

## LEIBNIZ INTERPOLATION PROPERTIES

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In the framework of Abstract Algebraic Logic (see [2] for a general survey) a number of different interpolation properties have been considered: Maehara Interpolation Property, Robinson Interpolation Property, and Deductive Interpolation Property among others (see [1] as well as the references therein).

In [3], Kihara and Ono introduce some new notions of interpolation for substructural logics. For instance, one of the interpolation properties that they study is the following: for any arbitrary pair of formulas  $\varphi, \psi$ , if  $\vdash (\varphi \backslash \psi) \wedge (\psi \backslash \varphi)$ , then there exists a formula  $\delta$  with variables  $\text{var}(\delta) \subseteq \text{var}(\varphi) \cap \text{var}(\psi)$  such that

$$\vdash (\varphi \backslash \delta) \wedge (\delta \backslash \varphi) \quad \text{and} \quad \vdash (\delta \backslash \psi) \wedge (\psi \backslash \delta).$$

In the same paper they present algebraic characterizations for all these notions and investigate their relation with the usual interpolation properties. The authors also remark: “only a few properties specific to FL-algebras or residuated lattices are used in our discussion,” and they propose the study of these interpolation properties in a more general setting. The objective of the present paper is to give an answer to that remark by exhibiting a general framework in which these interpolation properties can be expressed and characterized by means of algebraic tools.

First, we note that  $\nabla(x, y) = \{(x \backslash y) \wedge (y \backslash x)\}$  is a set of equivalence formulas (in the sense of Abstract Algebraic Logic) for the substructural logics, that is, it defines the Leibniz congruences of their models. In particular, for every theory  $\Gamma$  and every pair of formulas  $\varphi, \psi$ ,

$$\Gamma \vdash \nabla(\varphi, \psi) \quad \Leftrightarrow \quad \langle \varphi, \psi \rangle \in \mathbf{\Omega}\Gamma.$$

This motivates the choice of the equivalential logics, i.e., logics having a set of equivalence formulas, as our general setting.

We introduce a family of what we call ‘Leibniz’ interpolation properties, in which we are mainly interpolating pairs of elements in the Leibniz congruence of a particular free algebra. Applied to the case of substructural logics, some of them coincide with those introduced in [3].

The first notion of interpolation that we introduce is the *Leibniz Interpolation Property* (LIP). A sentential logic  $\mathcal{S}$  has the LIP if for every pair of formulas  $\varphi, \psi$ , if  $\langle \varphi, \psi \rangle \in \mathbf{\Omega}\mathcal{S}$ , then there exists a formula  $\delta$  with variables  $\text{var}(\delta) \subseteq \text{var}(\varphi) \cap \text{var}(\psi)$  such that  $\langle \varphi, \delta \rangle, \langle \delta, \psi \rangle \in \mathbf{\Omega}\mathcal{S}$ . Here,  $\mathbf{\Omega}\mathcal{S}$  stands for the Leibniz congruence of the set of theorems of  $\mathcal{S}$ .

By adding new conditions we can strengthen this property in several different ways. For instance, we say that a sentential logic  $\mathcal{S}$  has the *Leibniz Interpolation Property for Theories* (LIPT) if for every set of formulas  $\Gamma$  and every pair of formulas  $\varphi, \psi$ , if  $\langle \varphi, \psi \rangle \in \mathbf{\Omega}\bar{\Gamma}$ , then there exists a formula  $\delta$  with variables  $\text{var}(\delta) \subseteq \text{var}(\Gamma, \varphi) \cap \text{var}(\Gamma, \psi)$  such that  $\langle \varphi, \delta \rangle, \langle \delta, \psi \rangle \in \mathbf{\Omega}\bar{\Gamma}$ , where  $\bar{\Gamma} = \text{Cn}_{\mathcal{S}}(\Gamma)$  is the theory of  $\mathcal{S}$  generated by  $\Gamma$ .

In the context of equivalential logics, we find characterizations for these interpolations properties in terms of categorical properties of their classes of reduced

models, reduced algebras, or both. We say that a concrete category  $\mathbf{C}$  has *covering pullbacks for (pairs of) monos* if whenever  $f : A \hookrightarrow B$  and  $g : A \hookrightarrow C$  is a pair of monos in  $\mathbf{C}$ , there exists a pullback and onto morphisms  $\rho_A, \rho_B, \rho_C$  rendering commutative the following diagram:

$$(1) \quad \begin{array}{ccccc} & & B & \xleftarrow{\rho_B} & B' \\ & \nearrow f & & & \nearrow & \\ A & \xleftarrow{\rho_A} & A' & \xrightarrow{\rho_C} & C' & \searrow & \\ & \searrow g & & & \searrow & \\ & & C & \xleftarrow{\rho_C} & C' & \nearrow & \\ & & & & & & E \end{array}$$

We obtain the following result:

**Theorem 1.** *For an equivalential logic  $\mathcal{S}$ , the following statements are equivalent:*

- (i)  $\mathcal{S}$  has the *LIP*,
- (ii) the category  $\mathbf{Mod}^* \mathcal{S}$  of its reduced models has covering pullbacks for monos,
- (iii) the category  $\mathbf{Alg}^* \mathcal{S}$  of its reduced algebras has covering pullbacks for monos.

For the others interpolation properties we find similar characterizations, by strengthening the conditions in Diagram (1). For instance, the *LIPT* for an equivalential logic  $\mathcal{S}$  is equivalent to the existence in  $\mathbf{Mod}^* \mathcal{S}$  of covering pullbacks for monos of the form of (1) where, moreover,  $\rho_A$  is an isomorphism.

Three of the notions that we introduce and deserve a special attention are the *Relative Leibniz*, *Robinson-Leibniz*, and *Maehara-Leibniz Interpolation Properties*. They are the natural generalizations of three of the interpolation properties considered in [3] for substructural logics. We investigate the relation of these properties with those studied by Czelakowsky and Pigozzi in [1] and obtain the following result for the equivalential logics satisfying the *G-rule*, that is, the regularly algebraizable logics.

**Theorem 2.** *For every regularly equivalential logic  $\mathcal{S}$ , we have the following implications:*

- (i) *The Relative Interpolation Property implies Interpolation Property.*
- (ii) *The Robinson-Leibniz Interpolation Property implies the Robinson Interpolation Property.*
- (iii) *The Maehara-Leibniz Interpolation Property implies the Maehara Interpolation Property.*

## REFERENCES

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