

# Categorical approach to first-order modal logic

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## Background: Neighborhood-Sheaves?

an objective: to understand

*neighborhood-sheaf semantics* [Kishida '11]

- a semantics of first-order modal logic (FOML)
- works for modal logic MC (and stronger ones)
- doesn't seem to work for logics weaker than MC
- But why?

N.B. Neighborhood-sheaf semantics is NOT a topos-theoretic semantics for FOML.

# This talk

## Categorical/coalgebraic model of FOML?

- Intention: by considering general settings we could clarify what is really necessary
- The talk consists of two parts:
  - ① What do we need to define a semantics of FOML?
  - ② How do we construct such a structure?

# Outline

- 1 Semantics of FOML
  - Modality and Coalgebra
  - Terms and Cartesian Category
  - Combining Them Together
- 2 How to Construct  $\mathcal{S}$ 
  - Constant Domain Model
  - “Sheaves”
- 3 Summary

# What is needed to interpret FOML?

FOML = prop. logic + **modality** + **terms** (+ quantifiers)

**Modality** is interpreted using

- a  **$T$ -coalgebra** for a functor  $T : \mathbf{Sets} \rightarrow \mathbf{Sets}$ , and
- a natural transformation  $\sigma : \mathcal{P} \rightarrow \mathcal{P} \circ T$  (so-called predicate lifting)
  - ▶  $\mathcal{P}$ : contravariant powerset functor

**Terms** are interpreted using

- a **cartesian category**  $\mathcal{S}$

To combine these structures we need

- a functor  $U : \mathcal{S} \rightarrow \mathbf{CoAlg}(T)$

# Kripke frame as coalgebra

## Definition (coalgebra)

For a functor  $T$ , a  $T$ -coalgebra is a morphism  $c : X \rightarrow TX$ .

## Ordinary Kripke frame

Kripke frame is  $(W, R)$ , where  $R \subseteq W \times W$ .

## Coalgebraic form

Kripke frame is a  $\mathcal{P}$ -coalgebra  $R : W \rightarrow \mathcal{P}W$ .

Ordinary and coalgebraic form are equivalent:

$$\mathcal{P}(W \times W) \simeq 2^{W \times W} \simeq (2^W)^W \simeq (\mathcal{P}W)^W$$

# Kripke semantics in coalgebraic form

$(W, R)$ : a Kripke frame

formulas are interpreted as a subset  $\llbracket \varphi \rrbracket \subseteq W$

## Ordinary Kripke semantics

when  $R$  is seen as a relation,

- $w \in \llbracket \Box \varphi \rrbracket$  is defined as  $\forall v. (w R v \implies v \in \llbracket \varphi \rrbracket)$

## Coalgebraic form

when  $R$  is seen as a map,

- $w \in \llbracket \Box \varphi \rrbracket$  is defined as  $R(w) \subseteq \llbracket \varphi \rrbracket$

# Modality as Predicate Lifting

- $\llbracket \Box \varphi \rrbracket = \{w \mid R(w) \subseteq \llbracket \varphi \rrbracket\} = R^{-1}\{Q \mid Q \subseteq \llbracket \varphi \rrbracket\}$
- define  $\sigma : \mathcal{P}X \rightarrow \mathcal{P}(\mathcal{P}X)$  by  $\sigma(P) = \{Q \mid Q \subseteq P\}$
- then  $\llbracket \Box \rrbracket = R^{-1} \circ \sigma$
- similarly  $\tau(P) = \{Q \mid Q \cap P \neq \emptyset\}$  defines  $\llbracket \Diamond \rrbracket = R^{-1} \circ \tau$
- topological interior operator is induced from
  - ▶  $W$ : a topological space
  - ▶  $N(w)$ : the set of neighborhoods of  $w \in W$
  - ▶  $\rho(P) = \{\mathcal{U} \subseteq \mathcal{P}W \mid P \in \mathcal{U}\}$as  $\mathbf{int}(P) = \{w \mid P \in N(w)\} = N^{-1} \circ \rho(P)$   
(topological interpretation of S4 modal logic)



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# Modality for a general coalgebra

In general,

- a functor  $T : \mathbf{Sets} \rightarrow \mathbf{Sets}$
- a **coalgebra**  $c : X \rightarrow TX$
- a natural transformation (predicate lifting)  $\sigma : \mathcal{P} \rightarrow \mathcal{P} \circ T$

induces a modality  $c^{-1} \circ \sigma : \mathcal{P}X \rightarrow \mathcal{P}(TX) \rightarrow \mathcal{P}X$ .

# Interpreting terms in a cartesian category

Terms are interpreted as functions (or morphisms).

- $\llbracket t \rrbracket : D^n \rightarrow D$ , if  $t$  is a term with  $n$  variables
  - ▶ e.g.  $\llbracket x + y + 2 \rrbracket : \mathbb{N}^2 \rightarrow \mathbb{N}$

$\llbracket t \rrbracket$  is defined as:

- constants  $\llbracket c \rrbracket : 1 \rightarrow D$  (given)
- $m$ -ary functions  $\llbracket f \rrbracket : D^m \rightarrow D$  (given)
- variables  $\llbracket x_i \rrbracket = \pi_i : D^n \rightarrow D$
- $\llbracket f(t_1, \dots, t_m) \rrbracket$  is a composition  
 $\llbracket f \rrbracket \circ \langle \llbracket t_1 \rrbracket, \dots, \llbracket t_m \rrbracket \rangle : D^n \rightarrow D^m \rightarrow D$

This makes sense in any **cartesian category**  $\mathcal{S}$  and object  $D \in \mathcal{S}$ .

# Coalgebras and cartesian category

We use

- **coalgebra** to interpret a modality
- **cartesian category** to interpret terms

related in some way. (N.B.  $\mathbf{CoAlg}(T)$  is not cartesian.)

Consider

- a functor  $T : \mathbf{Sets} \rightarrow \mathbf{Sets}$ ,
- a cartesian category  $\mathcal{S}$ , and
- a functor  $U : \mathcal{S} \rightarrow \mathbf{CoAlg}(T)$
- a predicate lifting  $\sigma : \mathcal{P} \rightarrow \mathcal{P} \circ T$

(such that  $\mathcal{P}(|U(\cdot)|)$  is a hyperdoctrine)

Then we can define a semantics of FOML.

# Interpreting FOML

- terms are interpreted in  $\mathcal{S}$
- formula  $\varphi$  is interpreted as  $\llbracket \varphi \rrbracket \subseteq |U(D^n)|$ 
  - ▶  $|U(D^n)|$  is the underlying set of  $U(D^n)$
- $\wedge, \vee, \neg, \rightarrow$  are set-theoretic operations
- $\Box$  is interpreted as  $\mathcal{P} \circ \sigma$
- $\forall, \exists$  are

$$\forall \llbracket \varphi \rrbracket = \{x \in |U(D^n)| \mid (U\pi)^{-1}(x) \subseteq \llbracket \varphi \rrbracket\}$$

$$\exists \llbracket \varphi \rrbracket = \{x \in |U(D^n)| \mid (U\pi)^{-1}(x) \cap \llbracket \varphi \rrbracket \neq \emptyset\}$$

where  $\pi : D^{n+1} \rightarrow D^n$ .

In other words, the right and left adjoints to  $(U\pi)^{-1}$ .

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# Constructing a concrete model of FOML?

A problem: How do we construct a cartesian category  $\mathcal{S}$ ?

- Typically, choose a suitable subcategory of  $\mathbf{CoAlg}(T)/X$  for a fixed  $X \in \mathbf{CoAlg}(T)$ .
- We will
  - ▶ describe a constant domain model
  - ▶ observe how Kripke/neighborhood sheaves are defined



# Constant domain model

- Fix  $c : X \rightarrow TX$ .
- For a set  $S$ ,  $\pi_2 : S \times X \rightarrow X$  can be turned into an object of  $\mathbf{CoAlg}(T)/X$ :

$$S \times X \simeq \coprod_{a \in S} X \rightarrow \coprod_{a \in S} TX \rightarrow T \left( \coprod_{a \in S} X \right) \simeq T(S \times X)$$

More concretely,  $(a, x) \mapsto T(\iota_a)(c(x))$   
where  $\iota_a : X \ni x \mapsto (a, x) \in S \times X$ .

- $\mathcal{S} \subseteq \mathbf{CoAlg}(T)/X$  consisting of coalgebras of this form is cartesian, and this  $\mathcal{S}$  gives a model of FOML.

# Kripke/neighborhood sheaves

Occasionally "locally isomorphic" maps  $D \rightarrow X$  form a cartesian category  $\mathcal{S} \subseteq \mathbf{CoAlg}(T)/X$ .

## Example

- Kripke sheaf  $p : D \rightarrow X$  is a p-morphism s.t.  
 $\forall d \in D, p$  is injective on  $R_D(d)$
- neighborhood sheaf  $p : D \rightarrow X$  is a p-morphism s.t.  
 $\forall d \in D, \exists U \in N_D(d), p$  is injective on  $U$

## Fact

$p : D \rightarrow X$  is a Kripke/neighborhood sheaf iff  
 $\Delta : D \rightarrow D \times_X D$  is a p-morphism (homomorphism)

# Sheaves in $\mathbf{CoAlg}(T)/X$ ?

Similar formulation in a more general setting?

## A Problem

Given  $T$ -hom  $p : D \rightarrow X$ , is  $D \times_X D$  a  $T$ -coalgebra?

- we have  $D \times_X D \rightarrow TD \times_{TX} TD$
- but not  $D \times_X D \rightarrow T(D \times_X D)$

What do we need to define  $D \times_X D \rightarrow T(D \times_X D)$ ?

# An observation on sheaves

- in Kripke/neighborhood sheaf cases, there are weaker notion of maps between coalgebras
  - ▶ graph homs, continuous maps
- the categories of Kripke/neighborhood frames and such maps have pullbacks, hence their slices are cartesian
- so  $D \times_X D$  is equipped with coalgebra structure
- define sheaf as  $p : D \rightarrow X$  s.t.  $\Delta : D \rightarrow D \times_X D$  is a homomorphism
- $\mathcal{S}$  is defined as the subcategory of sheaves

Open question: does this argument work in general?

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## Related work

Topos semantics for higher-order S4  
[Awodey, Kishida, Kotzsch '14]

- uses a certain Heyting-algebra object instead of  $\Omega$
- geometric morphism induces an S4 modality

# Summary and further questions

Summary:

- first-order modal logic is interpreted by
  - ▶ cartesian category  $\mathcal{S}$  and
  - ▶ a functor  $\mathcal{S} \rightarrow \mathbf{CoAlg}(T)$

Open question:

- recipe for constructing  $\mathcal{S}$ ?
- a notion of “sheaf” for coalgebras?