

An introduction to the independent set reconfiguration problem (ISRECONF)

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The Reconfiguration Problem

- INSTANCE:
 - ① Collection of configurations.
 - ② Allowed transformation rule(s).
- QUESTION: For any two configurations A, B from the given collection, can A be transformed to B using the given rule(s)?

A classic example is the so-called 15-*puzzle*.

- INSTANCE:
 - ④ **Configuration:** 15 tiles numbered from 1 to 15 are arranged on an 4×4 grid, leaving one empty square.
 - ② **Rule:** A tile can move to the empty square if it is above, or below, or on the left, or on the right of that empty square.
- QUESTION: Starting from any *source* configuration, can we reach the *target* configuration, where all numbers are in order (right-hand picture in the figure below)?

8	15	13	3
10		14	7
5	1	2	4
9	12	11	6

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	

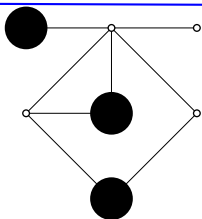


Figure: An independent set of a graph. Independent vertices are marked with black tokens.

- 1 TS: Slide tokens along edges.
- 2 TJ: A token “jumps” from one vertex to another.
- 3 TAR: Add or Remove tokens.

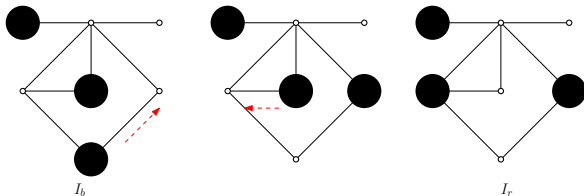


Figure: A YES-instance under TS rule.

- o INSTANCE:
 - 1 A graph $G = (V, E)$.
 - 2 Two independent sets I_b, I_r .
 - 3 “Reconfiguration” rules: TS, TJ, TAR.
- o QUESTION: Can I_b be transformed to I_r using one of the given rules such that all intermediate sets are independent?

- 1 “Independent set” is an important object in computational complexity and graph theory.

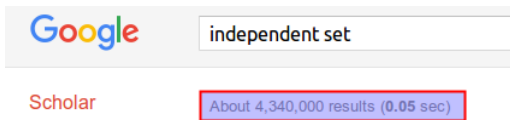


Figure: A search on Google Scholar with keyword “independent set”.

- 2 ISRECONF (under any of the three given rules) is PSPACE-complete for general graphs, perfect graphs, and even planar graphs. Several PSPACE-hardness results are shown using reduction from ISRECONF and its variants.

Marcin Kamiński, Paul Medvedev, and Martin Milanič (2012). “Complexity of independent set reconfigurability problems”. In: *Theoretical Computer Science* 439, pp. 9–15

Jan van den Heuvel (2013). “The complexity of change”. In: *Surveys in Combinatorics 2013*. Ed. by Simon R. Blackburn, Stefanie Gerke, and Mark Wildon. Cambridge University Press, pp. 127–160

Graph	Rule(s)	Complexity	Paper(s)
planar	TS, TJ, TAR	PSPACE-complete	Robert A. Hearn and Erik D. Demaine (2005). "PSPACE-completeness of Sliding-block Puzzles and Other Problems Through the Nondeterministic Constraint Logic Model of Computation". In: <i>Theoretical Computer Science</i> 343.1-2, pp. 72–96
general	TS, TJ, TAR	PSPACE-complete	Takehiro Ito, Erik D. Demaine, et al. (2011). "On the complexity of reconfiguration problems". In: <i>Theoretical Computer Science</i> 412.12-14, pp. 1054–1065
line	TJ, TAR	P	
perfect	TS, TJ, TAR	PSPACE-complete	
even-hole-free	TJ, TAR	P	Marcin Kamiński, Paul Medvedev, and Martin Milanič (2012). "Complexity of independent set reconfigurability problems". In: <i>Theoretical Computer Science</i> 439, pp. 9–15
cograph (P_4 -free)	TS	P	
cograph (P_4 -free)	TJ, TAR	P	Paul Bonsma (2014). "Independent Set Reconfiguration in Cographs". In: <i>WG 2014</i> . Ed. by Dieter Kratsch and Ioan Todinca. Vol. 8747. LNCS. Springer, pp. 105–116
bounded bandwidth	TJ, TAR	PSPACE-complete	Marcin Wrochna (2014). "Reconfiguration in bounded bandwidth and treedepth". In: <i>arXiv preprints arXiv:1405.0847</i>
claw-free	TS, TJ	P	Paul Bonsma, Marcin Kamiński, and Marcin Wrochna (2014). "Reconfiguring Independent Sets in Claw-Free Graphs". In: <i>SWAT 2014</i> . Ed. by R. Ravi and Ingelil Gørtz. Vol. 8503. LNCS. Springer, pp. 86–97
tree	TS	P	Erik D. Demaine et al. (2015). "Linear-time algorithm for sliding tokens on trees". In: <i>Theoretical Computer Science</i> 600, pp. 132–142
bipartite permutation	TS	P	Eli Fox-Epstein et al. "Sliding Token on Bipartite Permutation Graphs". In: <i>ISAAC 2015</i> , Nagoya, Japan, Dec. 9–11. (To be appeared)
bipartite distance-hereditary	TS	P	

Table: Recent results on studying ISRECONF.

In the process of studying reconfiguration problem and its variants (including ISRECONF), the following tools are particularly useful.

- 1 **NONDETERMINISTIC CONSTRAINT LOGIC (NCL):** This framework is motivated from the so-called *sliding-block puzzle*, which is a generalization of the classic 15-puzzle. NCL is a general framework built for proving PSPACE-hardness which simply requires the construction of a couple of gadgets that can be connected together in a planar graph. More details can be found in [Robert A. Hearn and Erik D. Demaine \(2005\)](#). “PSPACE-completeness of Sliding-block Puzzles and Other Problems Through the Nondeterministic Constraint Logic Model of Computation”. In: *Theoretical Computer Science* 343.1-2, pp. 72–96
- 2 **DYNAMIC PROGRAMMING:** Roughly speaking, DYNAMIC PROGRAMMING is a method for solving a complex problem by breaking it down into a collection of simpler subproblems. A general setting of DYNAMIC PROGRAMMING method for solving reconfiguration problems was recently presented in [Paul Bonsma and Daniel Paulusma \(2015\)](#). “Using Contracted Solution Graphs for Solving Reconfiguration Problems”. In: *arXiv preprints arXiv:1509.06357*

o INSTANCE:

- ① A graph $G = (V, E)$.
- ② Two cliques C_0, C_r .
- ③ Reconfiguration rules:
TS, TJ, TAR.

- o QUESTION: Can C_0 be transformed to C_r using one of the given rules such that all intermediate sets are clique?

- ① TS: Slide a token in clique C to a vertex in $V \setminus C$ along an edge of G .
- ② TJ: A token in clique C jump to a vertex in $V \setminus C$.
- ③ TAR: Add or Remove tokens.

Interestingly, in CLIQUE RECONFIGURATION, the three rules TS, TJ and TAR are equivalent, while in ISRECONF they are not. More details are in [Takehiro Ito, Hirotaka Ono, and Yota Otachi \(2015\)](#). "Reconfiguration of Cliques in a Graph". In: *TAMC 2015, Singapore, May 18-20, 2015, pp. 212–223*

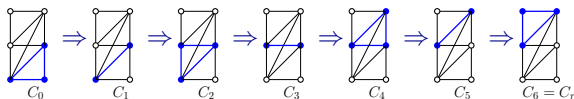


Figure: A sequence $\langle C_0, C_1, \dots, C_6 \rangle$ of cliques in the same graph, where the vertices in cliques are depicted by large (blue) circles (tokens). (© Ito, Ono, and Otachi 2015)

Some interesting directions for ISRECONF study are as follows:

- 1 There are a long list of open problems: What is the complexity of ISRECONF for ... graph under ... rule?
- 2 Also, it is natural to ask questions about the *length of the reconfiguration sequence* (number of intermediate sets required to transform the source configuration to the target one). Recently, the ISRECONF problem for trees was shown to be in P (under any of the three rules), but the corresponding “shortest reconfiguration sequence” problem is still open?
- 3 Many researchers are interested in studying the connection between a decision problem (such as the famous INDEPENDENT SET problem) and its corresponding reconfiguration version (such as ISRECONF). Interestingly, for a general graph, INDEPENDENT SET is NP-complete, and ISRECONF is PSPACE-complete.