

# Reverse Mathematics and Equilibria of Continuous Games

a joint work with NingNing Peng and Weiguang Peng

Takeshi Yamazaki

Mathematical Institute, Tohoku University,

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### *Definition of continuous games*

**Definition 1** (RCA<sub>0</sub>). *A continuous game (with  $n$  players) is a  $2n$ -tuple  $\langle A_1 \dots, A_n; f_1 \dots, f_n \rangle$  where*

- 1. each  $A_i$  is a nonempty compact metric space,*
- 2. each  $f_i$  is a continuous function form  $\prod A_i \rightarrow \mathbb{R}$  with a modulus of uniform continuity.*

From now, we'd like to consider an extension of Nash's theorem by Glicksberg on continuous games.

Fix a continuous game  $G = \langle A_1 \dots, A_n; f_1 \dots, f_n \rangle$ .

We call a probability measure  $\mu$  over  $A_i$ , a mixed strategy of player  $i$ . Recall that in Reverse Mathematics, a probability measure  $\mu$  over a compact metric space  $X$  is defined as a positive linear functional on  $C(X)$  with  $\mu(1)$  where  $1$  is a constant function whose output is  $1$ .

Of course, we define a mixed profile by a finite sequence of mixed strategies for all players. But to define an expected value of a pay-off function, we have to think of it as the product of measures.

Recall again that  $C(X)$  is defined as a complete separable metric space given from a countable set of "polynomials" on  $X$  with sup-norm. In fact, to have the same space, we can use diverse notions of "polynomial". So, to define a product measure, we use the following version of S-W theorem.

**Lemma 1** (RCA<sub>0</sub>). *Let  $X$  be a compact metric space with an evidence  $\langle \langle x_{ij} : i \leq n_j \rangle : j \in \mathbb{N} \rangle$  for compactness. Assume that a sequence  $\langle p_n : n \in \mathbb{N} \rangle$  of  $C(X)$  satisfies the following: (1)  $p_0 = 1$ ,*

*(2)  $\{p_n : n \in \mathbb{N}\}$  is closed under  $+$ ,  $\cdot$  and  $r \cdot$  for any  $r \in \mathbb{Q}$ .*

*(3) there exists a function  $h : \mathbb{N}^2 \rightarrow \mathbb{N}$  such that for each  $i$  and  $j$ ,  $p_{h(i,j)}$  satisfies  $0 \leq p_{h(i,j)} \leq 1$ ,  $p_{h(i,j)} = 1$  on  $\overline{B(x_{ij}, 2^{-j})}$  and  $p_{h(i,j)} = 0$  on  $B(x_{ij}, 2^{-j+1})^c$ .*

*Then, there exists an effective correspondence between  $\hat{A} = \langle \mathbb{N}, d_A \rangle$  and  $C(X)$  where  $d_A(n, m) = \|p_n - p_m\|_\infty$*

Let  $P_i$  be the “polynomials” over  $\mathbb{Q}$  from basic functions on  $A_i$ . Then  $P = \{\sum_{j < m} p_{1j} \cdots p_{nj} : \text{each } p_{ij} \in P_i\}$  can be regarded as satisfying the conditions of Lemma 1 for  $C(A)$ . ( $A = \prod A_i$ )

So we may assume that  $C(A) = \langle P, \|\cdot\|_\infty \rangle$ . Then, given a sequence  $\langle \mu_1, \dots, \mu_n \rangle$  such that each  $\mu_i$  is a probability measure on  $A_i$ , we can define the product measure  $\mu$  on  $A$  by

$$\mu(p_1 \cdots p_n) = \prod_i \mu_i(p_i),$$

for any  $p_i \in P_i$  with the canonical extension.

We identify a mixed profile as a product measure given from it.

Let  $\mu$  be a mixed profile of  $G$ . Then we define the expected value of the payoff functions  $f_i$  by

$$f_i(\mu) = \int f_i d\mu \text{ (i.e., } = \mu(f_i))$$

.

Note that  $\mu_{-i} = \langle \mu_j : j \neq i \rangle$  is a vector of mixed strategies for all players except  $i$ . We sometimes write  $(\mu_i, \mu_{-i})$  instead of  $\mu$ .

**Definition 2** ( $\text{RCA}_0$ ). *A mixed Nash equilibrium of a continuous game  $\langle A_1, \dots, A_n; f_1, \dots, f_n \rangle$  is a mixed profile  $\mu^*$  such that for all  $i$ ,  $1 \leq i \leq n$ ,  $f_i(\mu^*) \geq f_i(\mu_i, \mu_{-i}^*)$  for any mixed strategy  $\mu_i$  of player  $i$ .*

## *Glicksberg's theorem*

Glicksberg's theorem is proved in  $ACA_0$ .

**Theorem 2** ( $ACA_0$ ). *Every continuous game has a mixed Nash equilibrium.*

Our proof is essentially based on Ozdaglar's proof on Glicksberg's theorem in his lecture notes. In the proof, the following fact is the most important.

**Theorem 3.** *The following assertions are pairwise equivalent over  $RCA_0$ .*

(1)  $ACA_0$ .

(2) *Any sequence  $\langle \mu_n : n \in \mathbb{N} \rangle$  of probability measures on a compact metric space  $X$  has a weak convergent subsequence  $\langle \mu_{n_k} : k \in \mathbb{N} \rangle$ , that is, there exists a probability measure  $\mu$  such that  $\mu(f) = \lim_{k \rightarrow \infty} \mu_{n_k}(f)$  for all  $f \in C(X)$ .*

**proof.** (1)  $\rightarrow$  (2). Let  $X$  be a compact metric space and  $C(X)$  coded by  $(P, \|\cdot\|_\infty)$ . Let  $\langle p_i : i \in \mathbb{N} \rangle$  be an enumeration of the elements of  $P$ . For each  $n \in \mathbb{N}$  and  $\sigma \in 2^{<\mathbb{N}}$ , we define a closed interval  $I_\sigma^i$  by

$$I_\sigma^i = \left[ \left( \sum_{j < \text{lh}(\sigma)} \frac{\sigma(j)}{2^j} - 1 \right) \|p_i\|_\infty, \left( \sum_{j < \text{lh}(\sigma)} \frac{\sigma(j)}{2^j} - 1 + \frac{1}{2^{\text{lh}(\sigma)-1}} \right) \|p_i\|_\infty \right].$$

Take any sequence  $\langle \mu_n : n \in \mathbb{N} \rangle$  of probability measures on  $X$ . Then we have sequences  $\langle n_k : k \in \mathbb{N} \rangle$  and  $\langle \sigma_k^i : i, k \in \mathbb{N}, i \leq k \rangle$  such that

$$\text{lh}(\sigma_k^i) = k, \sigma_{k+1}^i \succeq \sigma_k^i$$

and  $\mu_{n_k}(p_i) \in I_{\sigma_k^i}^i$  for all  $i \leq k$ .



Then,  $|I_{\sigma_k^i}^i| = 2^{-k+1} \|p_i\|_\infty$  and  $I_{\sigma_{k+1}^i}^i \subset I_{\sigma_k^i}^i$  for all  $i \leq k$ .

So, we can get a probability measure  $\mu$  by

$$\mu(p) = \lim_{k \rightarrow \infty} \mu_{n_k}(p), \text{ for all } p \in P.$$

Thus, we have a sequence  $\langle \mu_{n_k} : k \in \mathbb{N} \rangle$  which converges to  $\mu$ .

**The converse direction (2)  $\rightarrow$  (1) is easy.**

Let  $\langle c_n : n \in \mathbb{N} \rangle$  be a sequence on  $[0, 1]$ . We may only show that it has a convergent subsequence.

Define a probability measure  $\mu_n$  on  $[0, 1]$  by  $\mu_n(f) = f(c_n)$ .

From our assumption (2), it has a subsequence

$$\mu_{n_k} \rightarrow \mu$$

for some probability measure  $\mu$ . Take  $f \in C([0, 1])$  with  $f(x) = x$ .

Then,  $\langle \mu_{n_k}(f) : k \in \mathbb{N} \rangle$  is a subsequence of  $\langle c_n : n \in \mathbb{N} \rangle$  which converges to  $\mu(f)$ .  $\square$ .

The rest of the proof of Theorem 2 is like this: Take any continuous game  $G = \langle A_1 \dots, A_n; f_1 \dots, f_n \rangle$ .

1. By compactness of  $A = \prod_i A_i$ , we can find a sequence of “essentially finite games”  $\langle G_m : m \in \mathbb{N} \rangle$  such that  $|f_i^m(\sigma) - f_i(\sigma)| \leq 2^{-m}$ , where  $G_m = \langle A_1 \dots, A_n; f_1^m \dots, f_n^m \rangle$
2. By the sequence version of Nash’s theorem, there exists a sequence  $\langle \mu_m : m \in \mathbb{N} \rangle$  of mixed strategies such that each  $\mu_m$  is Nash equilibrium of  $G_m$ .
3. Using Theorem 3, pick up a subsequence  $\langle \mu_{m_k} : k \in \mathbb{N} \rangle$  converging to  $\mu$ .
4. By 1,  $f_i(a_i, (\mu_{m_k})_{-i}) \leq f_i(\mu_{m_k}) + 2^{-m_k+1}$  for all  $a_i \in A_i$ . As  $k \rightarrow \infty$ ,  $f_i(a_i, (\mu)_{-i}) \leq f_i(\mu)$ , that is,  $\mu$  is a Nash equilibrium.

## Remark.

1. An “essentially finite game” may not be continuous. Indeed, it is coded by using simple functions with arithmetical sets in  $ACA_0$ .
2. The sequence version of Nash’s theorem can be proved in  $WKL_0$ . But, when we apply it to “essentially finite” games, so far, we require  $ACA_0$  because we have to construct a probability measure on  $A$  from a mixed strategies of finite games on the way of the proof.
3. I don’t know if Glicksberg’s theorem implies  $ACA_0$  or not.

Recall our version of Browder Fixed point theorem:

**Theorem 4.** *The following assertion is equivalent to  $WKL_0$  over  $RCA_0$ . Let  $\hat{A}$  be a nonempty compact convex “subspace” of  $\mathbb{R}^n$ . Let an open subset  $T$  of  $\hat{A} \times \hat{A}$  satisfy the following conditions:*

- (i) for any  $x \in \hat{A}$ , there exists  $y \in \hat{A}$  such that  $(x, y) \in T$ ,*
- (ii) for any  $x, y_1, y_2 \in \hat{A}$  and any positive reals  $r_1$  and  $r_2$ , if  $(x, y_1), (x, y_2) \in T$  and  $r_1 + r_2 = 1$ , then  $(x, r_1y_1 + r_2y_2) \in T$ .*

*Then there exists  $x \in \hat{A}$  such that  $(x, x) \in T$ .*

**Definition 3** (RCA<sub>0</sub>). *A pure Nash equilibrium of a continuous game  $\langle A_1, \dots, A_n; f_1 \dots, f_n \rangle$  is a profile  $a^* \in A$  such that for all  $i$ ,  $1 \leq i \leq n$ ,  $f_i(a^*) \geq f_i(a_i, a_{-i}^*)$  for all  $a_i \in A_i$ .*

A continuous function  $f : [0, 1] \rightarrow \mathbb{R}$  is concave if for any  $x, y \in [0, 1]$   $tf(x) + (1 - t)f(y) \leq f(tx + (1 - t)y)$ .

Then, we have the following result.

**Theorem 5.** *The following is provable in WKLO. Let  $G = \langle [0, 1] \dots, [0, 1]; f_1 \dots, f_n \rangle$  be a continuous game. Assume that any  $f_i(\cdot, x_{-i})$  is concave. Then there exists a pure Nash equilibrium.*

## References

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Thank you so much!!