

Compiler of Minila

Lecture Note 5

Topics

- Quick reminder of (part of) the previous two lectures
- Compilation of expressions
- Compilation of statements
- Compiler

Quick Reminder (1)

- A typical Minila program is in the form:
 $S_1 S_2 \dots S_n$, where each S_i is as follows:
 - `estm`
 - `X := E ;`
 - `if E then Sa else Sb fi`
 - `while E do S od`
 - `for X Ea Eb do S od`
- Expressions are as follows:
 - N (a natural number), $v(i)$ (a variable, where $i = 1, 2, \dots$)
 - $E_a \text{ op } E_b$, where `op` is `++`, `--`, `**`, `//`, `%%`, `===`, `!===`, `<<`, `>>`, `&&`, or `||`.

Quick Reminder (2)

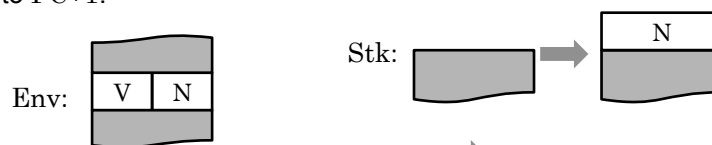
Let `PC` be a program counter, `Stk` a stack and `Env` an environment.

- `push(N)`:
 - Push N onto `Stk` and set `PC` to `PC+1`.



`PC: X` → `X+1`

- `load(V)`:
 - Find N corresponding to V in `Env`, push N onto `Stk`, and set `PC` to `PC+1`.

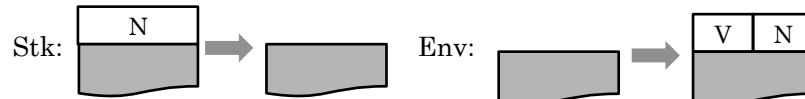


`PC: X` → `X+1`

Quick Reminder (3)

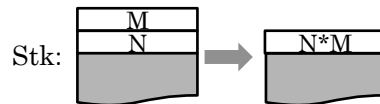
Let PC be a program counter, Stk a stack and Env an environment.

- $store(V)$:
 - Let N be the natural number at the top of Stk . Pop Stk , register V and N into Env , and set PC to $PC+1$.



$PC: X \rightarrow X+1$

- $multiply$:
 - Let M, N be the two natural numbers from the top of Stk . Pop Stk twice, push $N*M$ onto Stk , and set PC to $PC+1$.



$PC: X \rightarrow X+1$

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Quick Reminder (4)

- $divide$:
 - Let M, N be the two natural numbers from the top of Stk . Pop Stk twice, push the quotient of dividing N by M onto Stk , and set PC to $PC+1$.
- mod :
 - Let M, N be the two natural numbers from the top of Stk . Pop Stk twice, push the remainder of dividing N by M onto Stk , and set PC to $PC+1$.
- add :
 - Let M, N be the two natural numbers from the top of Stk . Pop Stk twice, push $N+M$ onto Stk , and set PC to $PC+1$.
- $minus$:
 - Let M, N be the two natural numbers from the top of Stk . Pop Stk twice, push the absolute value of the difference between N and M onto Stk , and set PC to $PC+1$.

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Quick Reminder (5)

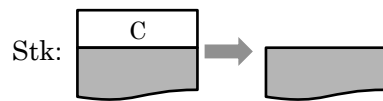
- lessThan:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 1 onto Stk if N is less than M and push 0 onto Stk otherwise, and set PC to PC+1.
- greaterThan:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 1 onto Stk if N is greater than M and push 0 onto Stk otherwise, and set PC to PC+1.
- equal:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 1 onto Stk if N equals M and push 0 onto Stk otherwise, and set PC to PC+1.

Quick Reminder (6)

- notEqual:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 0 onto Stk if N equals M and push 1 onto Stk otherwise, and set PC to PC+1.
- or:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 0 onto Stk if both N and M are 0 and push 1 onto Stk otherwise, and set PC to PC+1.
- and:
 - Let M, N be the two natural numbers from the top of Stk. Pop Stk twice, push 1 onto Stk if both N and M are not 0 and push 0 onto Stk otherwise, and set PC to PC+1.

Quick Reminder (7)

- `jump N`:
 - Set PC to $PC+N$.
- `bjump N`:
 - Set PC to $PC - N$.
- `jumpOnCond N`:
 - Let C be the natural number at the top of `Stk`. Pop `Stk` and set PC to $PC+N$ if C is not 0 and set PC to $PC+1$ otherwise.



PC: $X \rightarrow$ **if** $C = 0$ **then** $X+1$ **else** $X+N$

- `quit`:
 - Terminate the execution.

Quick Reminder (8)

- The compiler translates Minila programs into lists of instructions.
- When the compiler takes the program:

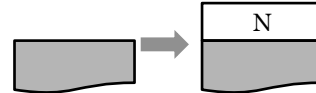
```
v(0) := 1 ;
v(1) := 1 ;
while v(1) << 10 || v(1) === 10 do
    v(0) := v(0) ** v(1) ;
    v(1) := v(1) ++ 1 ;
od
```

- it generates the list of instructions:

```
push(1) | store(v(0)) | push(1) | store(v(1)) |
load(v(1)) | push(10) | lessThan | load(v(1)) | push(10) |
equal | or | jumpOnCond(2) | jump(10) |
load(v(0)) | load(v(1)) | multiply | store(v(0)) |
load(v(1)) | push(1) | add | store(v(1)) |
bjump(17) | quit | cnil
```

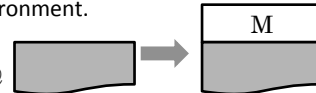
Compilation of Expressions (1)

- A natural number N :
 $\text{eq genForExp}(N) = \text{push}(N) \mid \text{cnil} .$



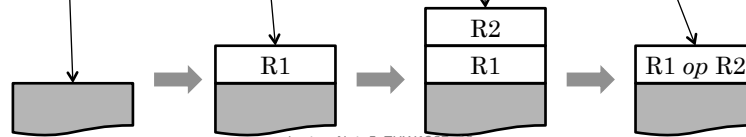
- A variable V :
 $\text{eq genForExp}(V) = \text{load}(V) \mid \text{cnil} .$

Let M be the natural number corresponding to V in the environment.



- $E1 \text{ op } E2$:
 $\text{eq genForExp}(E1 \text{ op } E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{instructionForOp} \mid \text{cnil}) .$

Let $R1$ and $R2$ be the results of evaluating $E1$ and $E2$.



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Compilation of Expressions (2)

- $E1 ** E2$:
 $\text{eq genForExp}(E1 ** E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{multiply} \mid \text{cnil}) .$
- $E1 ++ E2$:
 $\text{eq genForExp}(E1 ++ E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{add} \mid \text{cnil}) .$
- $E1 -- E2$:
 $\text{eq genForExp}(E1 -- E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{minus} \mid \text{cnil}) .$
- $E1 << E2$:
 $\text{eq genForExp}(E1 << E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{lessThan} \mid \text{cnil}) .$
- $E1 >> E2$:
 $\text{eq genForExp}(E1 >> E2) = \text{genForExp}(E1) @ \text{genForExp}(E2) @ (\text{greaterThan} \mid \text{cnil}) .$

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Compilation of Expressions (2)

- $E1 === E2$:
eq genForExp(E1 === E2)
= genForExp(E1) @ genForExp(E2) @ (equal | cnil) .
- $E1 !== E2$:
eq genForExp(E1 !== E2)
= genForExp(E1) @ genForExp(E2) @ (notEqual | cnil) .
- $E1 \&\& E2$:
eq genForExp(E1 \&\& E2)
= genForExp(E1) @ genForExp(E2) @ (and | cnil) .
- $E1 || E2$:
eq genForExp(E1 || E2)
= genForExp(E1) @ genForExp(E2) @ (or | cnil) .

Compilation of Expressions (3)

- $3 ++ 4$:
push(3) | push(4) | add | cnil
- $v(0) ++ 4$:
load(v(0)) | push(4) | add | cnil
- $3 ++ 4 ** 5$:
push(3) | push(4) | push(5) | multiply | add | cnil
- $v(0) ++ v(1) << 10 || v(0) ++ v(1) === 10$:
load(v(0)) | load(v(1)) | add | push(10) | lessThan |
load(v(0)) | load(v(1)) | add | push(10) | equal |
or | cnil

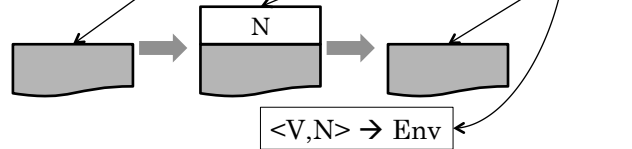
Compilation of Statements (1)

- The empty statement `estm`:
 $\text{eq generator}(\text{estm}, \text{CL}) = \text{CL}$.
- $V := E ; :$
 $\text{eq generator}(V := E ; S, \text{CL})$
 $= \text{generator}(S, \text{CL} @ \text{genForExp}(E) @ (\text{store}(V) | \text{cnil}))$.

The program following the assignment.

The sequence of instructions generated so far.

Let N be the result of evaluating E .



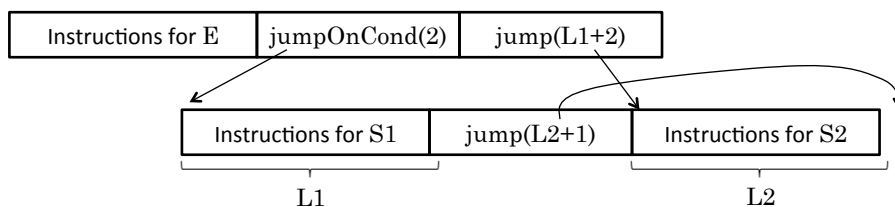
Compilation of Statements (2)

- Examples
 - $v(0) := 1 ; :$
`push(1) | store(v(0)) | cnil`
 - $v(0) := v(0) ++ 4 ; :$
`load(v(0)) | push(4) | add | store(v(0)) | cnil`
 - $v(0) := 3 ++ 4 ** 5 ; :$
`push(3) | push(4) | push(5) | multiply | add | store(v(0)) | cnil`
 - $v(0) := v(0) ++ v(1) << 10 || v(0) ++ v(1) == 10 ; :$
`load(v(0)) | load(v(1)) | add | (push(10) | lessThan | load(v(0)) | load(v(1)) | add | push(10) | equal | or | store(v(0)) | cnil`

Compilation of Statements (3)

- `if E then S1 else S2 fi` :

$$\begin{aligned} & \text{eq generator}(\text{if } E \text{ then } S1 \text{ else } S2 \text{ fi } S, CL) \\ &= \text{generator}(S, CL @ \text{genForExp}(E) \\ & \quad @ (\text{jumpOnCond}(2) | \\ & \quad \quad \text{jump}(\text{len}(\text{generator}(S1, \text{cnil})) + 2) | \text{cnil}) \\ & \quad @ \text{generator}(S1, \text{cnil}) \\ & \quad @ (\text{jump}(\text{len}(\text{generator}(S2, \text{cnil})) + 1) | \text{cnil}) \\ & \quad @ \text{generator}(S2, \text{cnil})). \end{aligned}$$



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Compilation of Statements (4)

- `if 1 then estm else estm fi` :

$$\text{push}(1) | \text{jumpOnCond}(2) | \text{jump}(2) | \text{jump}(1) | \text{cnil}$$
- `if v(0) then v(0) := 0 ; else v(0) := 1 ; fi` :

$$\begin{aligned} & \text{load}(v(0)) | \text{jumpOnCond}(2) | \text{jump}(4) | \\ & \text{push}(0) | \text{store}(v(0)) | \text{jump}(3) | \\ & \text{push}(1) | \text{store}(v(0)) | \text{cnil} \end{aligned}$$
- `if v(0) << v(1) ++ 1 then v(0) := v(0) ++ 1 ; else estm fi` :

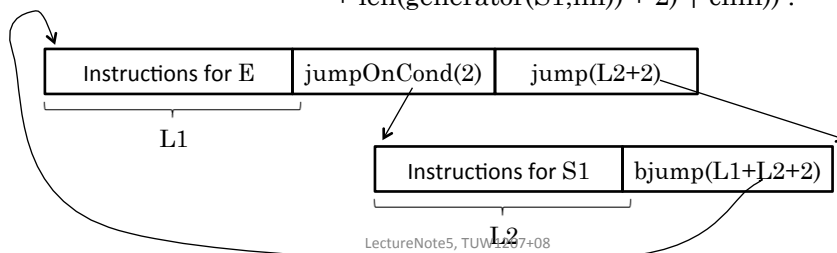
$$\begin{aligned} & \text{load}(v(0)) | \text{load}(v(1)) | \text{push}(1) | \text{add} | \text{lessThan} | \\ & \text{jumpOnCond}(2) | \text{jump}(6) | \\ & \text{load}(v(0)) | \text{push}(1) | \text{add} | \text{store}(v(0)) | \\ & \text{jump}(1) | \text{cnil} \end{aligned}$$

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Compilation of Statements (5)

- `while E do S1 od :`
`eq generator(while E do S1 od S,CL)`
`= generator(S,`
`CL @ genForExp(E)`
`@ (jumpOnCond(2) |`
`jump(len(generator(S1,nil)) + 2) | cnil)`
`@ generator(S1,cnil)`
`@ (bjump(len(genForExp(E))`
`+ len(generator(S1,nil)) + 2) | cnil)) .`



Compilation of Statements (6)

- `while 0 do estm od :`
`push(0) | jumpOnCond(2) | jump(2) | bjump(3) | cnil`
- `while v(0) << 10 do v(1) := v(1) ++ v(0) ; od :`
`load(v(0)) | push(10) | lessThan |`
`jumpOnCond(2) | jump(6) |`
`load(v(1)) | load(v(0)) | add | store(v(1)) | bjump(9) | cnil`
- `while v(0) << 9 || v(0) === 10 do v(1) := v(1) ++ v(0) ; od :`
`load(v(0)) | push(9) | lessThan |`
`load(v(0)) | push(10) | equal | or |`
`jumpOnCond(2) | jump(6) |`
`load(v(1)) | load(v(0)) | add | store(v(1)) | bjump(13) | cnil`

Compilation of Statements (7)

- When the compiler takes the program:

```
v(0) := 1 ; v(1) := 1 ;
while v(1) << 10 || v(1) === 10 do
  v(0) := v(0) ** v(1) ;
  v(1) := v(1) ++ 1 ;
od
```

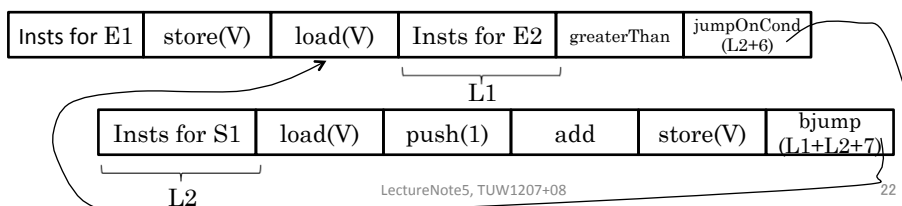
- it generates the list of instructions:

```
push(1) | store(v(0)) |
push(1) | store(v(1)) |
load(v(1)) | push(10) | lessThan |
load(v(1)) | push(10) | equal | or |
jumpOnCond(2) | jump(10) |
load(v(0)) | load(v(1)) | multiply | store(v(0)) |
load(v(1)) | push(1) | add | store(v(1)) |
bjump(17) |
quit | cnil
```

Compilation of Statements (8)

- for V E1 E2 do S1 od :

```
eq generator(for V E1 E2 do S1 od S,CL)
= generator(S, CL @ genForExp(E1)
@ (store(V) | load(V) | cnil)
@ genForExp(E2)
@ (greaterThan |
jumpOnCond(len(generator(S1,cnil)) + 6) | cnil)
@ generator(S1,cnil)
@ (load(V) | push(1) | add | store(V) |
bjump(len(generator(E2)) +
len(generator(S1,cnil)) + 7) | cnil) ) .
```



Compilation of Statements (9)

- for v(0) 0 1 do estm od :
 push(0) | store(v(0)) |
 load(v(0)) | push(1) | greaterThan | jumpOnCond(6) |
 load(v(0)) | push(1) | add | store(v(0)) | bjump(8) | cnil
- for v(0) 0 10 do v(1) := v(1) ++ v(0) ; od :
 push(0) | store(v(0)) |
 load(v(0)) | push(10) | greaterThan | jumpOnCond(10) |
 load(v(1)) | load(v(0)) | add | store(v(1)) |
 load(v(0)) | push(1) | add | store(v(0)) | bjump(12) | cnil
- for v(0) v(1) (v(2) ++ v(1)) do v(3) := v(3) ** v(0) ; od :
 load(v(1)) | store(v(0)) |
 load(v(0)) | load(v(2)) | load(v(1)) | add | greaterThan |
 jumpOnCond(10) |
 load(v(3)) | load(v(0)) | multiply | store(v(3)) |
 load(v(0)) | push(1) | add | store(v(0)) | bjump(14) | cnil

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Compilation of Statements (10)

- When the compiler takes the program:

```
v(0) := 1 ;  
for v(1) 1 10 do  
  v(0) := v(1) ** v(0) ;  
od
```

- it generates the list of instructions:

```
push(1) | store(v(0)) |  
push(1) | store(v(1)) | load(v(1)) | push(10) |  
greaterThan | jumpOnCond(10) |  
load(v(1)) | load(v(0)) | multiply | store(v(0)) |  
load(v(1)) | push(1) | add | store(v(1)) |  
bjump(12) | quit | cnil
```

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Compiler

Function compile:

```
op compile : Stm -> Clist
eq compile(S:Stm)
    = generator(S,cnil) @ (quit | cnil) .
```

Exercises

1. Define a compiler of the calculator in Exercise 1 of Lecture 3 for the virtual machine in Exercise 2 in Lecture 4.