



# Simulation Theorems in Multi-valued Modal $\mu$ -Calculus

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# Motivation

"Refinement of Models" in Model Checking

- Model Checking = Modeling + Checking
- Tatsumi and Kameyama tried to get minimal one among models checked successfully.
- They needed a number of model checking.



They wanted to perform a number of model checking all at once.





## Superposition of Models







# From 2={T,F} to general L

- Transition System, Kripke Model, Simulation
- State semantics of Modal  $\mu$  -Calculus, Simulation Theorem
  - De Morgan algebra [Tatsumi-Kameyama 2006]
  - Complete Heyting algebra [This talk]
- Path semantics of Linear Modal  $\mu$  -calculus, Simulation Theorem
  - Complete Heyting algebra + condition [This talk]





# Why complete Heyting algebra ?

- Sets and binary relations form a category.
- L must be a complete Heyting algebra for sets and binary L-valued relations to form a category [Johnstone 2002].





# Complete Heyting algebra

- is  $(L, \leq, \vee, \wedge, \Rightarrow)$  satisfying the following.
- 1.  $(L, \leq)$  is a partially ordered set.
- 2. An arbitrary subset of L has the join (so, also the meet).
- 3. a∧b≦c ⇔ b≦a⇒c

Example:

2,  $2 \times 2$ , ...,  $2^n$ , ...

The open sets of a topological space





# Category of L-valued relations

- Objects are sets
- Arrows from A to B are functions from  $A \times B$  to L







#### Composition: L=2 and L= $2 \times 2$

















# L-valued Transition System







## L-valued Kripke model

#### consists of the following L-relations.

































Transition

































# L-valued State Semantics

#### Modal $\mu$ -Calculus $\psi ::= p |\perp| \top | \psi \lor \psi | \psi \land \psi | \psi \Rightarrow \psi$ $| x | \mu x. \psi | \nu x. \psi | \diamondsuit \psi | \Box \psi$

K,s,V 
$$\vDash \psi$$
 is an element of L

- Natural definition (no details in this talk)
- Intuitionistic version

 $\mathsf{K},\mathsf{s},\mathsf{V}\vDash\psi\quad\neq\quad\mathsf{K},\mathsf{s},\mathsf{V}\vDash(\psi\!\Rightarrow\!\bot)\!\Rightarrow\!\bot$ 





# **Simulation Theorem**

- For any simulation,
- if the abstract model satisfies  $\psi$ ,
- then the concrete model satisfies  $\psi$  .
  - When  $\psi$  has no  $\Box$  in the negative positions and no  $\diamondsuit$  in the positive positions
  - Example:  $\nu X.P \land \Box X$ 
    - "P always globally holds".

#### This theorem holds in L-valued context.





## L-valued Path Semantics

Linear Modal  $\mu$  -Calculus (generalization of LTL)  $\psi ::= p |\perp| T | \psi \lor \psi | \psi \land \psi | \psi \Rightarrow \psi$  $| x | \mu x. \psi | \nu x. \psi | Next \psi$ 

K,  $\pi$ , V  $\vDash \psi$  is defined for a path  $\pi$ .





### 2-valued Path Semantics

Path Semantics = Path Construction + State Semantics







































#### L-valued Path Semantics







#### L-valued Path Semantics







#### Is simulation lifted ?







# Is simulation lifted ?

- No. We found a counterexample.
- We gave a sufficient condition:
  - A simulation is lifted if L is the open sets of a topological space and closed for countable intersections.
  - Examples: power sets, Nat  $\cup \{\omega\}$
- Under the condition, the simulation theorem for path semantics holds.





# Conclusion

- Complete Heyting algebra valued
  - Transition System, Kripke Model, Simulation
  - State Semantics for Modal  $\mu$  -Calculus, Simulation Theorem
- Under our new condition
  - Path Semantics for Linear Modal  $\mu$  -Calculus, Simulation Theorem





#### Future Work

• To relate this work to fuzzy relations or probabilistic relations