Coding for Super Dense Networks 1

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PROBLEMS & MOTIVATIONS



- **Problem 1**: Billions connected devices in 2020², forming *super dense networks*.
- **Problem 2**: High complexity of fixed transmission scheduling in multiway relay systems.
- **Problem 3**: Low throughput of random access technology.
- Prospective applications: disaster areas networks, wireless sensor networks, high density areas networks, satellite communications.



Objective

Design uncordinated multiway relay networks serving *massive* number of users and having high throughput attained by employing multiuser detection.

 $^{^2\,\}mathrm{Ericsson},$ More Than 50 Billion Connected Devices, Ericsson White Paper, February, 2011.

System Model





- All users want to exchange information among themselves with the relay's help.
- Relay: Amplify-and-Forward, Half-Duplex.
- Synchronous (slot) transmission.
- Transmission phases: Multiple Access Channel (MAC) and Broadcast Channel (BC).
- *Pair-of-Time-Slot* (PTS): a MAC phase followed by a BC phase.
- $\bullet~1$ contention period is composed of N PTSs.
- Offered traffic (logical) load delivered to **one user** :

$$G = \frac{M}{N}.$$
 (1)

• The users randomly select the code type $c_h \sim (h, 1)$ according to the given probability mass function (pmf)

$$\Lambda = \{\Lambda_2, \Lambda_3, ..., \Lambda_n\}$$

- Code $c_h \sim (h, 1)$ is packet-oriented repetition coded with order h.
- Networks Rate per user: $R_N = \frac{1}{\sum \Lambda_h h} = \frac{1}{\overline{n}}$, where \overline{n} is the expected length of the code.

TRANSMISSION SCHEME - NETWORKS





PTS 5





Received by all users

→ MAC phase ····· BC phase

PTS 4

Decoding & Asymptotic Analysis

- Find s node with degree d=1, then do:
 - Local Decoding in s node and then u node.
 - SIC in the connected *s* nodes.
- If no s node with d=1, find s node with d=2, then do :
 - Local decoding in s node using iterative demapping $(IDM)^3$ algorithm to decode two users messages.
 - Local decoding in *u* nodes.
 - $\bullet\,$ SIC in the connected s nodes.
- Stop if all information are decoded or max iteration.

Asymptotic Analysis

- Degree distribution of u nodes is $\lambda(x) = \sum_{h=2}^{n} \lambda_h x^{h-1}$
- Degree distribution of s nodes is $\rho(x) = \exp\left(-\frac{\overline{n}}{G}(1-x)\right)$
- The erasure probability from u nodes to s nodes, and from s nodes to u nodes, respectively:

$$q = \sum_{h=2}^{n} \lambda_h p^{h-1} := f_u(p) \tag{2}$$

$$p = 1 - \left(1 + q\frac{G}{R_N}\right)e^{-q\frac{G}{R_N}} := f_s(q) \tag{3}$$





 $^{^3}$ Anwar, K. and Matsumoto, T., Three-way relaying systems using iterative spatial demapping, ISTC, Aug. 2012

EXIT CHART & UPPER BOUND





(a) EXIT Chart

(b) Upper Bound

- The evolution of $f_u(p)$ and $f_s(q)$ are plotted into the EXIT chart to characterize the convergence behavior of decoding process.
- Threshold value G^* is defined as the maximum value of G such that there is open tunnel between two curves.
- The upper bound of the system derived from EXIT chart area theorem is given by

$$R_N + \left(1 + \frac{2R_N}{G}\right)e^{-G/R_N} - \frac{2R_N}{G} < 0, \tag{4}$$

SIMULATION RESULTS - THROUGHPUT & PLR





- Simulation parameters :
 - N = 200 PTSs

 - *M* is variable depends on *G* $\Lambda_a(x) = 0.5x^2 + 0.28x^3 + 0.22x^8$ $\Lambda_b(x) = 0.25x^2 + 0.6x^3 + 0.15x^8$

• The proposed systems achieve throughput of more than 1 packet/slot and very low PLR.



- We have proposed uncoordinated multiway relaying systems that can serve massive number of users.
- ② Employment of IDM as the multiuser detector increases the throughput of the systems significantly. It make possible to utilize number of slots less than number of users.
- (2) The proposed systems can be optimized by choosing a proper degree distribution of u nodes such that G^* close to the bound.