### A Survey of Determinacy of Infinite Games in Second Order Arithmetic,

dedicating to 60's birthday of Professor Tanaka.

Keisuke Yoshii National Institute of Technology, Okinawa College

### Introduction

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Some statements are not provable in second order arithmetic, (or even ZFC). Focus on determinacy on  $Z_2$  and see some related results.

Def. of Infinite Game (i) Let  $A \subseteq \mathbb{N}^{\mathbb{N}}$  be a set in  $\mathcal{C}$ . (ii) Player I and II alternately choose natural

I  $n_0$   $n_2$  ...

numbers as follows: II  $n_1 \dots$ 

(iii) I wins if  $n_0, n_1, n_2 \cdots \in A$ . II wins if not.

(iv) Game on A is determinate if one of the players has a winning strategy.

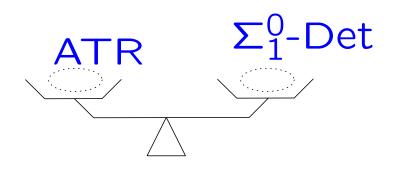
#### Notes

• Difficulties to prove the determinacy of game on *A*, (or, compute the winning strategy)

⇒ Depending on the complexities of set *A*. (e.g. *A* could be  $\Delta_1^0$  (clopen),  $\Sigma_1^0$  (open),  $\Sigma_2^0, \Delta_3^0, \Pi_3^0, \dots \Delta_1^1$  (Borel) .... What determinacy asserts?

Game G<sub>A</sub> is determinate means that one of the players have a winning strategy.
It asserts the existence of real number with certain complexity.

(e.g.)



In early age of reverse mathematics:

•  $ZFC^{-} \not\vdash Borel determinacy (\Delta_1^1-Det)$ . (F.

Friedman, 1971)

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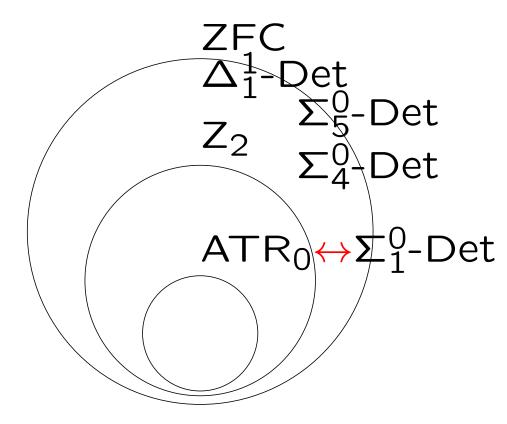
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- ZFC⊢ Borel determinacy. (D. Martin, 1975)

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- $Z_2 \not\vdash \Sigma_4^0$ -Det. (D. Martin, 1974)
- OZFC⊢ Borel determinacy. (D. Martin, 1975)
- And, in 1976 J. Steel showed that one of

the earliest results of Reverse Mathematics: (RCA<sub>0</sub>) ATR $\leftrightarrow \Sigma_1^0$ -Det.



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Defference b/w Boldface and Lightface version.

(Boldface)  $\mathsf{RCA}_0 \vdash \Pi_1^1 - \mathsf{CA} \leftrightarrow \Sigma_1^0 \land \Pi_1^0 - \mathsf{Det}.$ 

### Importances of those theorems (Lightface) (i) $ACA_0 \vdash (\Sigma_1^0 \land \Pi_1^0)$ -Det $\rightarrow \Pi_1^1$ -CA (ii) $ATR_0 \vdash \Pi_1^1$ -CA $\rightarrow (\Sigma_1^0 \land \Pi_1^0)$ -Det Base theory of (ii) can not be weaker than $ATR_0$

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(·.·) Set of sure winning positions for player I can be constructed by  $ATR_0$  with  $\Pi_1^1$ -oracle.

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(e.g.) In the proof of " $\Sigma_1^0 \wedge \Pi_1^0$ -Det  $\rightarrow \Pi_1^1$ -CA", an infinite sequence of natural numbers satisfying  $\Pi_1^1$ -formula will be constructed in the game by players.

Letting players construct sets in the games.

(e.g.) In the proof of " $\Sigma_1^0 \wedge \Pi_1^0$ -Det  $\rightarrow \Pi_1^1$ -CA", For any  $\varphi(n) : \Pi_1^1$ , there exists  $\theta(n, X) : \Delta_1^0$ s.t.

$$\varphi(n) \leftrightarrow \forall f \exists m \theta(n, f[m]).$$

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 $\Sigma_1^{\perp}$ -ID<sub>0</sub>, a new axiom system of Z<sub>2</sub> It asserts the existence of inductively defined pre-well-ordering. (We will See next) Output the comprehension axioms as before. (Some reason?) • Some Varieties of  $\Sigma_1^1$ -ID<sub>0</sub>? (Future studies)

## Def. of $\Sigma_1^1$ -ID<sub>0</sub> >An operator $\Gamma$ is a function from $\Gamma : \mathcal{P}(\mathbb{N}) \to \mathcal{P}(\mathbb{N})$ .

▶ If " $x \in \Gamma(X)$ " is represented by a  $\Sigma_1^1$  formula, then  $\Gamma$  is called  $\Sigma_1^1$  operator. ●  $\Sigma_1^1$ -ID: for any  $\Sigma_1^1$ -operator  $\Gamma$ , there exists pre-wellordering  $V \subset \mathbb{N} \times \mathbb{N}$  s.t. the following holds:

### **Def. Cont** • $\Sigma_1^1$ -ID: for any $\Sigma_1^1$ -operator $\Gamma$ , there exists pre-wellordering $V \subset \mathbb{N} \times \mathbb{N}$ s.t. the following holds:

$$\forall x \in F(V_x = \Gamma(V_{\leq x}) \cup V_{\leq x}), \\ \triangleright \Gamma(F) \subset F.$$

where

$$V_x = \{ y \in F : y \leq_V x \},\$$
  
$$V_{< x} = \{ y \in F : y <_V x \}, F = \{ x : x \leq_V x \}.$$

### Image of $\Sigma_1^1$ -ID $\triangleright$ Apply $\Gamma$ to $\emptyset$

### Image of $\Sigma_1^1$ -ID > Apply $\Gamma$ to $\Gamma(\emptyset)$ and take the union:

### $\Gamma(\emptyset) \cup \Gamma(\Gamma(\emptyset))$

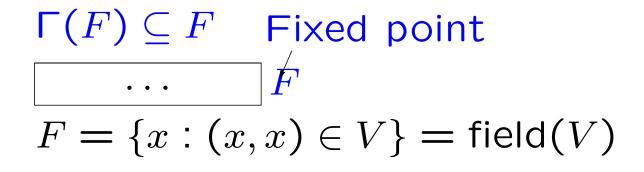


### Image of $\Sigma_1^1$ -ID $\triangleright$ Keep doing this until ...

### $\Gamma(\emptyset) \cup \Gamma(\Gamma(\emptyset)) \cup \Gamma(\Gamma(\emptyset)) \cup \Gamma(\Gamma(\emptyset)))$



### Image of $\Sigma_1^1$ -ID $\triangleright$ until the fixed point:



Σ<sup>1</sup><sub>1</sub>-ID<sub>0</sub>, a new axiom system of Z<sub>2</sub>
It asserts the existence of inductively defined pre-well-ordering. (We will See next)
Different with comprehension axioms as before. (Some?)
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•  $C\text{-ID}_0 \rightarrow C\text{-CA}_0$ 

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⇒ Logical equivalence can not be obtained by CA for  $\Sigma_2^0$ -Det.

Even Borel det. doesn't deduce  $\Delta_2^{\perp}$ -CA (MedSalem, Tanaka)[2007] (:) By  $\beta$ -model reflection and 2nd Imcomp.  $\cdot \Delta_1^1$ -Det+ $\Sigma_2^1$ -DC<sub>0</sub>  $\vdash$  "exist. of c.c.  $\beta$ -model M of  $\Delta_1^1$ -Det." (by  $\Sigma_2^1$ -DC<sub>0</sub>  $\vdash \Sigma_4^1$ -RFN<sup>1</sup>)  $\cdot \mathsf{M} \models \Delta_1^1 - \mathsf{Det} + \Delta_2^1 - \mathsf{CA}_0 (\mathsf{By} \ \Delta_1^1 - \mathsf{Det} \vdash \Delta_2^1 - \mathsf{CA}_0)$  $\cdot M \models \Delta_1^1 - \text{Det} + \Sigma_2^1 - \text{DC}_0 \text{ (by } \Sigma_2^1 - \text{DC}_0 \leftrightarrow \Sigma_\infty^1 - \Sigma_\infty^1 - \Sigma_\infty^1 + \Sigma_\infty^1 - \Sigma_$  $IND + \Delta_2^1 - CA_0$  $\square$ More details, see M.T[2007], Simpson[2009], VII

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# Some Varieties of $\Sigma_1^{\perp}$ -ID<sub>0</sub> • $\Sigma_1^1$ -ID $\leftrightarrow \Sigma_2^0$ -Det (1991) • $\Sigma_1^1$ -IDTR<sub>0</sub> $\leftrightarrow \Delta((\Sigma_2^0)_2)$ -Det (2012) • $[\Sigma_1^1]^2$ -ID<sub>0</sub> $\leftrightarrow$ $(\Sigma_2^0)_2$ -Det (2008) • $[\Sigma_1^1]^k$ -ID<sub>0</sub> $\leftrightarrow$ $(\Sigma_2^0)_k$ -Det (2008) • $[\Sigma_1^1]^k$ -IDTR<sub>0</sub> $\leftrightarrow \Delta((\Sigma_2^0)_{k+1})$ -Det (2012) (k > 3)ł

•  $(\Pi_3^1 - \mathsf{TI}_0)[\Sigma_1^1]^{\mathsf{TR}} - \mathrm{ID}_0 \leftrightarrow \Delta_3^0 - \mathrm{Det} (2008)$ 

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 $\varphi$  is a  $\Delta(\mathcal{C})$  formula if

$$\varphi \leftrightarrow \psi \wedge \neg \varphi \leftrightarrow \eta$$

,where  $\psi, \eta \in C$ . For any k > 1,  $(\Sigma_n^0)_k = \Sigma_n^0 \wedge (\Pi_n^0)_{k-1}$ . e.g.  $(\Sigma_2^0)_2 = \Sigma_2^0 \wedge \Pi_2^0$ 

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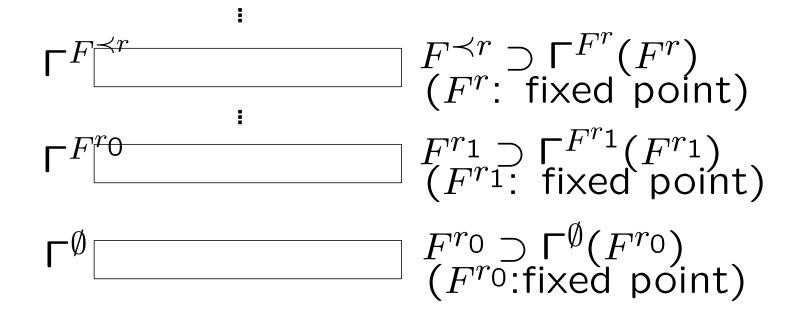
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•  $[\Sigma_1^1]^k ext{-}ID_0 \leftrightarrow (\Sigma_2^0)_k ext{-}Det (2008)$ •  $[\Sigma_1^1]^k ext{-}IDTR_0 \leftrightarrow \Delta((\Sigma_2^0)_{k+1}) ext{-}Det (2012)$ :  $(k \ge 3)$ •  $(\Pi_3^1 ext{-}TI_0)[\Sigma_1^1]^{\mathsf{TR}} ext{-}ID_0 \leftrightarrow \Delta_3^0 ext{-}Det (2008)$ 

### Image of $\Sigma_1^1$ -IDTR<sub>0</sub>



where  $F^{\prec r} = \cup \{F^{r'} : r' \prec r\}]$ 

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- (Det\*: determinacy on Cantor sapce).
- $\triangleright$  WKL $\leftrightarrow \Sigma_1^0$ -Det\*
- ▷ ATR  $\leftrightarrow \Sigma_1^0$ -Det  $\leftrightarrow \Sigma_2^0$ -Det\* ▷  $[\Sigma_1^1]^k$ -ID  $\leftrightarrow (\Sigma_2^0)_k \leftrightarrow (\Sigma_2^0)_{k-1}$ -Det\* (for  $k \ge 2$ )

The limit of determinacy in second order arithmetic: (Montalbán, Shore, 2012)
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• The limit of determinacy in second order arithmetic: (Monralbán, Shore,2012) (i)  $\Pi^{1}_{m+2}$ -CA  $\rightarrow (\Sigma^{0}_{3})_{m}$ -Det  $(m \ge 1)$ (ii)  $\Delta^{1}_{m+2}$ -CA  $\not\rightarrow (\Sigma^{0}_{3})_{m}$ -Det  $(m \ge 1)$ 

• The limit of determinacy in second order arithmetic: (Monralbán, Shore, 2012) (i)  $\Pi^1_{m+2}$ -CA  $\rightarrow (\Sigma^0_3)_m$ -Det  $(m \ge 1)$ (ii) $\Delta^1_{m+2}$ -CA  $\not\rightarrow (\Sigma^0_3)_m$ -Det  $(m \ge 1)$ > Note that the reversal of (i) does not hold since  $\Delta^1_1$ -Det  $\not\rightarrow \Delta^1_2$ -CA

Set Theory

Open determinacy for class games for Con(ZFC), (Hamkins)

 $\triangleright$ 

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## Set Theory

- Open determinacy for class games for Con(ZFC) (Hamkins)
- ▷ A Mathematician's Year in Japan, by Hamkins
- ▷ Characterization of  $\Sigma_4^0$ -Det (S. Hatchman)

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Then, ...

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Then, thank very much Professor Tanaka and.