}$-完全集合は$\mathcal{P} \neq \mathcal{NP}$である限り, $\mathcal{NP} \cap$ co-$\mathcal{NP}$には入らない$\mathcal{NP}$集合である.



Ex. 6.7. Meaning of Theorem 6.3 (class $\mathcal{EXP}$)

Let D be any $\mathcal{EXP}$-complete set.

Contraposition of Theorem 6.3(1)
 ($\mathcal{C} \not\subset \mathcal{P}$ → $A \notin \mathcal{P}$, where $\mathcal{EXP} \not\subset \mathcal{P}$ → $D \notin \mathcal{P}$)
 $\mathcal{P} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset \mathcal{P}$ ($\because \mathcal{P} \subseteq \mathcal{EXP}$) → $D \notin \mathcal{P}$
Contraposition of Theorem 6.3(2) ($\mathcal{C} \not\subset \mathcal{NP}$ → $A \notin \mathcal{NP}$,
 Here, $\mathcal{EXP} \not\subset \mathcal{NP}$ → $D \notin \mathcal{NP}$)
 $\mathcal{NP} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset \mathcal{NP}$ ($\because \mathcal{NP} \subseteq \mathcal{EXP}$) → $D \notin \mathcal{NP}$
Contraposition of Theorem 6.3(3) ($\mathcal{C} \not\subset$ co-$\mathcal{NP}$ → $A \not\subset$ co-$\mathcal{NP}$,
 here, $\mathcal{EXP} \not\subset$ co-$\mathcal{NP}$ → $D \notin$ co-$\mathcal{NP}$)
 co-$\mathcal{NP} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset$ co-$\mathcal{NP}$ → $D \notin$ co-$\mathcal{NP}$
But, by Theorem 5.7, since we know $\mathcal{P} \subsetneq \mathcal{EXP}$, we have $D \notin \mathcal{P}$.

$\mathcal{EXP}$-complete sets are not computable in polynomial time.

例6.7. 定理6.3の意味(クラス$\mathcal{EXP}$)

Dを$\mathcal{EXP}$-完全集合とする.

定理6.3(1)の対偶($\mathcal{C} \not\subset \mathcal{P}$ → $A \notin \mathcal{P}$, ここでは$\mathcal{EXP} \not\subset \mathcal{P}$ → $D \notin \mathcal{P}$)
 $\mathcal{P} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset \mathcal{P}$ ($\because \mathcal{P} \subseteq \mathcal{EXP}$) → $D \notin \mathcal{P}$
定理6.3(2)の対偶($\mathcal{C} \not\subset \mathcal{NP}$ → $A \notin \mathcal{NP}$,
 ここでは$\mathcal{EXP} \not\subset \mathcal{NP}$ → $D \notin \mathcal{NP}$)
 $\mathcal{NP} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset \mathcal{NP}$ ($\because \mathcal{NP} \subseteq \mathcal{EXP}$) → $D \notin \mathcal{NP}$
定理6.3(3)の対偶($\mathcal{C} \not\subset$ co-$\mathcal{NP}$ → $A \notin$ co-$\mathcal{NP}$,
 ここでは$\mathcal{EXP} \not\subset$ co-$\mathcal{NP}$ → $D \notin$ co-$\mathcal{NP}$)
 co-$\mathcal{NP} \neq \mathcal{EXP}$ → $\mathcal{EXP} \not\subset$ co-$\mathcal{NP}$ → $D \notin$ co-$\mathcal{NP}$
ところが定理5.7から $\mathcal{P} \subsetneq \mathcal{EXP}$ であるから, $D \notin \mathcal{P}$.

$\mathcal{EXP}$-完全集合は多項式時間では計算不可能.

Theorem 6.4. A: any $\mathcal{C}$-complete set
 For any set B we have
 (1) $A \leq_m^P B$ → B is $\mathcal{C}$-hard.
 (2) $A \leq_m^P B \wedge B \in \mathcal{C}$ → B is $\mathcal{C}$-complete.

Proof:
By Def. 6.2 $\forall L \in \mathcal{C}[L \leq_m^P A]$
By Theorem 6.2, $L \leq_m^P A \wedge A \leq_m^P B \rightarrow L \leq_m^P B$
Therefore, $\forall L \in \mathcal{C}[L \leq_m^P B]$
That is, B is $\mathcal{C}$-hard.

定理6.4. A: 任意の$\mathcal{C}$-完全集合
 すべての集合Bに対し,
 (1) $A \leq_m^P B$ → Bは$\mathcal{C}$-困難.
 (2) $A \leq_m^P B \wedge B \in \mathcal{C}$ → Bは$\mathcal{C}$-完全.

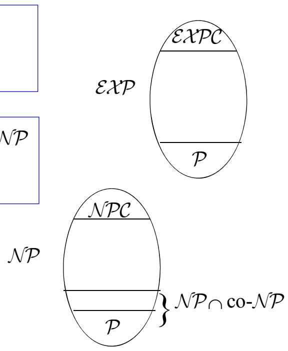
証明:
定義6.2より, $\forall L \in \mathcal{C}[L \leq_m^P A]$
定理6.2より, $L \leq_m^P A \wedge A \leq_m^P B \rightarrow L \leq_m^P B$
したがって, $\forall L \in \mathcal{C}[L \leq_m^P B]$

すなわち, Bは$\mathcal{C}$-困難.

$\mathcal{EXPC} \equiv \{L$: L is $\mathcal{EXP}$-complete$\}$
$\mathcal{NPC} \equiv \{L$: L is $\mathcal{NP}$-complete$\}$
Then, we have the following theorems.

Theorem 6.5.
 (1) $\mathcal{EXPC} \cap \mathcal{P} = \phi$
 (2) $\mathcal{EXP} - (\mathcal{EXPC} \cup \mathcal{P}) \neq \phi$

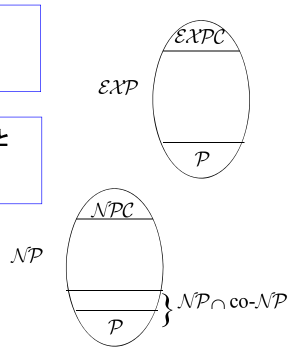
Theorem 6.6: Assuming $\mathcal{P} \neq \mathcal{NP}$
 (1) $\mathcal{NPC} \cap \mathcal{P} = \phi$
 (2) $\mathcal{NP} - (\mathcal{NPC} \cup \mathcal{P}) \neq \phi$



$\mathcal{EXPC} \equiv \{L$: Lは$\mathcal{EXP}$-完全$\}$
$\mathcal{NPC} \equiv \{L$: Lは$\mathcal{NP}$-完全$\}$
とすると, 次の定理が成り立つ.

定理6.5.
 (1) $\mathcal{EXPC} \cap \mathcal{P} = \phi$
 (2) $\mathcal{EXP} - (\mathcal{EXPC} \cup \mathcal{P}) \neq \phi$

定理6.6: $\mathcal{P} \neq \mathcal{NP}$**を仮定すると**
 (1) $\mathcal{NPC} \cap \mathcal{P} = \phi$
 (2) $\mathcal{NP} - (\mathcal{NPC} \cup \mathcal{P}) \neq \phi$



Schedule(残りの予定)

- 7/16(Thu) Office Hour (*Today!!*):
  - Comments on the report (レポートの解答と解説)
- 7/23(Thu): Last class (後半最後の講義)
  - Course Evaluation Questionnaire（授業アンケート）
  - Misc.（その他）
- 7/27(Mon): Final exam (期末試験)
  - 40 points ✖Textbook, Copy, Printout,…
  - You can bring your own hand-written notebook (手書きノートのみ持ち込み可)
  - Lesson 3~Lesson 6（講義3~講義6）