

A Galois connection between basic covers and binary positivities

Francesco Ciraulo
(j.w.w. Giovanni Sambin)



Department of Mathematics and Computer Science
University of PALERMO (Italy)
ciraulo@math.unipa.it
www.math.unipa.it/~ciraulo

Workshop on Constructive Aspects of Logic and Mathematics
Kanazawa, 8 - 12 March 2010

Overview

1 BASIC topologies

- basic covers = closure operators = saturations
- binary positivities = interior operators = reductions
- “compatibility”

Overview

1 BASIC topologies

- basic covers = closure operators = saturations
- binary positivities = interior operators = reductions
- “compatibility”

2 DENSE and THICK basic topologies

- dense \approx representable (= pointwise definable)
- thick \approx generated (inductively-coinductively)

Overview

1 BASIC topologies

- **basic covers** = closure operators = saturations
- **binary positivities** = interior operators = reductions
- “compatibility”

2 DENSE and THICK basic topologies

- **dense** \approx representable (= pointwise definable)
- **thick** \approx generated (inductively-coinductively)

3 Properties

- **Galois** connection between covers and positivities
- **Dense** and **thick** as **(co)reflections**
- **Adjunction** between dense and thick

Part I

Basic topologies

Basic covers

a **COVER** relation over a set S is

an (infinitary) relation \triangleleft between elements $a, b, \dots \in S$ and subsets $U, V, \dots \subseteq S$, such that:

$$\underbrace{\frac{a \in U}{a \triangleleft U} \quad \frac{a \triangleleft U \quad \forall u (u \in U \Rightarrow u \triangleleft V)}{a \triangleleft V}}_{\text{BASIC COVER}} \quad \frac{a \triangleleft U \quad a \triangleleft V}{a \triangleleft U \dots V} \quad \dots$$

Saturations (closure operators)

Basic cover \simeq saturation (closure operator)

$$\begin{array}{l} a \triangleleft U \iff \mathcal{A}(U) \\ a \in \mathcal{A}(U) \quad \{a \in S \mid a \triangleleft U\} \end{array}$$

a **SATURATION** is

$$\mathcal{A} : \mathcal{P}(S) \rightarrow \mathcal{P}(S)$$

$$U \subseteq \mathcal{A}(U) \quad \mathcal{A}\mathcal{A} = \mathcal{A} \quad \frac{U \subseteq V}{\mathcal{A}(U) \subseteq \mathcal{A}(V)}$$

Binary positivities and Reductions (interior operators)

$$\times : S \times \mathcal{P}(S) \rightarrow Prop$$

$$a \times U$$

$$\{a \in S \mid a \times U\}$$

$$\frac{a \times U}{a \in U}$$

$$\frac{a \times U \quad \forall v (v \times U \Rightarrow v \in V)}{a \times V}$$

$$\mathcal{J} : \mathcal{P}(S) \rightarrow \mathcal{P}(S)$$

$$a \in \mathcal{J}(U)$$

$$\mathcal{J}(U)$$

$$\mathcal{J}(U) \subseteq U$$

$$\frac{\mathcal{J}(U) \subseteq V}{\mathcal{J}(U) \subseteq \mathcal{J}(V)}$$

$$\mathcal{J}(U) \subseteq U$$

$$\mathcal{J}\mathcal{J} = \mathcal{J}$$

$$\frac{U \subseteq V}{\mathcal{J}(U) \subseteq \mathcal{J}(V)}$$

Basic (formal) topologies

Covers and positivities linked by COMPATIBILITY
(saturations) (reductions)

a **BASIC TOPOLOGY** is

$(S, \mathcal{A}, \mathcal{J})$

$(S, \triangleleft, \times)$

(set, *saturation*, *reduction*)

(set, *cover*, *positivity*)

+

compatibility :

$$\frac{\mathcal{A}(U) \text{ } \mathcal{J}(V)}{U \text{ } \mathcal{J}(V)} \quad 1$$

$$\frac{a \triangleleft U \quad a \times V}{(\exists u \in U)(u \times V)}$$

¹Where: $(U \text{ } V) \stackrel{\text{def}}{\iff} (U \cap V \text{ is inhabited})$.

What's next?

Two ways to construct a basic topology:

- inductive-coinductive generation (Martin-Löf & Sambin)
- representation via operators induced by a relation (Basic Picture)

Aim of the talk : to generalise these two methods and show that they are “DUAL” (in some precise sense).

Questions we want to address : *Is \times always determined by \triangleleft ?
(and to answer) When does this happen?
What about the converse?*

Part II

Thick basic topologies

inductively-coinductively generated
basic topologies

Inductive generation of \triangleleft

Problem : given some (= a set-indexed family of) “axioms”
define the **least** \triangleleft satisfying the axioms.

Solution: inductive generation

For any (set-indexed) family of axioms $a \prec U_i$, define \triangleleft by:

$$\boxed{\frac{a \in V}{a \triangleleft V}} \quad \& \quad \boxed{\frac{a \prec U_i \quad (\forall x \in U_i)(x \triangleleft V)}{a \triangleleft V}} \quad \&$$

$$\boxed{\frac{a \triangleleft V \quad V \subseteq P \quad \forall i(U_i \subseteq P \Rightarrow a \in P)}{a \in P}} \quad (\text{induction})$$

\triangleleft is *the least cover* such that:

$$\frac{a \prec U}{a \triangleleft U}$$

Generating positivity by co-induction

(Martin-Löf & Sambin)

$$\boxed{\frac{a \times V}{a \in V}} \quad \& \quad \boxed{\frac{a \prec U_i \quad a \times V}{(\exists x \in U_i)(x \times V)}} \quad \&$$

$$\boxed{\frac{a \in P \quad P \subseteq V \quad \forall i(a \in P \Rightarrow U_i \not\prec P)}{a \times V} \quad (\text{coinduction})}$$

\times is *the greatest positivity* which is **compatible** with \triangleleft

$$\boxed{\frac{a \in V}{a \triangleleft V}}$$

$$\boxed{\frac{a \prec U_i \quad (\forall x \in U_i)(x \triangleleft V)}{a \triangleleft V}}$$

$$\boxed{\frac{a \triangleleft V \quad V \subseteq P \quad \forall i(U_i \subseteq P \Rightarrow a \in P)}{a \in P}}$$

... with operators...

(Set-indexed) family of axioms \rightsquigarrow Operator $\mathcal{F} : \mathcal{P}(S) \rightarrow \mathcal{P}(S)$
 $a \prec U \rightsquigarrow a \in \mathcal{F}(U)$

$$\mathcal{A}(V) \stackrel{\text{def}}{=} \bigcap \left\{ P \mid V \subseteq P \text{ and } \frac{X \subseteq P}{\mathcal{F}(X) \subseteq P} \text{ (for all } X) \right\}$$

\mathcal{A} is the least saturation which contains \mathcal{F}

$$\mathcal{J}(V) = \bigcup \left\{ P \mid P \subseteq V \text{ and } \frac{\mathcal{F}(X) \not\subseteq P}{X \not\subseteq P} \text{ (for all } X) \right\}$$

\mathcal{J} is the greatest reduction which is compatible with \mathcal{A}

The greatest reduction compatible with a given saturation

(Impredicatively)

Let \mathcal{A} be a saturation; then:

(impredicative)

$$\mathbf{J}(\mathcal{A})(V) \stackrel{\text{def}}{=} \bigcup \left\{ L \mid L \subseteq V \text{ and } \frac{\mathcal{A}(X) \not\subseteq L}{X \not\subseteq L} \text{ (for all } X) \right\}$$

is *the greatest reduction* compatible with \mathcal{A}

Classically: $\mathbf{J}(\mathcal{A})(U) = \neg\mathcal{A}(\neg U)$.

Intuitionistically, $\neg\mathcal{A}\neg$ is not a reduction.²

²As for a counterexample, take \mathcal{A} to be the identity operator. 

Thick basic topologies

A basic topology $(S, \mathcal{A}, \mathcal{J})$ is:

generated : if \mathcal{A} and \mathcal{J} are generated **inductively** and **coinductively**, respectively, by the **same** axioms;

thick : if $\mathcal{J} = \mathbf{J}(\mathcal{A})$.

GENERATED \implies **THICK**

Impredicatively the two notions coincide.

One saturation, many reductions...

$$(S, \mathcal{A}, \dots)$$

Problem: *In how many ways can we fill the spaces?*

The greatest solution: $\mathbf{J}(\mathcal{A})$

The least solution: \perp , where $\perp(V) = \emptyset$

Each \mathcal{J} , a different class of formal points.
subspace

$$\underbrace{(S, \mathcal{A}, \perp)}_{\text{no points}}$$

$$\underbrace{(S, \mathcal{A}, \dots) \quad (S, \mathcal{A}, \dots) \quad (S, \mathcal{A}, \dots)}_{\text{proper subspaces}}$$

$$\underbrace{(S, \mathcal{A}, \mathbf{J}(\mathcal{A}))}_{\text{the whole space}}$$

Part III

Dense basic topologies

When points form a set.

Representable basic topologies

(see Sambin's Basic Picture. . .)

Another way to construct a (basic) topology:

Idea : given $\alpha_i \subseteq S (i \in I)$ (the future formal points)
define a basic topology “pointwise”:

$$a \triangleleft U \stackrel{\text{def}}{\Leftrightarrow} (\forall i \in I)(a \in \alpha_i \Rightarrow \alpha_i \not\subseteq U)$$

$$a \times U \stackrel{\text{def}}{\Leftrightarrow} (\exists i \in I)(a \in \alpha_i \ \& \ \alpha_i \subseteq U)$$

Intuitively³ : this is the **least** basic topology
which has the α_i 's as points.

³The exact statement require some extra condition.

Sambin's Basic Pairs

Start from $\Vdash: X \times S \rightarrow Prop$ (an arbitrary binary **relation**) and define the following four operators:

$$\mathcal{P}(X) \rightleftharpoons \mathcal{P}(S)$$

$$\begin{array}{lcl} \{x \mid (\exists u)(u \in U \ \& \ x \Vdash u)\} = \text{ext}(U) & \leftarrow & U \\ & D \rightarrow & \diamond D = \{a \mid (\exists d)(d \in D \ \& \ d \Vdash a)\} \\ \{x \mid (\forall u)(x \Vdash u \Rightarrow u \in U)\} = \text{rest}(U) & \leftarrow & U \\ & D \rightarrow & \square D = \{a \mid (\forall d)(d \Vdash a \Rightarrow d \in D)\} \end{array}$$

Since $\text{ext} \dashv \square$ and $\diamond \dashv \text{rest}$ we get that $\mathcal{A} = \square \text{ext}$ is a saturation (the monade of the former adjunction) and $\mathcal{J} = \diamond \text{rest}$ is a reduction (the comonade of the latter adjunction). Moreover, they are compatible, that is: $(S, \mathcal{A}, \mathcal{J}) = (S, \square \text{ext}, \diamond \text{rest})$ is a basic topology (*represented* by the relation \Vdash).

Equivalently...

$$a \in \mathcal{A}(U) \Leftrightarrow a \in \square(\text{ext}(U)) \Leftrightarrow (\forall x \in X)(a \in \diamond\{x\} \Rightarrow \diamond\{x\} \not\subseteq U)$$

$$a \in \mathcal{J}(U) \Leftrightarrow a \in \diamond(\text{rest}(U)) \Leftrightarrow (\exists x \in X)(a \in \diamond\{x\} \ \& \ \diamond\{x\} \subseteq U)$$

So, given a set-indexed family $\overbrace{\{F_i \subseteq S \mid i \in I\}}^{\{\diamond\{x\} \subseteq S \mid x \in X\}}$, define:

$$a \in \mathcal{A}(U) \stackrel{\text{def}}{\Leftrightarrow} (\forall i \in I)(a \in F_i \Rightarrow F_i \not\subseteq U)$$

$$a \in \mathcal{J}(U) \stackrel{\text{def}}{\Leftrightarrow} (\exists i \in I)(a \in F_i \ \& \ F_i \subseteq U)$$

Properties:

- $(S, \mathcal{A}, \mathcal{J})$ is a basic topology (\mathcal{A} and \mathcal{J} are compatible)
- \mathcal{J} is *the least reduction* s.t. each F_i is **formal closed** (= \mathcal{J} -fixed)
- \mathcal{A} is *the greatest saturation compatible with \mathcal{J}*

The greatest cover compatible with a given positivity (Impredicatively)

Impredicatively, every \mathcal{J} is “representable”: take the family $\{\mathcal{J}(Z) \mid Z \subseteq S\}$ and check:

$$a \in \mathcal{J}(U) \Leftrightarrow (\exists Z \subseteq S)(a \in \mathcal{J}(Z) \ \& \ \mathcal{J}(Z) \subseteq U)$$

(impredicative)

For any reduction \mathcal{J} there exists the greatest saturation compatible with it:

$$a \in \mathbf{A}(\mathcal{J})(U) \stackrel{\text{def}}{\Leftrightarrow} (\forall Z \subseteq S)(a \in \mathcal{J}Z \Rightarrow U \not\subseteq \mathcal{J}Z)$$

Classically: $\mathbf{A}(\mathcal{J})(U) = \neg \mathcal{J}(\neg U)$.

Intuitionistically: $\neg \mathcal{J} \neg$ is a saturation, BUT not compatible with \mathcal{J} .

Dense basic topologies

Definition

A basic topology $(S, \mathcal{A}, \mathcal{J})$ is **dense** if $\mathcal{A} = \mathbf{A}(\mathcal{J})$.

Proposition

$$\text{REPRESENTABLE} \implies \text{DENSE}$$

(and they coincide impredicatively)

Summing up

$\mathbf{J}(\mathcal{A})$
the greatest reduction
compatible with
 \mathcal{A}

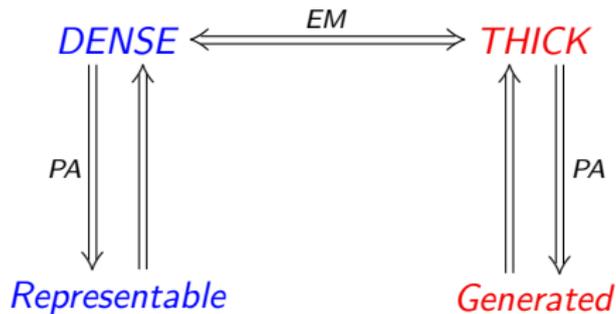
$\mathbf{A}(\mathcal{J})$
the greatest saturation
compatible with
 \mathcal{J}

IMPREDICATIVELY: they ALWAYS exist!

Predicatively, they exist if...

... if \mathcal{A} is generated
(by axioms)

... if the \mathcal{J} -closed subsets
are set-based
(obtained by union
from a set-indexed subfamily)



$$\frac{\text{Dense (Representable)} \longleftarrow \text{Generated}}{\text{Markov}}$$

$$\text{Representable} \xRightarrow{?} \text{Thick (Generated)}$$

Part IV

The Galois connection

... between saturations and reductions
(on the same set S)

TFAE

- \mathcal{A} and \mathcal{J} are compatible
- $\mathcal{A} \subseteq \mathbf{A}(\mathcal{J})$
- $\mathcal{J} \subseteq \mathbf{J}(\mathcal{A})$

$$\mathcal{J} \subseteq \mathbf{J}(\mathcal{A}) \iff \mathcal{A} \subseteq \mathbf{A}(\mathcal{J})$$

$$\{\text{reductions over } S\} \begin{array}{c} \xrightarrow{\mathbf{A}} \\ \xleftarrow{\mathbf{J}} \end{array} \{\text{saturation s over } S\}$$

is a (antitone) Galois connection
(or dual adjunction)

For any basic topology $\mathcal{S} = (S, \mathcal{A}, \mathcal{J})$ we put:

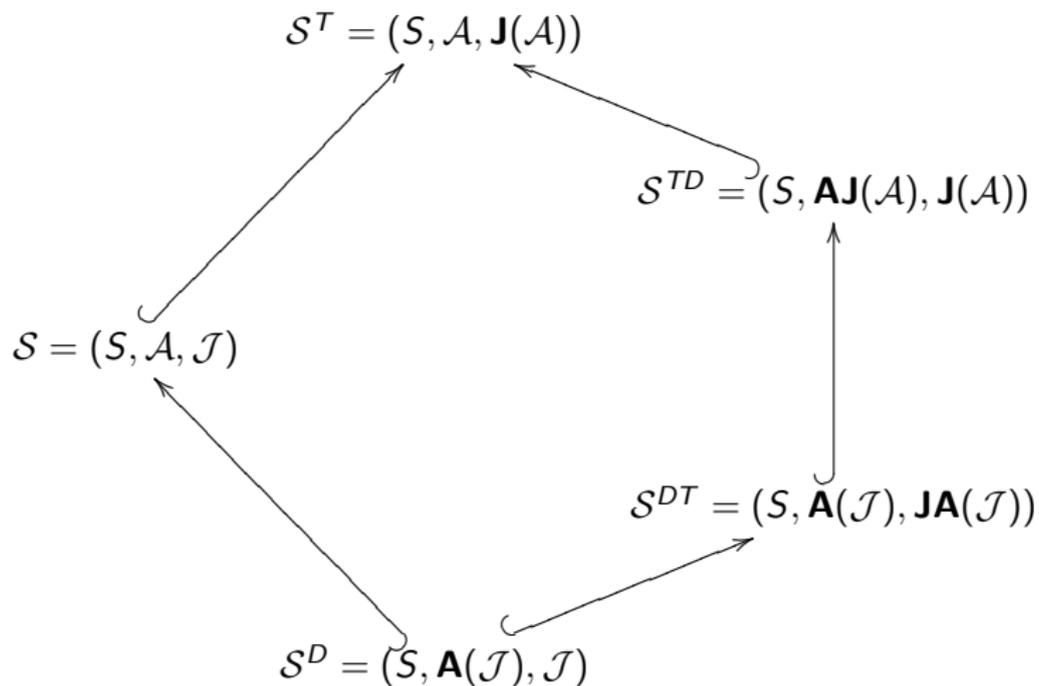
(“densification”) $\mathcal{S}^D = (S, \mathbf{A}(\mathcal{J}), \mathcal{J})$

the greatest *dense* (*representable*) basic topology
which is contained in \mathcal{S}

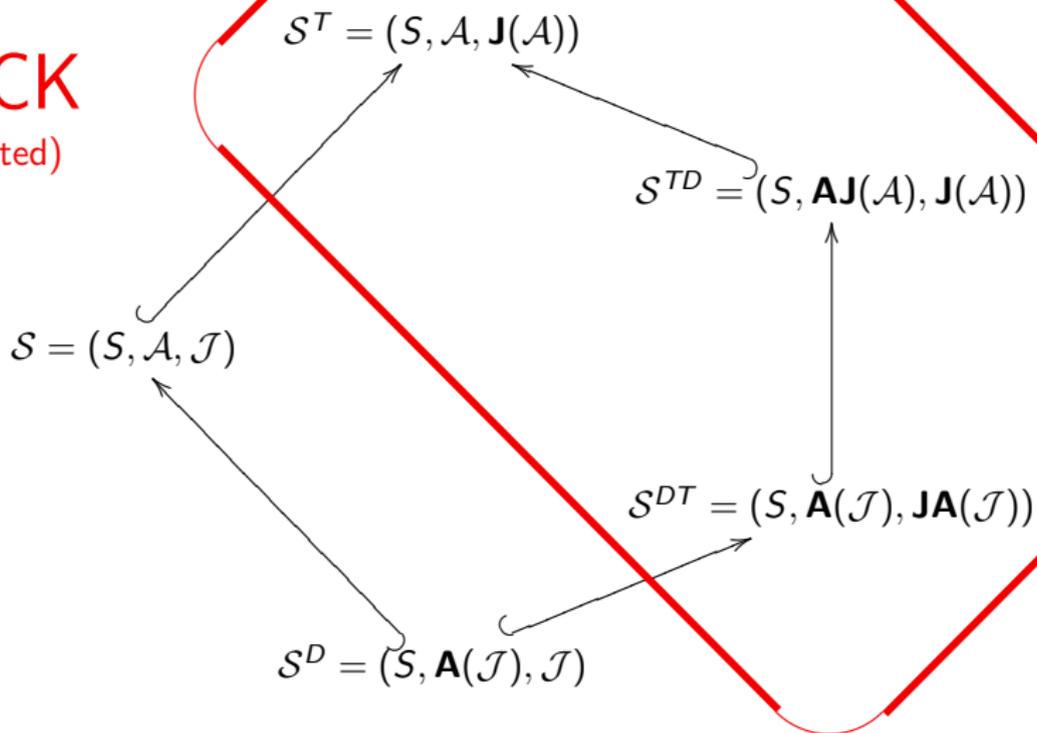
(“thickification”) $\mathcal{S}^T = (S, \mathcal{A}, \mathbf{J}(\mathcal{A}))$

the least *thick* (*generated*) basic topology
which contains \mathcal{S}

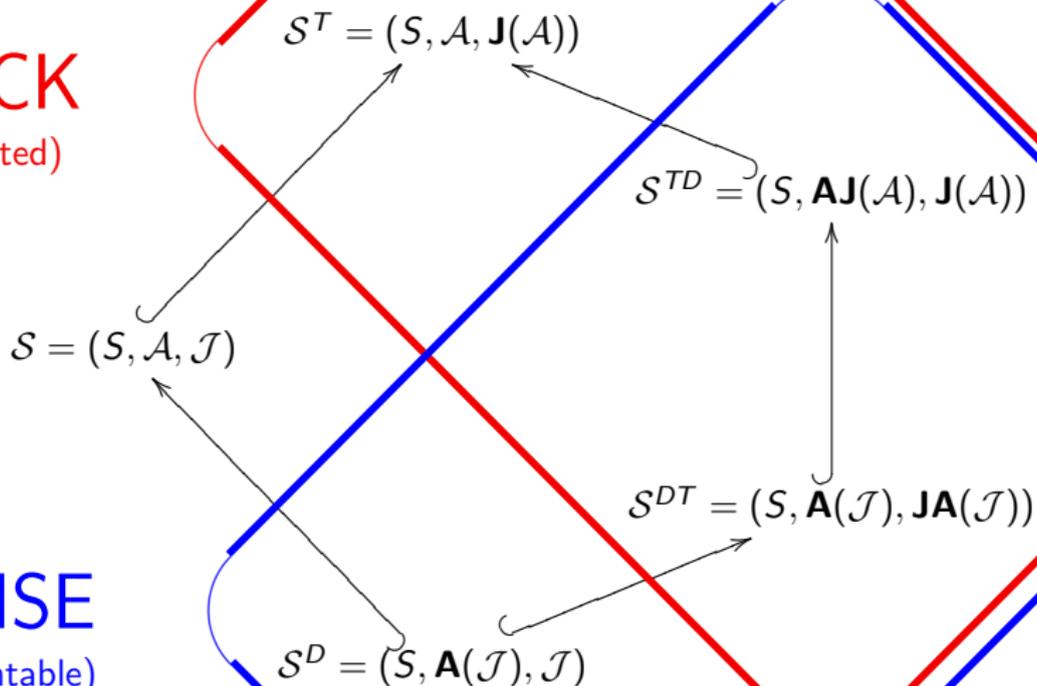
$$\mathcal{S}^D \hookrightarrow \mathcal{S} \hookrightarrow \mathcal{S}^T$$



THICK
(generated)



THICK
(generated)



DENSE
(representable)

Part V

Continuity for dense and thick topologies

Continuous relations

(morphisms of basic topologies)

A morphism $(S_1, \mathcal{A}_1, \mathcal{J}_1) \xrightarrow{r} (S_2, \mathcal{A}_2, \mathcal{J}_2)$ is given by a binary relation r between S_1 and S_2 such that:

(covers/saturations)

$$\frac{b \triangleleft_2 V}{r^{-}[\{b\}] \triangleleft_1 r^{-}[V]}$$

\Updownarrow

$$r^{-}[\mathcal{A}_2(U)] \subseteq \mathcal{A}_1(r^{-}[V])$$

Where: $a \in r^{-}[V] \stackrel{\text{def}}{\Leftrightarrow} (\exists v \in V)(a r v)$
(inverse image of r)

(positivities/reductions)

$$r[\mathcal{J}_1(U)] \subseteq \mathcal{J}_2(r[U])$$

Where: $b \in r[U] \stackrel{\text{def}}{\Leftrightarrow} (\exists u \in U)(u r b)$
(direct image of r)

Morphisms of dense and thick topologies

$\underbrace{(S_1, \mathbf{A}(\mathcal{J}_1), \mathcal{J}_1)}_{(dense)} \xrightarrow{r} (S_2, \mathcal{A}_2, \mathcal{J}_2)$	$(S_1, \mathcal{A}_1, \mathcal{J}_1) \xrightarrow{r} \underbrace{(S_2, \mathcal{A}_2, \mathbf{J}(\mathcal{A}_2))}_{thick}$
is continuous if and only if	
$r\mathcal{J}_1 \subseteq \mathcal{J}_2 r$	$r^- \mathcal{A}_2 \subseteq \mathcal{A}_1 r^-$

In particular a **morphism** between two
dense thick
topologies **is determined by** the condition about
reductions saturations

Corollary

The **full** subcategory
DBTop **TBTop**
whose objects are all the
DENSE **THICK**
basic topologies
is a
coreflective **reflective**
subcategory of **BTop**
(basic topologies and continuous relations)

$$\text{DBTop} \begin{array}{c} \hookrightarrow \\ \xleftarrow{(\)^D} \end{array} \text{BTop} \begin{array}{c} \xrightarrow{(\)^T} \\ \hookleftarrow \end{array} \text{TBTop}$$

 S^D S S^T

Adjunction between **dense** and **thick**

The following two adjunctions...

$$\text{DBTop} \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{(\)^D} \end{array} \text{BTop} \begin{array}{c} \xrightarrow{(\)^T} \\ \xleftarrow{\quad} \end{array} \text{TBTop}$$

...yields also the following (by composition)...

$$\text{DBTop} \begin{array}{c} \xrightarrow{(\)^T |_{\text{DBTop}}} \\ \xleftarrow{(\)^D |_{\text{TBTop}}} \end{array} \text{TBTop}$$

References

-  F. Ciraulo and G. Sambin, *A Galois connection between closure and interior operators through the notion of basic topology*, in preparation.
-  T. Coquand - S. Sadocco - G. Sambin - J. Smith, *Formal topologies on the set of first-order formulae*, J. Symb. Log. **65**, No. 3, 1183-1192 (2000).
-  T. Coquand - G. Sambin - J. Smith - S. Valentini, *Inductively generated formal topologies*, Ann. Pure Appl. Logic **124**, No. 1-3, 71-106 (2003).
-  P. Martin-Löf, G. Sambin *Generating positivity by coinduction* in G. Sambin, *The Basic Picture: Structures for Constructive Topology*, Oxford University Press, to appear.
-  G. Sambin, *The Basic Picture: Structures for Constructive Topology*, Oxford University Press, to appear.

ありがとう

[arigatō]

(Thank you!)