

極微細LSIのタイミング設計

Timing Issues in Nanotechnology LSI

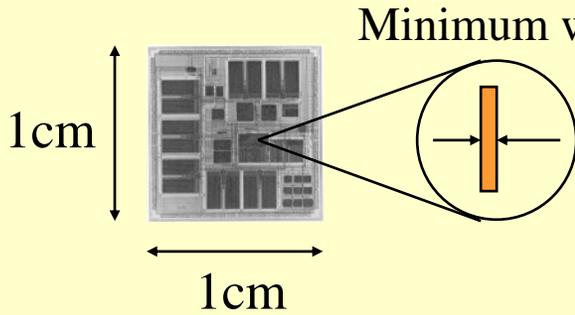
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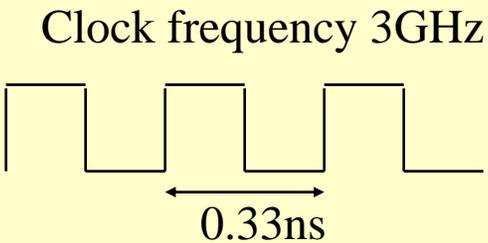
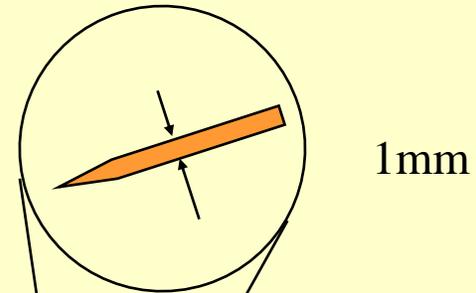
Mineo Kaneko

Information Science, JAIST

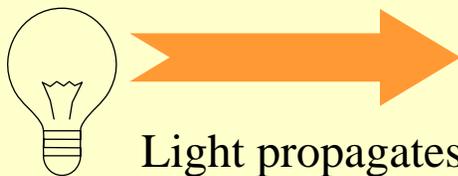
VLSI in the Year 2007



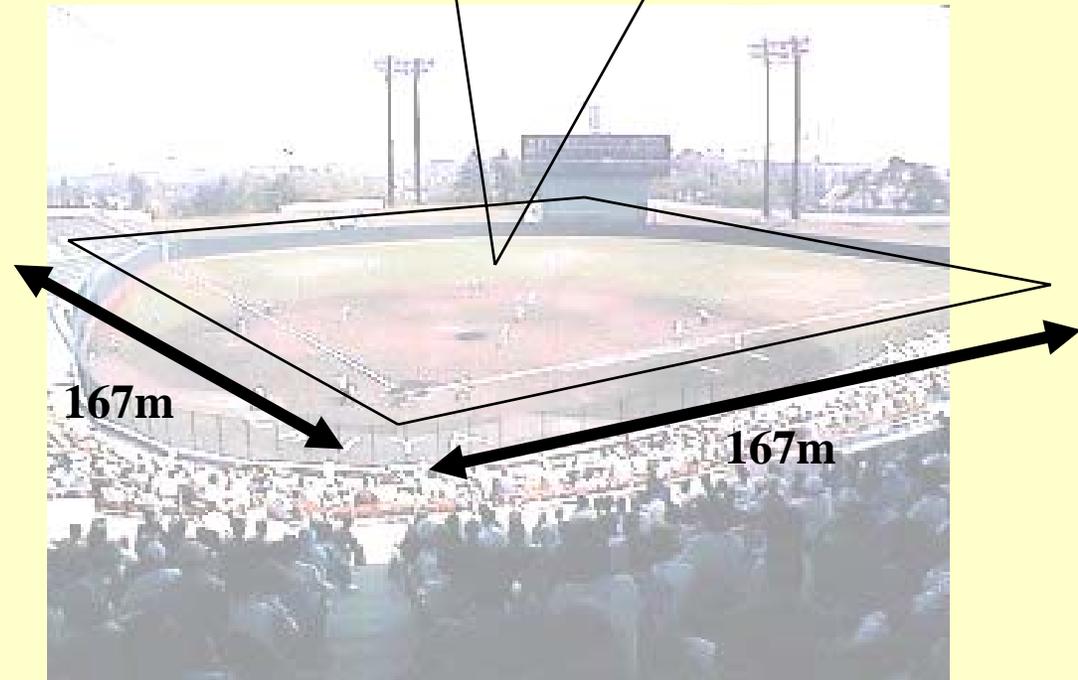
$$\frac{60 \text{ [nm]}}{1 \text{ [cm]}} = \frac{1 \text{ [mm]}}{167 \text{ [m]}}$$



$$(3 \times 10^8 \text{ [m]}) \times (0.33 \times 10^{-9} \text{ [s]}) = 0.1 \text{ [m]}$$

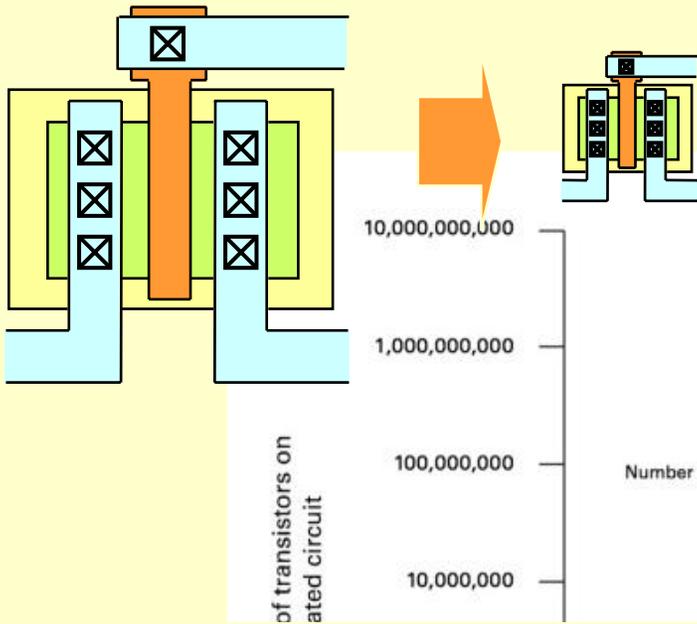


Light propagates 10cm in 0.33ns.



Arranging 1mm ϕ wire in Baseball ground

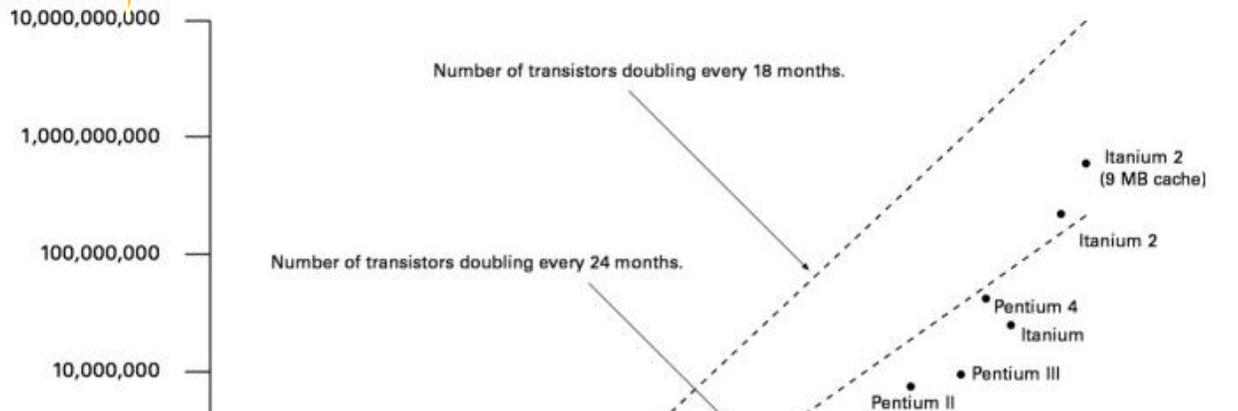
History of IC = History of Shrinking



Moore's Law (Gordon E. Moore)

Complexity grows double in every 18-24 months

"Cramming more components onto integrated circuits", Electronics Magazine 19 April 1965



Shrink → High Space-Density → More Transistors in a chip

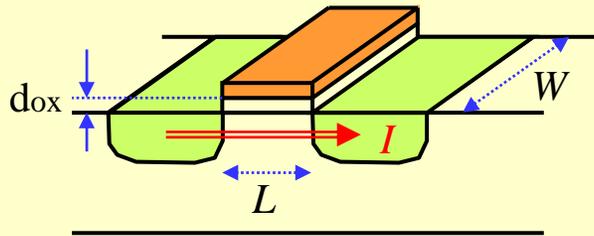
Shrink → Improved Tr. performance → High Speed IC

Further shrinking

- Large propagation delay,
- Inaccuracy in delay estimation,
- Static and dynamic delay fluctuation

Electrical Aspect of VLSI

MOS Transistor



Current I , Voltage V , Capacitance C_g

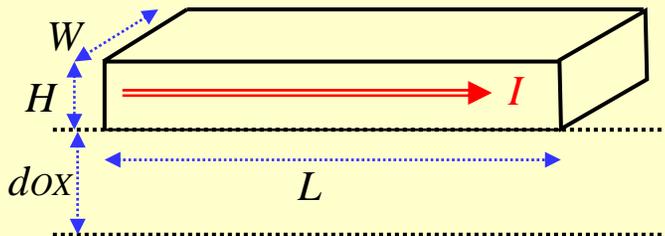
$$I \propto d_{ox}^{-1} \cdot W \cdot L^{-1} \cdot V^2$$

$$C_g \propto d_{ox}^{-1} \cdot W \cdot L$$

Switching delay; *delay*

$$delay \propto \frac{C_g \cdot V}{I} = \frac{d_{ox}^{-1} \cdot W \cdot L \cdot V}{d_{ox}^{-1} \cdot W \cdot L^{-1} \cdot V^2} = L^2 \cdot V^{-1}$$

Wire



Resistance R_w , Capacitance C_w

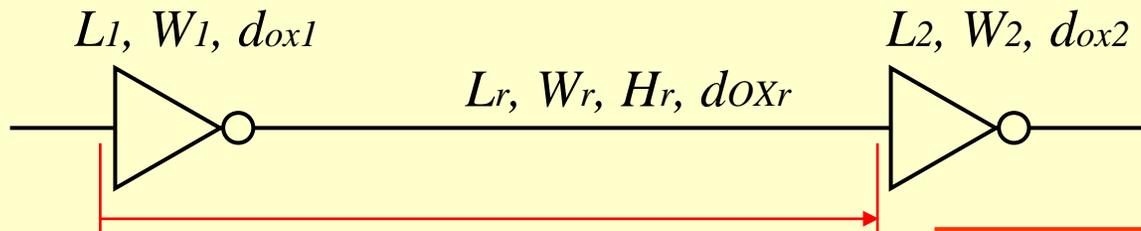
$$R_w \propto L \cdot H^{-1} \cdot W^{-1}$$

$$C_w \propto W \cdot L \cdot d_{ox}^{-1}$$

Propagation delay; *delay*

$$delay \propto \frac{C_w \cdot V}{I} = C_w \cdot R_w = \frac{W \cdot L^2}{H \cdot W \cdot d_{ox}}$$

Electrical Aspect of VLSI



Delay = Switching delay + Propagation delay

Propagation delay is not improved by shrinking

$$\text{Elmore delay (First Moment Model)} \quad \text{delay} = \lambda \frac{(d_{oxr}^{-1} W_r L_r + d_{ox2}^{-1} W_2 L_2) L_1}{d_{ox1}^{-1} W_1 \cdot V} + \kappa \frac{L_r}{W_r H_r} \left(\frac{d_{oxr}^{-1} W_r L_r}{2} + d_{ox2}^{-1} W_2 L_2 \right)$$

Difficulty in delay estimation: Need higher-order model

Various parasitic effects

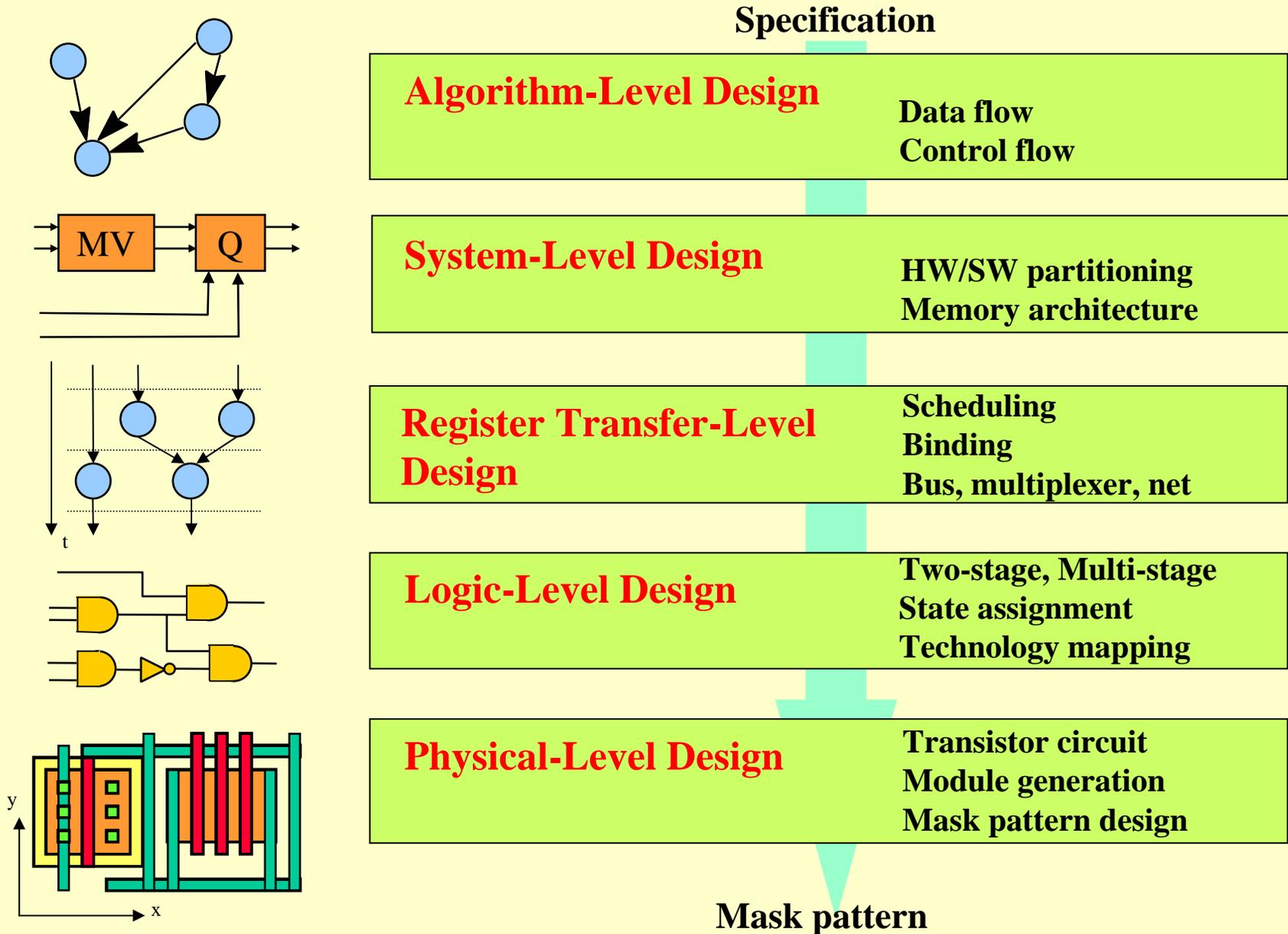
Static/Dynamic delay fluctuation

Static/Dynamic delay fluctuation: Fluctuations of chemical density and physical size in the fabrication process

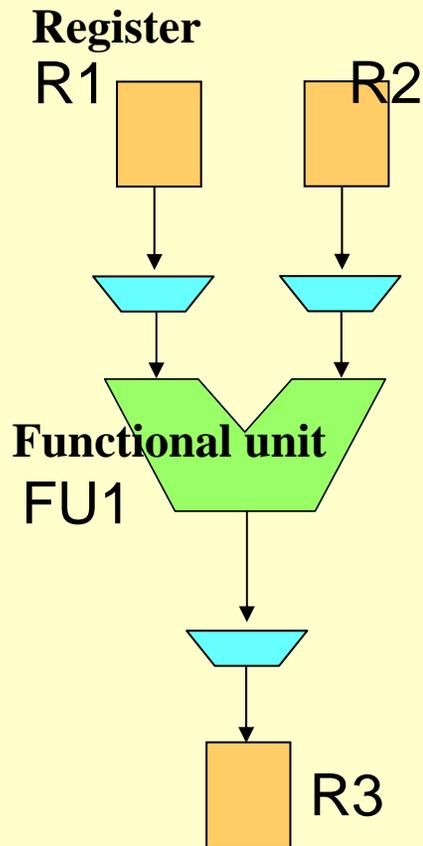
Noise on supply voltage

Cross-talk noise

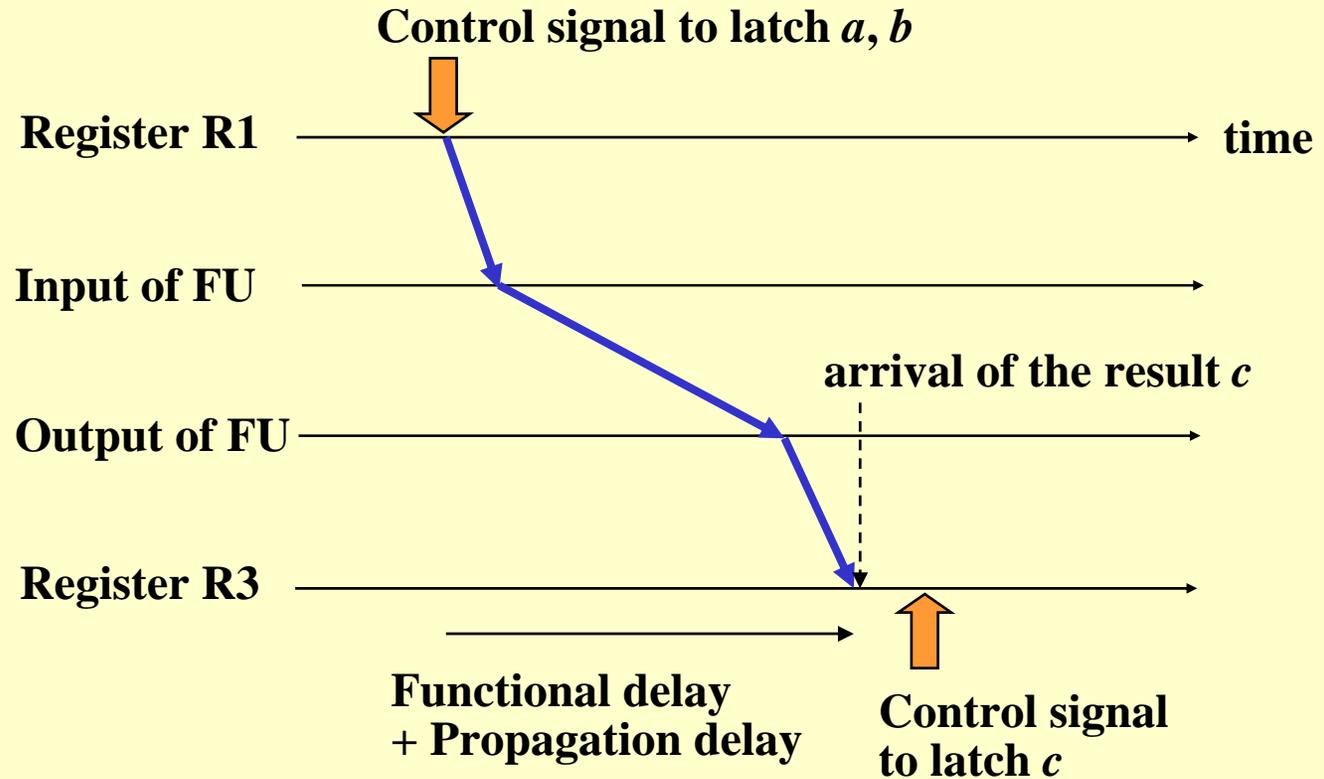
Top-down Hierarchical Design of VLSI



Timing issues in Data-path



Execution of $c = a + b$

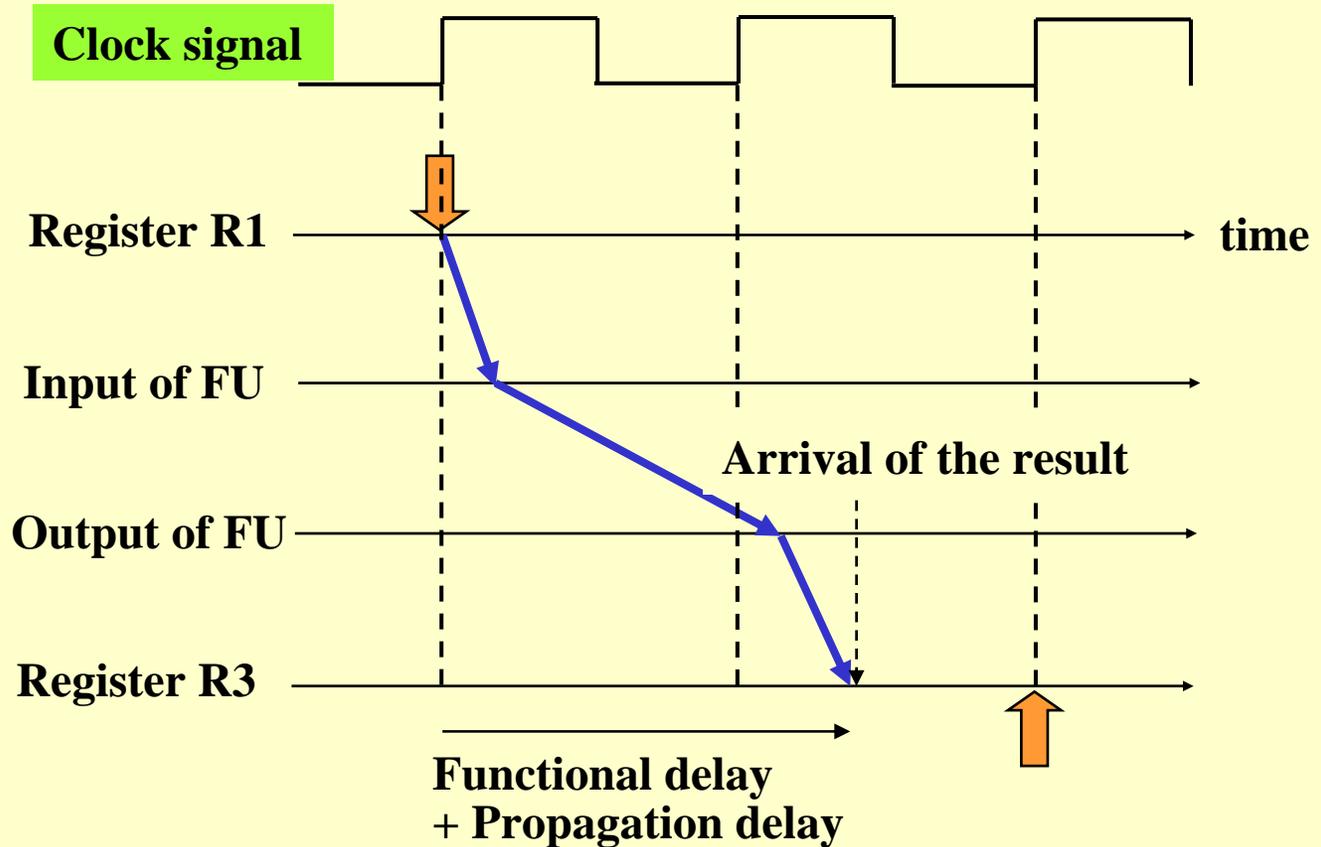
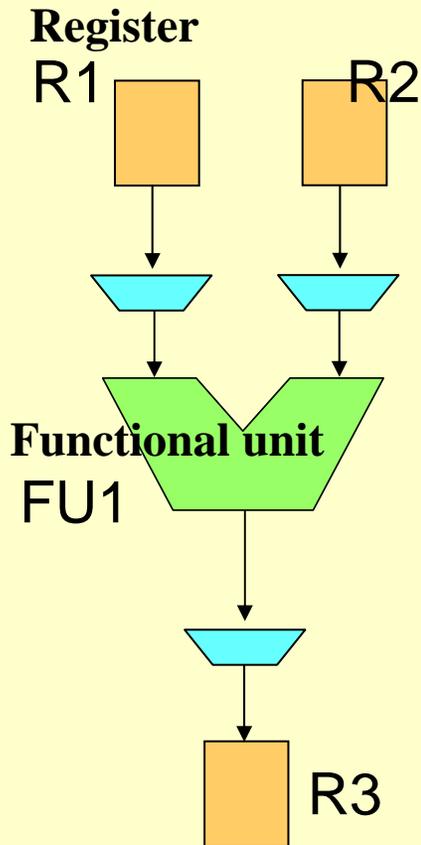


Timing of control signals (\Downarrow , \Uparrow) determines data-path behavior

- Many operations share the same FU, many data share the same register.
- Various different delay values
- Delay may vary from its nominal value statically and dynamically.

Timing issues in Data-path

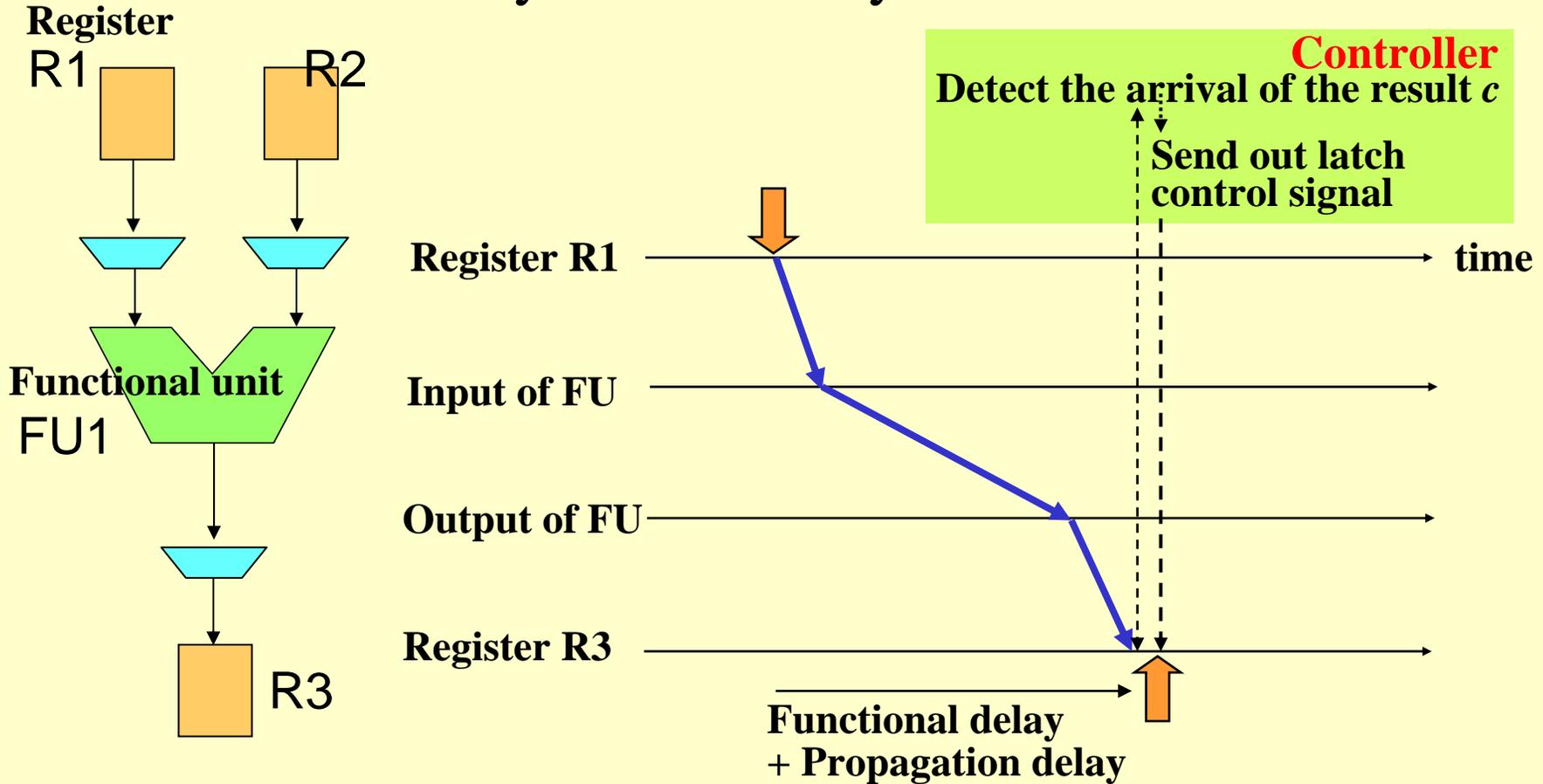
— Synchronous System —



- Design needs "Delay estimation" + "Timing margin"
- Easy to implement as a circuit
- Worst-case estimation + Sufficient margin = Low performance

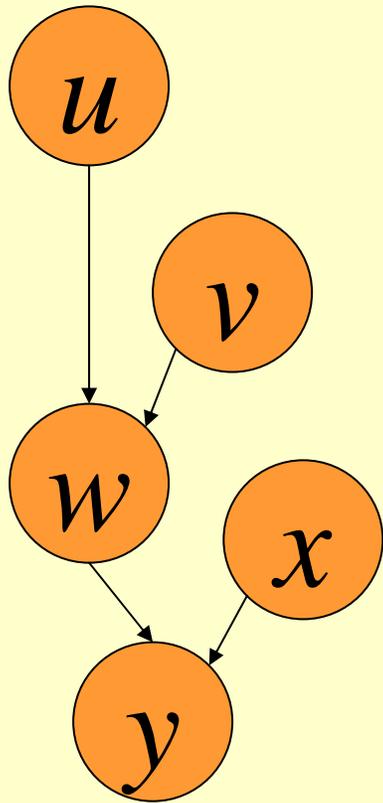
Timing issues in Data-path

— Asynchronous System —

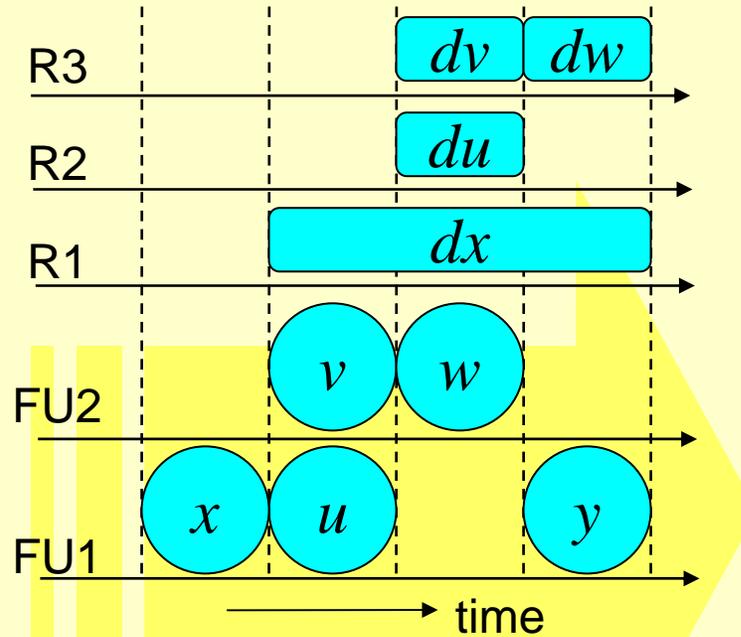


- No delay-estimation, no timing-margin
- Tolerance to a large range of delay fluctuation
- Large area (circuit) overhead in detecting-circuitry

What is High-Level Synthesis?



Behavioral description
of an application
algorithm

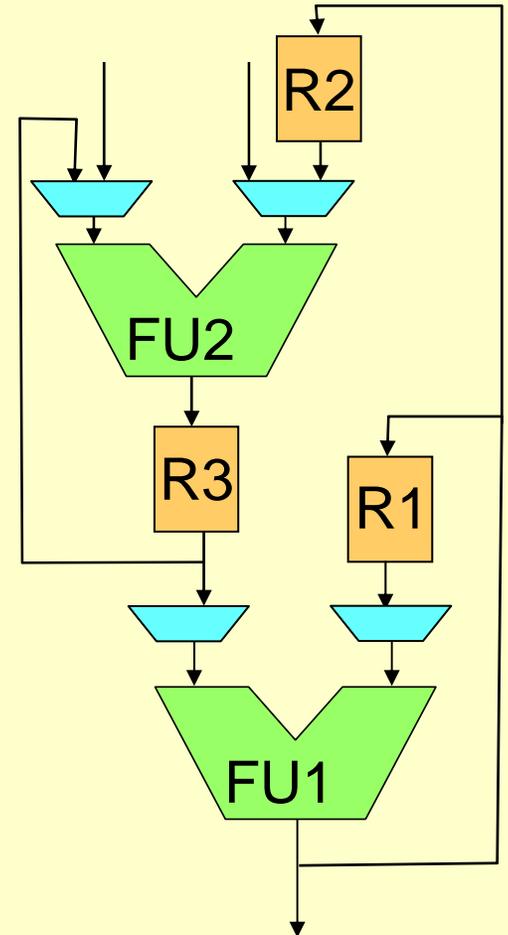


Scheduling;

determines the start time of
each operation

Resource Binding;

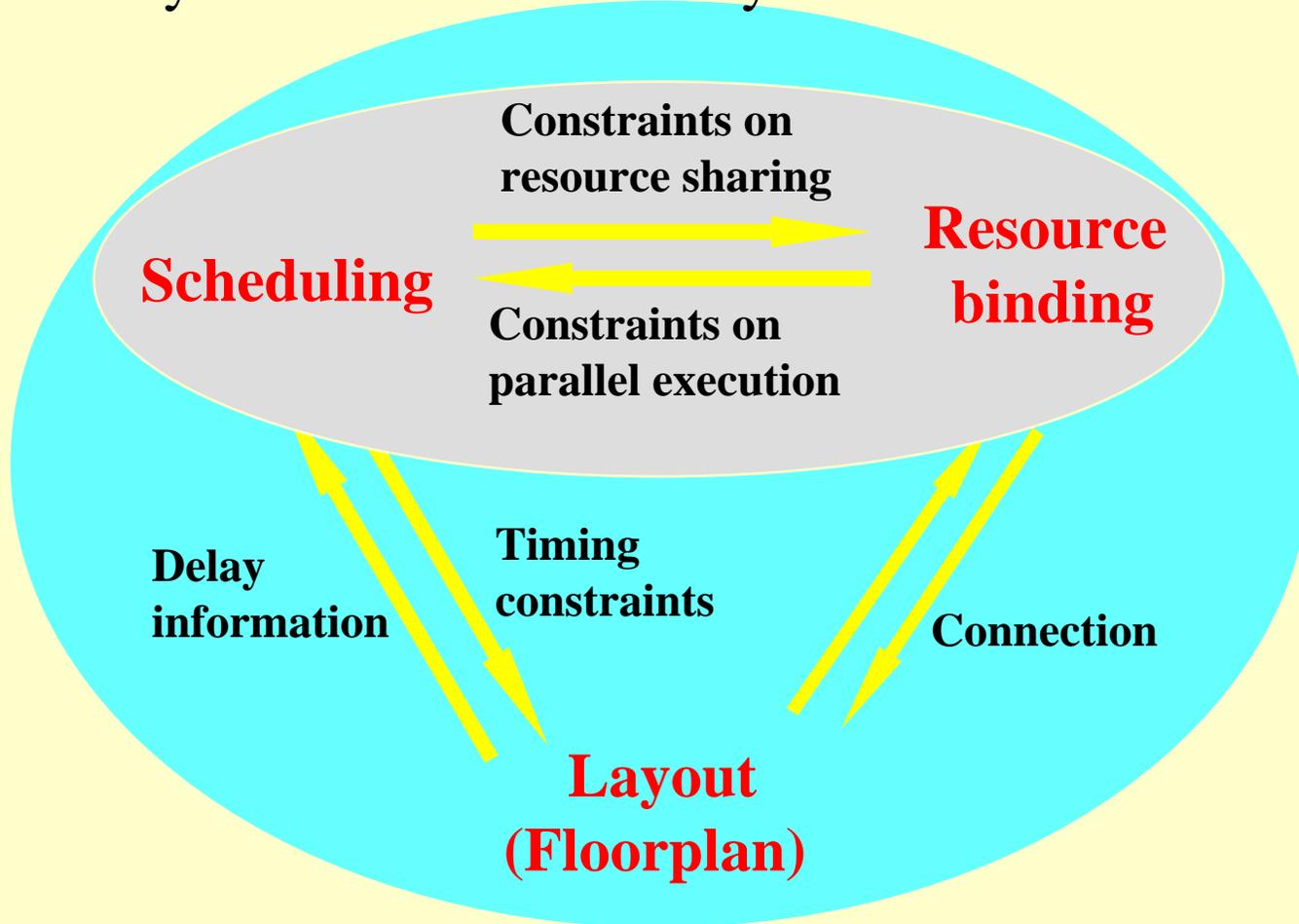
assigns each operation to one of available
functional units, and assigns each data to
one of available registers



**Data-path part
+ Control part**

A. New Approach to High-Level Synthesis

High-level synthesis under FU-delay dominant situation

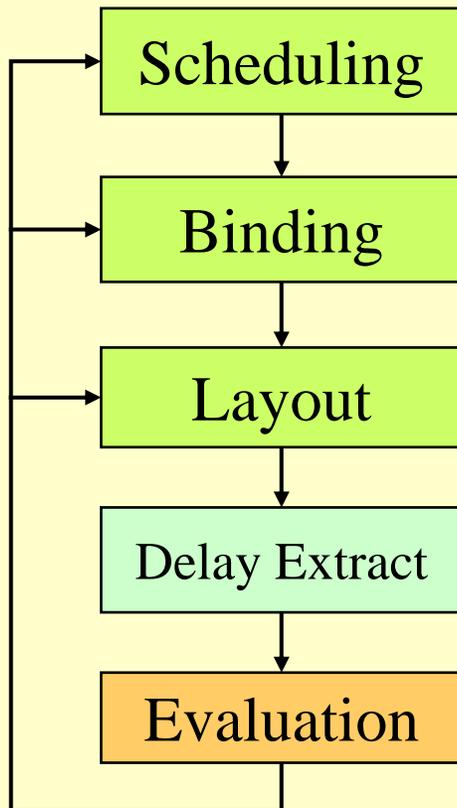


Wire-delay aware high-level synthesis:

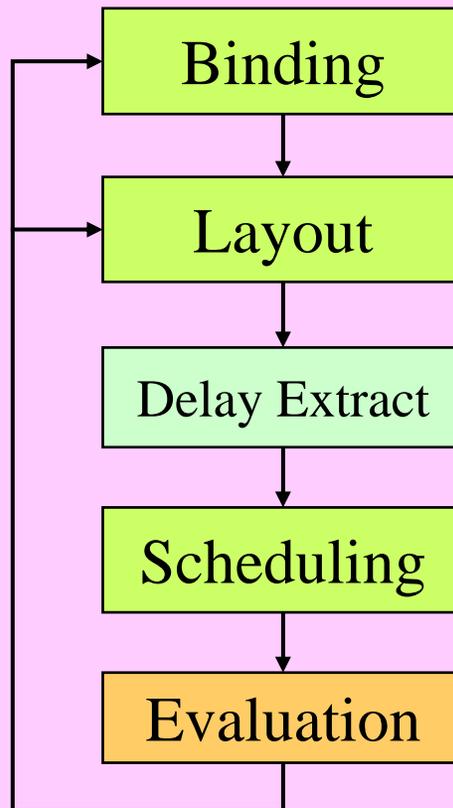
A. New Approach to High-Level Synthesis

Wire-delay aware high-level synthesis:

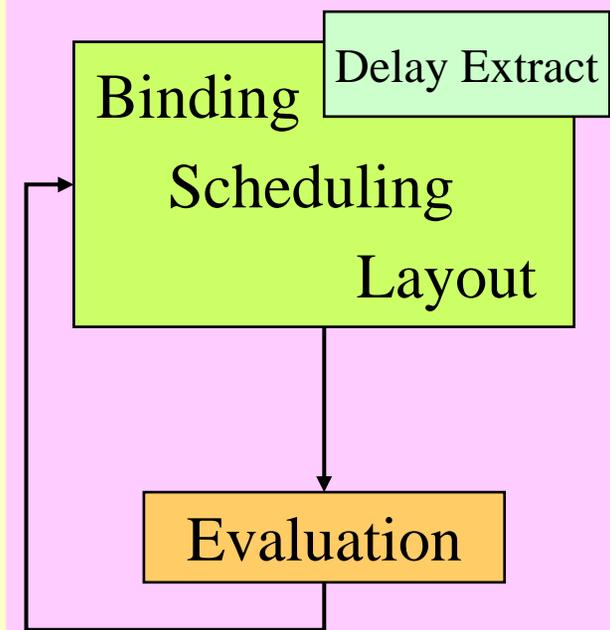
Scheduling-Centric



Binding-Centric

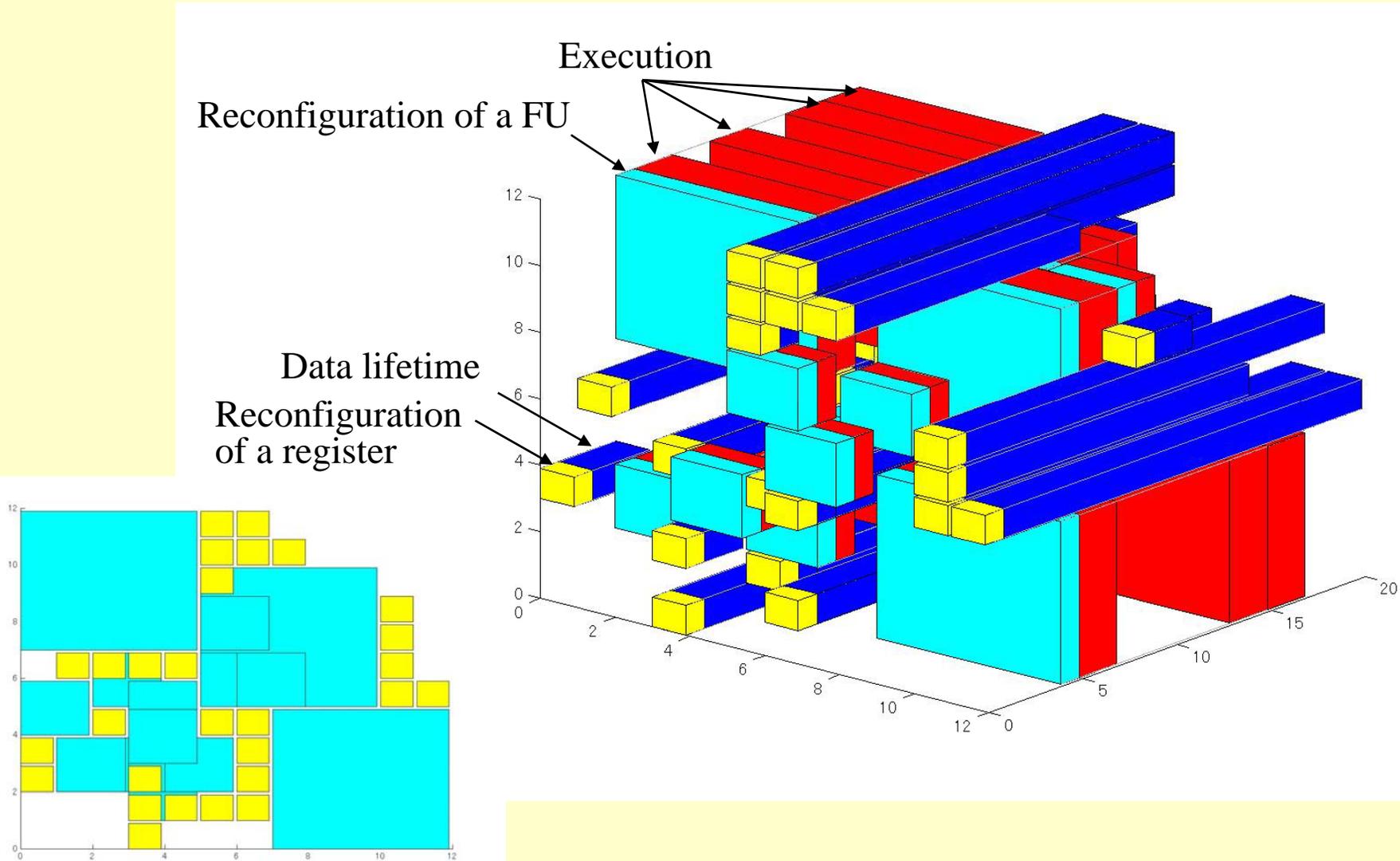


3D-Approach



A. 3D-Approach High-Level Synthesis

Application to Dynamically Reconfigurable LSI



A. 3D-Approach High-Level Synthesis

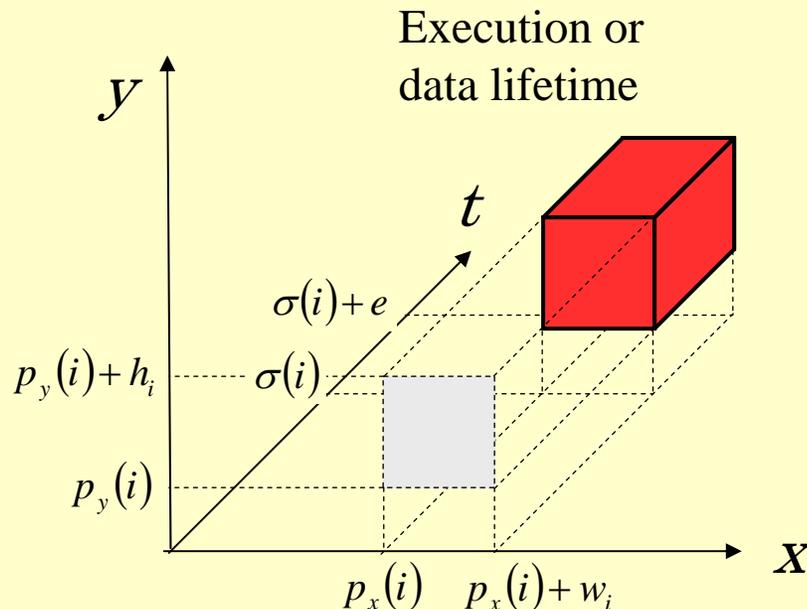
Computation algorithm to be implemented (Dependence Graph)

$DG = (O, D, A)$ O, D : Set of operations, and set of data

A : Dependency

$e: O \rightarrow N$; Operation delay

Sizes of functional units and registers



For each operation/data

$(p_x(i), p_y(i), \sigma(i))$

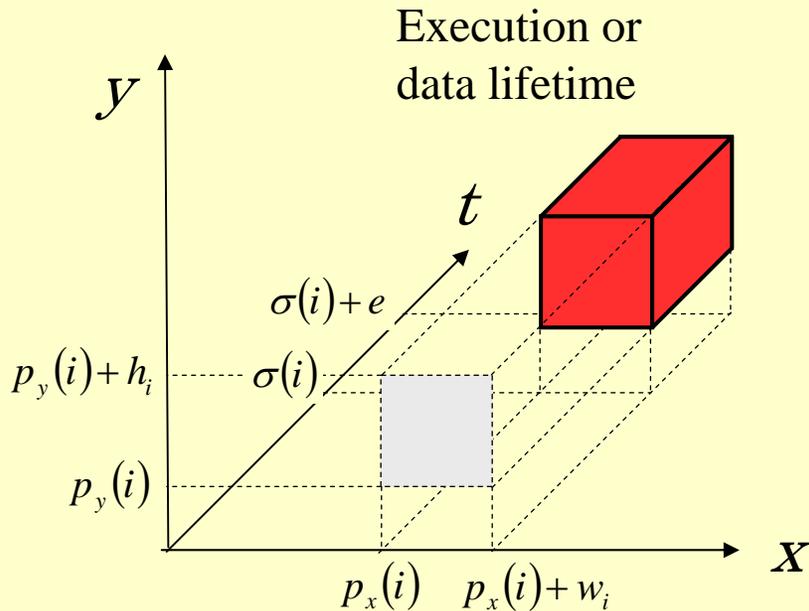
$p_x(i) \in N, p_y(i) \in N, \sigma(i) \in N$

Naive solution space $N^{3(|O|+|D|)}$

Need to check conflicts

Need to check timing constraints

A. 3D-Approach High-Level Synthesis



Constrained Sequence-Quintuple

5-tuple $(\Gamma_1, \Gamma_2, \Gamma_3, \Gamma_4, \Gamma_5)$

- Each of $\Gamma_1, \Gamma_2, \Gamma_3, \Gamma_4$ is a permutation of elements in O, D
 - Γ_5 is a permutation of elements in O (a topological order w.r.t. DG)
- $(\Gamma_1, \Gamma_2, \Gamma_3, \Gamma_4, \Gamma_5)$ represents relative spatial relation in x - y - t space.
 - $O((|O|+|D|)^2)$ computation-time algorithm to compute $\{(p_x(i), p_y(i), \sigma(i)) \mid i \in O \cup D\}$ which has the minimum layout area and the minimum makespan among all solutions satisfying the spatial relation specified by the code.
 - The size of the solution space $((|O|+|D|)!)^5$

A. New Approach to High-Level Synthesis

Basic Theory

- Condition for feasible binding
- Efficient solution space for 3D-Approach to High-Level Synthesis

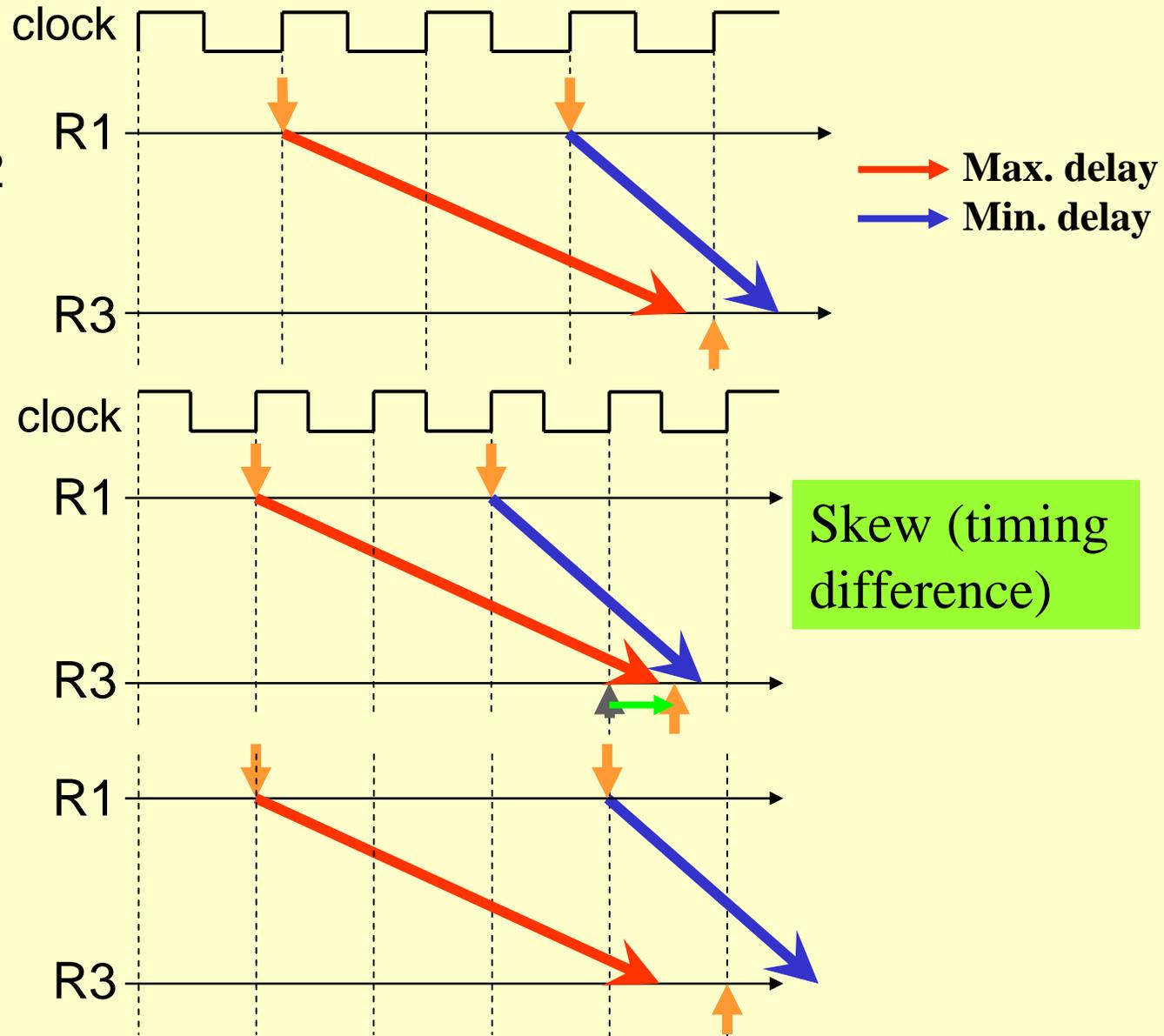
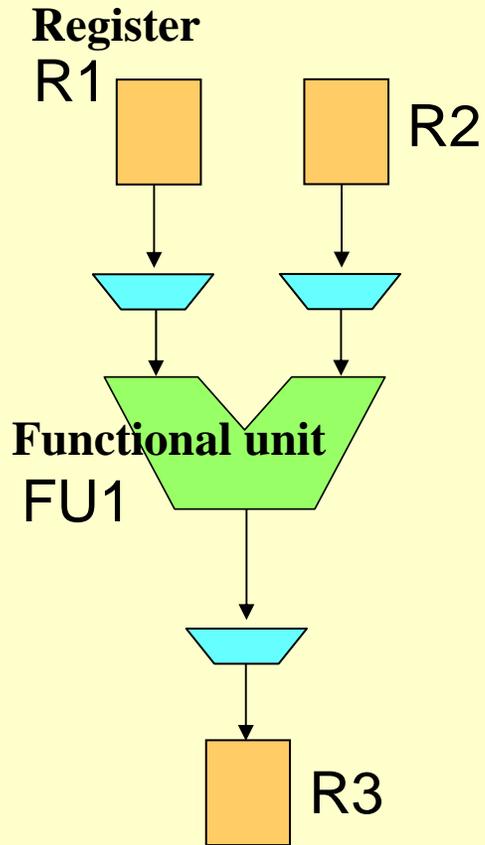
Elementary Technology

- Binding constrained scheduling
- Data-path layout, performance estimation

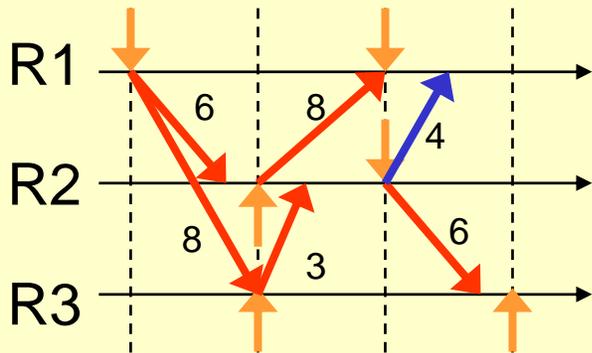
Synthesis System

- Synthesis system considering wire delay
- Synthesis system for reconfigurable systems
- Synthesis system considering control skew
- Synthesis system for asynchronous systems

B. Design Considering Skew



B. Schedule and Skew Optimization

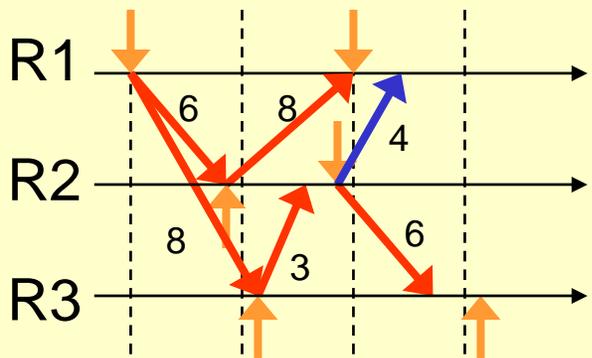


Optimum schedule under zero-skew

Minimum clock period = 8

Schedule length = 3

Total computation time = $8 \times 3 + 0 = 24$



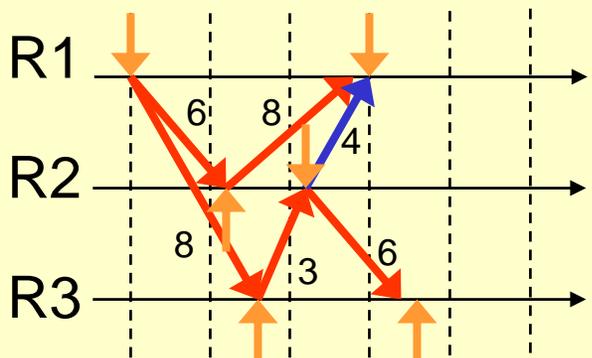
Applying skew optimization

→skew values $(\tau_{r_1}, \tau_{r_2}, \tau_{r_3}) = (0, -1, 1)$

Minimum clock period = 7

Schedule length = 3

Total computation time = $7 \times 3 + 1 = 22$



Simultaneous schedule and skew optimization

→skew value $(\tau_{r_1}, \tau_{r_2}, \tau_{r_3}) = (0, 1, 3)$

Minimum clock period = 5

Schedule length = 3

Total computation time = $5 \times 3 + 3 = 18$

B. Skew-aware High-Level Synthesis

Basic Theory

- Computational Complexity:
Fixed Schedule, Optimize Skew \rightarrow P
Simultaneous Schedule and Skew Optimization
(even if the execution order is fixed) \rightarrow NP-hard

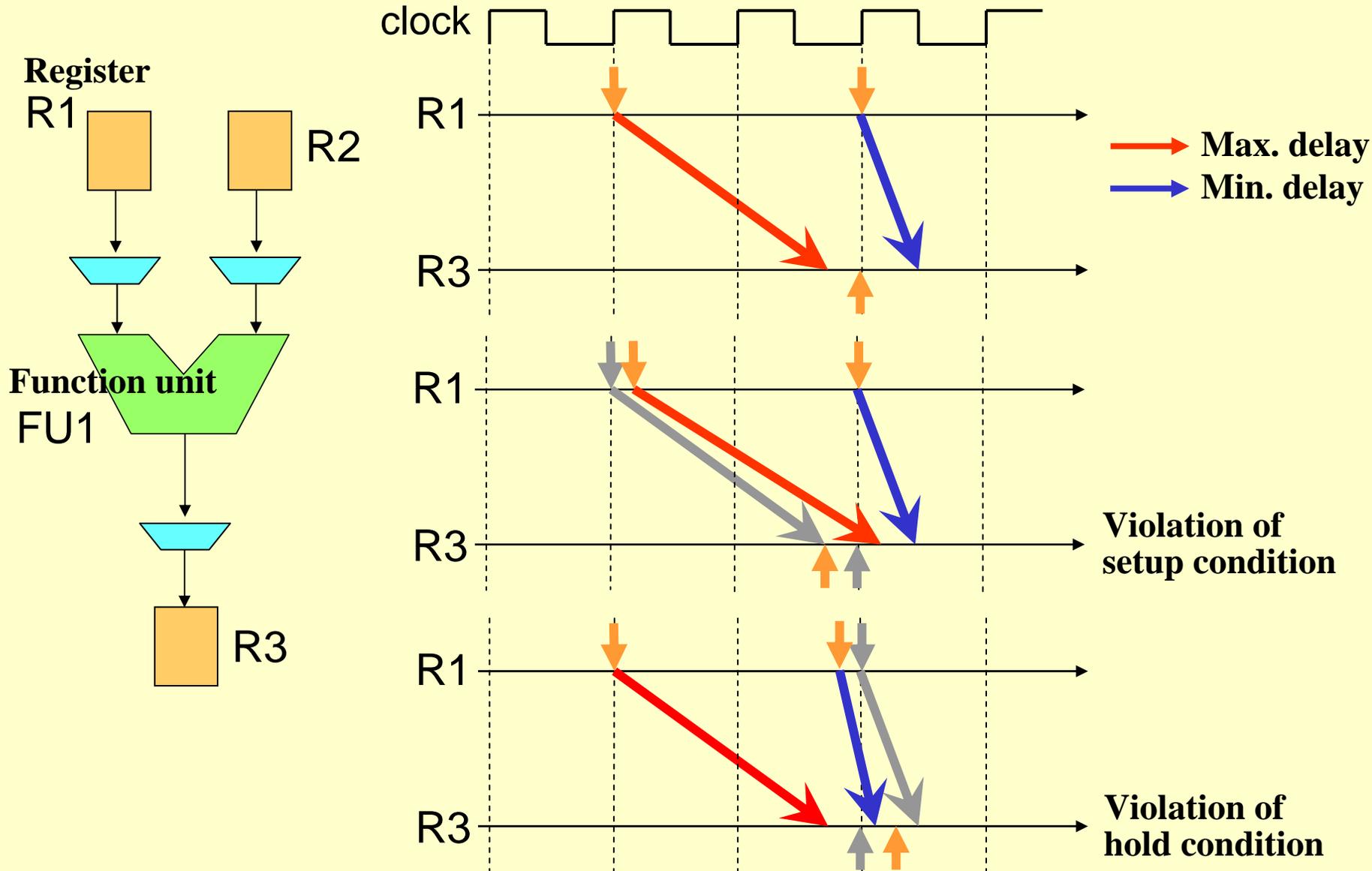
Elementary Technology

- Exact algorithm to compute optimum skew
- Heuristic algorithm for simultaneous schedule and skew optimization

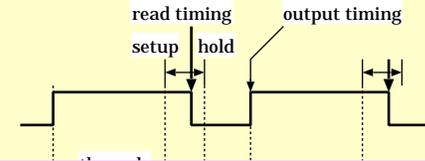
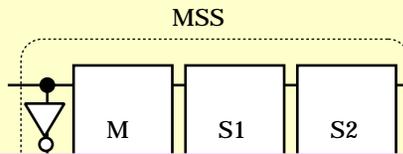
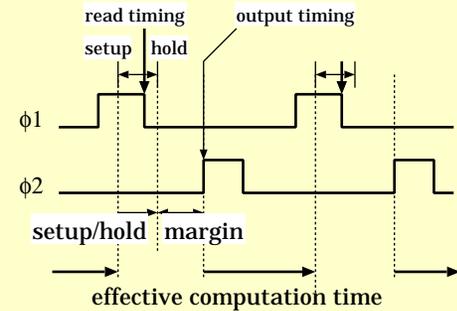
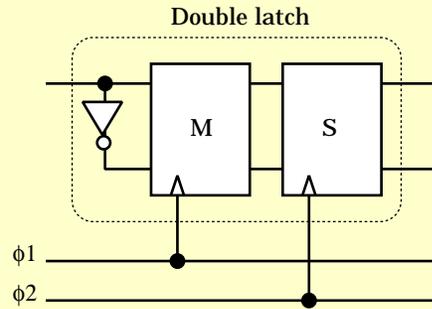
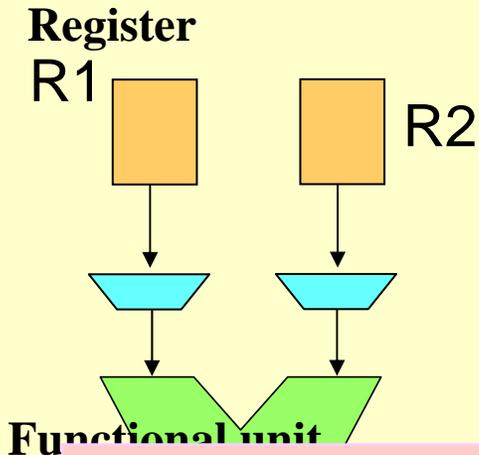
Synthesis System

- Binding-centric approach/3D approach to skew-aware data-path synthesis

C. Delay Fluctuation



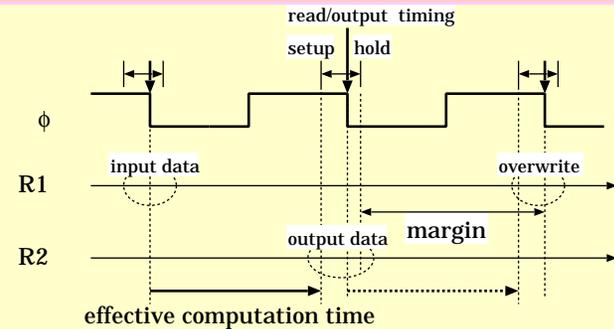
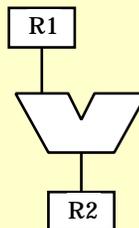
C. Delay Fluctuation



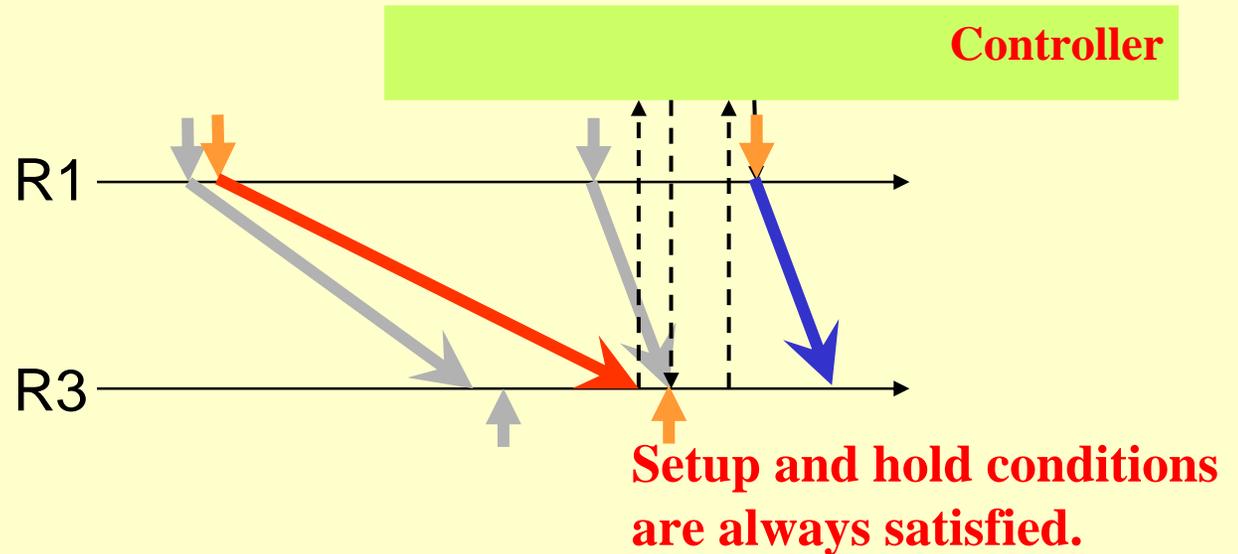
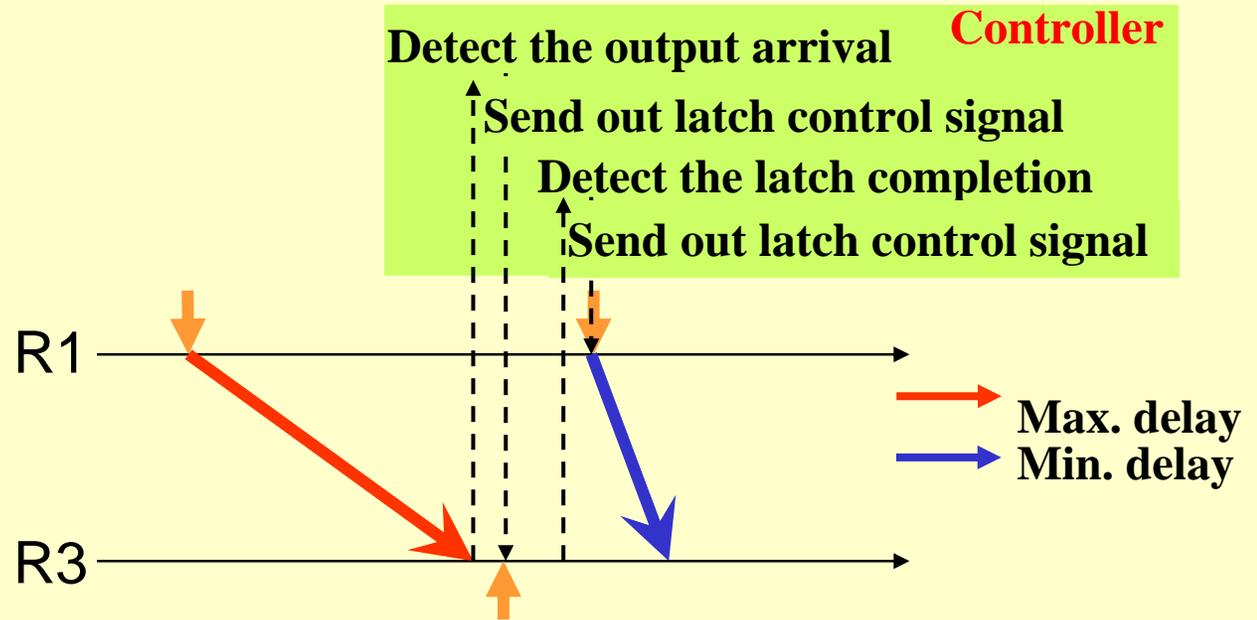
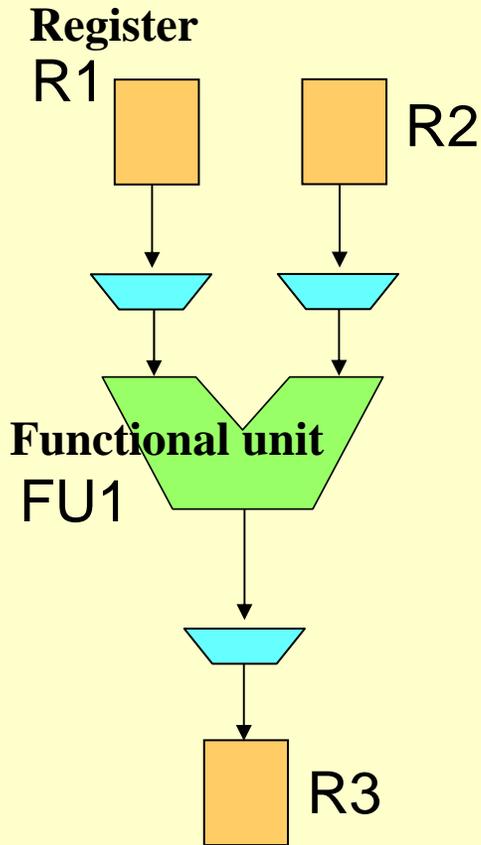
High-Level, Logic-Level, Circuit-Level Synthesis for VLSI which has the Robustness to Delay Fluctuation



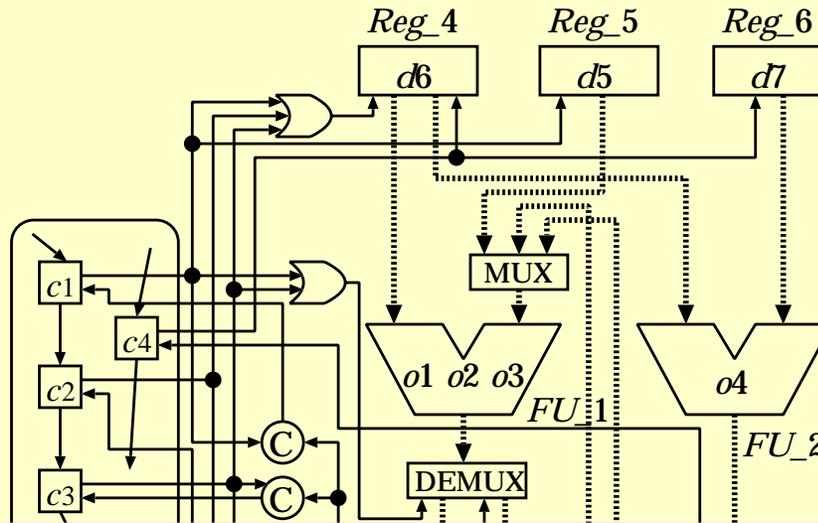
Proposed method



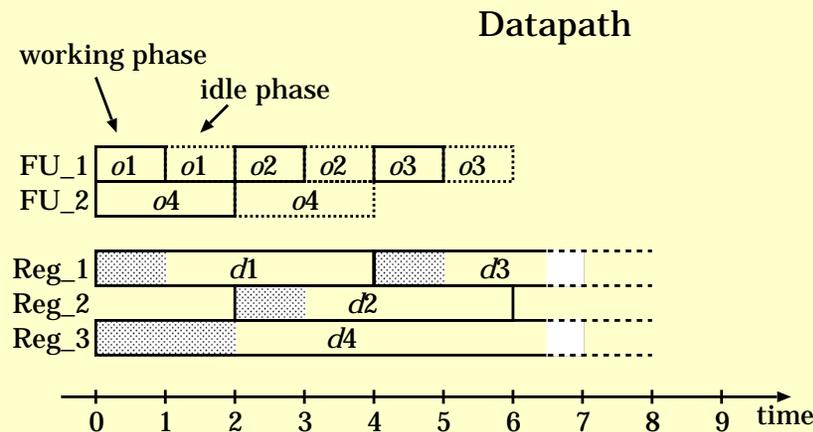
D. Asynchronous System



D. Asynchronous System



**High-Level, Logic-Level, Circuit-Level
Synthesis for Asynchronous System**



VLSI: a core device for reliable e-society

High speed, low power

- Propagation delay and power consumption on signal/clock wires
- Static/dynamic delay fluctuation

➡ **Considering layout in high-level design**

➡ **Robustness, tolerance, insensitiveness to delay fluctuation**

Large scale, system on chip

- Huge size of optimization problems

➡ **Design methodologies to break through the design crisis**

➡ **Efficient algorithms for huge size of problems**

Reliability

- Complex design constraints, a large number of design variables
- Increasing circuit complexity, low voltage, chemical/physical fluctuation

➡ **Reliable design: Reliable EDA tools, 100% automation**

➡ **Reliable chip: VLSI test, fault-tolerance**

High performance/Low power/Reliable System on Chip