Contents

Domain Engineering

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Before software can be designed we must know its requirements. Before requirements can be expressed we must understand the domain. So it follows, from our dogma, that we must first establish precise descriptions of domains; then from such, "derive" at least domain requirements; and from those and other (interface and machine) requirements design the software, or, more generally, the computing systems. The preceding was an engineering dogma. Now follows a science dogma: *Just as physicists have studied this universe for centuries (and more), and will continue to do so for centuries (or more), so it is about time that we also study such man-made universes as air traffic, the financial service industry, healthcare, manufacturing, "the market", railways, indeed transportation in general, and so forth. Just in-and-by-themselves. No need to excuse such a study by stating only engineering concerns. To understand is all. And helps engineering.*

In the main part of this chapter we shall outline what goes into a domain description, not so much how we acquire what goes in. That is: before we can acquire domain "knowledge" we must know what are suitable structures of domain descriptions. Thus we shall outline ideas of modelling the intrinsics (of a domain), the support technologies (of ...), the management and organisation (of ...), the rules and regulations (including [license or contract] scripts) (of ...), and the human behaviours (of a domain).

Before delving into the main part we shall, however first overview what we see as basic principles of describing phenomena and concepts of domains.

At the basis of all descriptional, prescriptional and specificational modeling work is abstraction. Mathematics is a reliable carrier of abstraction. Hence our domain modeling will be presented both as informal, yet short and precise, that is, concise narratives as well as formally. In this chapter we primarily express the formalisations in the RAISE [28] specification language, RSL [26]. We refer to [6, 7] for a comprehensive coverage of formal abstractions and models.

Two remarks are now in order. Firstly, there are other specification (cum development or design) languages, Alloy [44], ASM [14, 59], B [1], CafeOBJ [25, 24, 21], CASL [5, 18, 53], VDM-SL [10, 11, 23] and Z [68, 69, 75, 38]. But, secondly, none of these suffices. Each, including RSL, have their limitations in what they were supposed to express with ease. So one needs to combine, to integrate either of the above formal notations with for example the notations of Duration Calculus [77, 78] (DC), Message [41, 42, 43] or Live Sequence Charts (MSCs and LSCs) [19, 36, 47], Petri Nets [46, 54, 56, 57, 58], Statecharts [32, 33, 35, 37, 34] (SCs), TLA+ [48, 49, 51] etcetera.

Chapters 12–15 of [7] presents an extensive coverage of Petri Nets, MSCs and LSCs, SCs and DCs, respectively. This chapter presents an essence of chapters 5, 6 and 11 of [8].

The forthcoming book "Logics of Specification Languages" [9] covers the following formal notations ASM, B, CafeOBJ, CASL, DC, RAISE, VDM-SL, TLA+ and Z. [9] represents an extensive revision of the following published papers: [59, 16, 21, 53, 27, 51, 38].

1 Introduction

1.1 Application cum Business Domains

By domain we shall loosely speaking understand the same as by apoplication (or business) domain: a universe of discourse, an area of human and societal activitity for which, eventually some support in the form of computing and (electronic) communication (c&c) may be desired. Once such c&c, that is, hardware and software has been installed the environment for that machine is the former domain, and the new domain includes the machine, i.e., the hardware and software c&c. The machine interacts with its new domain (with its environment). But we can speak of a domain without ever thinking, or having to think about c&c applications.

Examples of domains are: air traffic, airports, the financial service industry (clients, banks, brokers, securities exchange, portfolio managers, insurance companies, credit card companies, financial service industry "watchdogs", etc.), freight logistics, health care, manufacturing, and transportation (air lines, railways, roads (private automobiles, busses, taxis, trucking), shipping). These examples could be said to be "grand scale", and to reflect infrastructure components of a society. Less "grand" examples of domains are: the interconnect cabling of and the electronic and electro-mechanical boxes of either a sound & noise measuring test set-up or a so-called "intelligent" home, the interlocking of groups of rail points (switches) of a railway station, an automobile (with all its mechanical, electro-mechanical and electronic parts and the composition of all these parts, and with the driver and zero, one or more passengers), or a set or laws (or just rules and regulations) of business enterprise (the domain would then, in this latter cases, include abstractions, relevant to the laws (etc.) of whatever these laws or rules and regulations referred to). In all of these latter cases we usually include the human or technological monitoring and control of these domains.

1.2 Physics, Domains and Engineering

Physics, "mother nature" has been studied for millenia. Physicists continue to unravel deeper and deeper understandings of the physically observable universe around us. Classical engineering builds on physics. Every aeronautical & aerospace, chemical, civil, electrical & electronics, and mechanical & control engineer is expected to know all the laws of physics relevant to their field, and much more, and is expected to model, using various branches of mathematics (calculus, statistics, probability theory, graph theory, combinatorics, etc.), phenomena of the domain in which their engineering artifacts are placed (as well as, of course, these artifacts themselves).

Software engineers sometimes know how to model their own artifacts (compilers, operating systems, database management systems, data communication systems, web services, etc.), but they seldomly, if ever are expected to model, and they mostly cannot, i.e., do not know how to model the domain in which their software operates.

2 Domain Engineering: The Engineering Dogma

- *Before software can be designed we must know its requirements.*
- *Before requirements can be expressed we must understand the domain.*
- *So it follows, from our dogma, that we must*
  - *first establish precise descriptions of domains;*
  - *then from such, "derive" at least domain requirements;*
  - *and from those and other (interface and machine) requirements design the software, or, more generally, the computing systems.*

That is, we propose what we have practiced for many years, that the software engineering process be composed — and that it be iterated over, in a carefully monitored and controlled manner — as follows:

- domain engineering,
- requirements engineering, and
- software design.

We see the domain engineering process as composed from, and iterated over:

1. identification of and regular interaction with stakeholders
2. domain (knowledge) acquisition
3. domain analysis
4. domain modelling

5. domain verification
6. domain validation
7. domain theory formation

In this chapter we shall only look at the principles and techniques of domain modeling, that is, item 4. To pursue items 2.–3. one must know what goes into a domain description, i.e., a domain model.

- *A major part of the domain engineering process is taken up by finding and expressing suitable abstractions, that is, descriptions of the domain.*
- *Principles for identifying, classifying and describing domain phenomena and concepts are therefore needed.*

This chapter focuses on presenting some of these principles and techniques.

3 Entities, Functions, Events and Behaviours

In the domain we observe phenomena. From usually repeated such observations we form (immediate, abstract) concepts. We may then "lift" such immediate abstract concepts to more general abstract concepts.

Phenomena are manifest. They can be observed by human senses (seen, heard, felt, smelled or tasted) or by physical measuring instruments (mass, length, time, electric current, thermodynamic temperature, amount of substance, luminous intensity). Concepts are defined.

We shall analyse phenomena and concepts according to the following simple, but workable classification: **entities**, **functions** (over entities), **events** (involving changes in entities, possibly as caused by function invocations, i.e., **actions**, and/or possibly causing such), and **behaviours** as (possibly sets of) sequences of actions (i.e., function invocations) and events.

3.1 Entities

- *By an* **entity** *we shall understand something static; although that "thing" may move, after it has moved it is essentially the same thing, an entity*

Entities are either atomic or composite. The decision as to which entities are considered what is a decision solely taken by the describer.

Atomic Entities

- *By an* **atomic entity** *we intuitively understand an entity which "cannot be taken apart" (into other, the sub-entities).*

Attributes — Types and Values:

With any entity we can associate one or more attributes.

- *By an* **attribute** *we understand a pair of a* **type** *and a* **value***.*

Example 1. **Atomic Entities:**

Entity: **Person**		Entity: **Bank Account**	
Type	**Value**	**Type**	**Value**
Name	Dines Bjørner	number	212 023 361 918
Weight	118 pounds	balance	1,678,123 Yen
Height	179 cm	interest rate	1.5 %
Gender	male	credit limit	400,000 Yen

•

"Removing" an attribute from an entity destroys its "entity-hood".

Composite Entities

- *By a composite entity we intuitively understand an entity (i) which "can be taken apart" into sub-entities, (ii) where the composition of these is described by its* **mereology***, and (iii) which further possess one or more attributes.*

Example 2. **Transport Net, A Narrative:**

Entity: **Transport Net**		
Subentities:	Segments	
	Junctions	
Mereology:	"set" of one or more s(egment)s and	
	"set" of two or more j(unction)s	
such that	each s(egment) is delimited by two j(unctions)	
and such that	each j(unction) connects one or more s(egments)	
Attributes		
	Types:	**Values:**
	Multimodal	Rail, Roads
	Transport Net of	Denmark
	Year Surveyed	2006

•

To put the above example of a composite entity in context we give an example of both an informal narrative and a correspondin formal specification:

Example 3. **Transport Net, A Formalisation:** A transport net consists of one or more segments and two or more junctions. With segments [junctions] we can associate the following attributes: segment [junction] identifiers, the identifiers of the two junctions to which segments are connected [the identifiers of the one or more segments connected to the junction], the mode of a segment [the modes of the segments connected to the junction].

type
  N, S, J, Si, Ji, M
value
  obs_Ss: N → S-**set**,   obs_Js: N → J-**set**
  obs_Si: S → Si,   obs_Ji: J → Ji
  obs_Jis: S → Ji-**set**,   obs_Sis: J → Si-**set**
  obs_M: S → M,   obs_Ms: J → M-**set**
axiom
  ∀ n:N • **card** obs_Ss(n) ≥ 1 ∧ **card** obs_Js(n) ≥ 2
  ∀ n:N • **card** obs_Ss(n) ≡ **card** {obs_Si(s)|s:S • s ∈ obs_Ss(n)}
  ∀ n:N • **card** obs_Js(n) ≡ **card** {obs_Ji(c)|j:J • j ∈ obs_Js(n)}
  ...
type
  Nm, Co, Ye
value
  obs_Nm: N → Nm, obs_Co: N → Co, obs_Ye: N → Ye

Si, Ji, M, Nm, Co, Ye are not entities. They are names of attribute types and, as such, designate attribute values. N is composite, S and J are considered atomic. •

- *By* **mereology** *we shall understand a theory of part-hood relations. That is, of the relations of part to whole and the relations of part to part within a whole.*

The term mereology seems to have been first used in the sense we are using ot my the Polish mathematical logician Stanisław Leshniewski [50, 52, 70, 71, 72].

States

- *By a domain* **state** *we shall understand a collection of domain entities chosen by the domain engineer.*

The pragmatics of the notion of state is that states are recurrent arguments to functions and are changed by function invocations.

3.2 Functions

- *By a* **function** *we shall understand something which when applied to some arguments (i.e., entities) yield some entities called the result of the function (application).*
- *By an* **action** *we shall understand the same things as applying a state-changing function to its arguments (including the state).*

[The observer functions of the formal example above are not the kind of functions we are seeking to identify in the domain. They are mere technicalities: needed, due to the way in which we formalise — and are deployed in order to express sub-entities, mereologies and attributes.]

Function Signatures

By a function signature we mean the name and type of a funtion.

type
  A, B, ..., C, X, Y, .., Z
value
  f: A $\times$ B $\times$ ... $\times$ C $\to$ X $\times$ Y $\times$ ... $\times$ Z

The last line above expresses a schematic function signature.

Function Descriptions

By a function description we mean a function signature and something which describes the relationship between function arguments (the a:A's, b:B's, ..., c:C's and the x:X's, y:Y's, ..., z:Z's).

Example 4. **Well Formed Routes:**

type
  P = Ji $\times$ Si $\times$ Ji   /$*$ path: triple of identifiers $*$/
  R$'$ = P* /$*$ route: sequence of connected paths $*$/
  R = {| r:R$'$ • wf_R(r) |} /$*$ subtype of R$'$: those r$'$s satisfying wf_R(r) $*$/
value
  wf_R: R$'$ $\to$ **Bool**
  wf_R(r) $\equiv$
    $\forall$ i:**Nat**•{i,i+1}$\subseteq$**inds** r$\Rightarrow$**let** (,,ji$'$)=r(i),(ji$''$,,)=r(i+1) **in** ji$'$=ji$''$ **end**

•

The last line above describes the route wellformedness predicate. [The meaning of the "(,," and ",,)" is that the omitted path components "play no rôle.]

3.3 Events

- *By an* **event** *we shall understand an instantaneous change of state not directly brought about by some explicitly willed action in the domain, but either by "external" forces. or implicitly as a non-intended result of an explicitly willed action.*

Events may or may not lead to the initiation of explicitly issued operations.

Example 5. **Events:** A 'withdraw' from a positive balance bank account action may leave a negative balance bank account. A bank branch office may have to temporarily stop actions, i.e., close, due to a bank robbery. •

Internal events: The first example above illustrates an internal action. It was caused by an action in the domain, but was not explicitly the main intention of the "woithdraw" function.

External events: The second example above illustrates an external action. We assume that we have not explicitly modelled bank robberies!

Formal modelling of events: With every event we can associate an event label. And event label can be thought of as a simple identifier. Two or more event labels may be the same.

3.4 Behaviours

- *By a* **behaviour** *we shall understand a structure of actions (i.e., function invocations) and events. The structure may typically be a set of sequences of actions and events.*

A behaviour is either a simple behaviour, or is a concurrent behaviour, or, if the latter, can be either a communicating behaviour or not.

- *By a* **simple behaviour** *we shall understand a sequence of actions and events.*

Example 6. **Simple Behaviours:** The opening of a bank account, the deposit into that bank account, zero, one or more other such deposits, a withdrawal from the bank account in question, etc. (deposits and withdrawals), ending with a closing of the bank account. Any prefix of such a sequence is also a simple behaviour. Any sequence in which one or more events are interspersed is also a simple behaviour. •

- *By a* **concurrent behaviour** *we shall understand a set of behaviours (simple or otherwise).*

Example 7. **Concurrent Behaviours:** A set of simple behaviours may result from two or more distinct bank clients, each operating of their own, distinct, that is, non-shared accounts. •

- *By a* **communicating behaviour** *we shall understand a set of two or more behaviours where otherwise distinct elements (i.e., behaviours) share events.*

The sharing of events can be identified via the event labels.

Example 8. **Communicating Behaviours:** To model that two or more clients can share the same bank account one could model the bank account as one behaviour and each client as a distinct behaviour. Let us assume that only one client can open an account and that only one client can close an account. Let us further assume that sharing is brought about by one client, say the one who opened the account, identifying the sharing clients. Now, in order to make sure that at most one client accesses the shared account at one one time (in any one "smallest" transaction interval) one may model "client access to account" as a pair of events such that during the interval between the first (begin transaction) and the second (end transaction) event no other client can share events with the bank account behaviour. Now the set of behaviours of the bank account and one or more of the client behaiours is an example of a communicating behavior. •

Formal modelling of behaviours: Communicating behaviours, the only really interesting behaviours, can be modelled in a great variety of ways: from set-oriented models in B [1], RSL [26, 28, 6, 7, 8], VDM [10, 11, 23], or Z [68, 69, 75, 38], to models using for example CSP [39, 63, 66, 40] (as for example "embedded" in RSL [26]), or, to diagram models using, for example, Petri nets [46, 54, 56, 57, 58], message [41, 42, 43] or live sequence charts [19, 36, 47], or statecharts [32, 33, 35, 37, 34].

3.5 Discussion

The main aim of Sect. 3 is to ensure that we have a clear understanding of the modelling concepts of entities, functions, events and behaviours. To "reduce" the modelling of phenomena and concepts to these four kinds of phenomena and concepts is, of course, debatable. Our point is that it works, that further classification, as is done in for example John F. Sowa's [67], is not necessary, or rather, is replaced by how we model attributes of for example entities[3] and how we model facets, such as we shall call them. The modelling of facets is the main aim of this chapter.

4 Domain Facets

- *By a domain* **facet** *we shall understand one amongst a finite set of generic ways of analysing a domain: a view of the domain, such that the different facets cover conceptually different views, and such that these views together cover the domain.*

The hedge here is "finite set of generic ways". Thus there is an assumption, a conjecture to be possibly refuted. Namely the postulate that there is a finite number of facets. We shall offer the following facets: intrinsics, sipport technology, management and organisation, rules and regulations (and scripts), and human behaviour.

4.1 Intrinsics

- *By domain* **intrinsics** *we shall understand those phenomena and concepts of a domain which are basic to any of the other facets (listed earlier and treated,*

[3] For such issues as static and dynamic attributes, dimensionality, tangibility, time and space, etc., we refer to Michael A. Jackson's [45] or Chap. 10 of [8].

in some detail, below), with such domain intrinsics initially covering at least one specific, hence named, stakeholder view.

Example 9. **Railway Net Intrinsics:** We narrate and formalise three railway net intrinsics.

- From the view of *potential train passengers* a railway net consists of lines, stations and trains. A line connects exactly two distinct stations.
- From the view of *actual train passengers* a railway net — in addition to the above — allows for several lines between any pair of stations and, within stations, provides for one or more platform tracks from which to embark or alight a train.
- From the view of *train operating staff* a railway net — in addition to the above — has lines and stations consisting of suitably connected rail units. A rail unit is either a simple (i.e., linear, straight) unit, or is a switch unit, or is a simple crossover unit, or is a switchable crossover unit, etc. Simple units have two connectors. Switch units have three connectors. Simple and switchable crossover units have four connectors. A path (through a unit) is a pair of connectors of that unit. A state of a unit is the set of paths, in the direction of which a train may travel. A (current) state may be empty: The unit is closed for traffic. A unit can be in either one of a number of states of its state space.

Railway Net Intrinsics A summary formalisation of the three narrated railway net intrinsics could be:

- *Potential train passengers:*

  scheme N0 =
    class
      type
        N, L, S, Sn, Ln
      value
        obs_Ls: N → L-**set**, obs_Ss: N → S-**set**
        obs_Ln: L → Ln, obs_Sn: S → Sn
        obs_Sns: L → Sn-**set**, obs_Lns: S → Ln-**set**
      axiom
        ...
    end

  N, L, S, Sn and Ln designate nets, lines, stations, station names and line names. One can observe lines and stations from nets, line and station names from lines and stations, pair sets of station names from lines, and lines names (of lines) into and out from a station from stations. Axioms ensure proper graph properties of these concepts.

- *Actual train passengers:*

scheme N1 = **extend** N0 **with**
  class
    type
      Tr, Trn
    value
      obs_Trs: S → Tr-**set**, obs_Trn: Tr → Trn
    axiom
      ...
  end

The only additions are that of track and track name sorts, related observer functions and axioms.

- *Train operating staff:*

scheme N2 = **extend** N1 **with**
  class
    type
      U, C
      $P' = U \times (C{\times}C)$
      P = {| p:P′ • **let** (u,(c,c′))=p **in** (c,c′)∈ ∪ obs_Ω(u) **end** |}
      Σ = P-**set**
      Ω = Σ-**set**
    value
      obs_Us: (N|L|S) → U-**set**
      obs_Cs: U → C-**set**
      obs_Σ: U → Σ
      obs_Ω: U → Ω
    axiom
      ...
  end

Unit and connector sorts have been added as have concrete types for paths, unit states, unit state spaces and related observer functions, including unit state and unit state space observers. The reader is invited to compare the three narrative descriptions with the three formal descriptions, line by line. •

Different stakeholder perspectives, not only of intrinsics, as here, but of any facet, lead to a number of different models. The name of a phenomenon of one perspective, that is, of one model, may coincide with the name of a "similar" phenomenon of another perspective, that is, of another model, and so on. If the intention is that the "same" names cover comparable phenomena, then the developer must state the comparison relation.

Example 10. **Comparable Intrinsics:** We refer to Example 9. We claim that the concept of nets, lines and stations in the three models of Example 9 must relate. The simplest possible relationships are to let the third model be the common "unifier" and to mandate

- that the model of nets, lines and stations of the *potential train passengers* formalisation is that of nets, lines and stations of the *train operating staff* model; and
- that the model of nets, lines, stations and tracks of the *actual train passengers* formalisation is that of nets, lines, stations of the *train operating staff* model.

Thus the third model is seen as the definitive model for the stakeholder views initially expressed. •

Example 11. **Intrinsics of Switches:** The intrinsic attribute of a rail switch is that it can take on a number of states. A simple switch $(^{c_|}Y_c^{c_/})$ has three connectors: $\{c, c_|, c_/\}$. c is the connector of the common rail from which one can either "go straight" $c_|$, or "fork" $c_/$ (Fig. 1). So we have that a possible state space of such a switch could be ω_{g_s} :

$$
\begin{array}{l}
\{\{\},\\
\{(c,c_|)\},\{(c_|,c)\},\{(c,c_|),(c_|,c)\},\\
\{(c,c_/)\},\{(c_/,c)\},\{(c,c_/),(c_/,c)\},\{(c_/,c),(c_|,c)\},\\
\{(c,c_|),(c_|,c),(c_/,c)\},\{(c,c_/),(c_/,c),(c_|,c)\},\{(c_/,c),(c,c_|)\},\{(c,c_/),(c_|,c)\}\}
\end{array}
$$

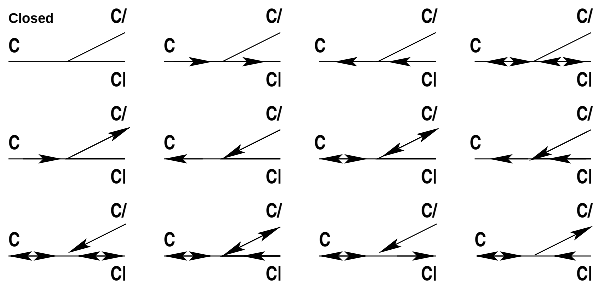

Fig. 1. Possible states of a rail switch

The above models a general switch ideally. Any particular switch ω_{p_s} may have $\omega_{p_s} \subset \omega_{g_s}$. Nothing is said about how a state is determined: who sets and resets it, whether determined solely by the physical position of the switch gear, or also by visible or virtual (i.e., invisible, intangible) signals up or down the rail, away from the switch. •

Conceptual Versus Actual Intrinsics

In order to bring an otherwise seemingly complicated domain across to the reader, one may decide to present it piecemeal:[4] First, one presents the very

[4] That seemingly complicated domain may seem very complicated, containing hundreds of entities, functions and behaviours. Instead of presenting all the entities,

basics, the fewest number of inescapable entities, functions and behaviours. Then, in a step of enrichment, one adds a few more (intrinsic) entities, functions and behaviours. And so forth. In a final step one adds the last (intrinsic) entities, functions and behaviours. In order to develop what initially may seem to be a complicated domain, one may decide to develop it piecemeal: We basically do as for the presentation steps: Steps of enrichments — from a big lie, via increasingly smaller lies, till one reaches a truth!

On Modelling Intrinsics

Domains can be characterised by intrinsically being entity, or function, or event, or behaviour intensive. Software support for activities in such domains then typically amount to database systems, computation-bound systems, real-time embedded systems, respectively distributed process monitoring and control systems. Modelling the domain intrinsics in respective cases can often be done property-oriented specification languages (like CafeOBJ [25, 24, 21] or CASL [5, 18, 53]), model-oriented specification languages (like B [1], VDM-SL [10, 11, 23], RSL [26], or Z [68, 69, 75, 38]), event-based languages (like Petri nets [46, 54, 56, 57, 58] or CSP [39, 63, 66, 40]), respectively process-based specfication languages (like MSCs [41, 42, 43], LSCs [19, 36, 47], statecharts [32, 33, 35, 37, 34], or CSP [39, 63, 66, 40]).

4.2 Support Technologies

- *By a domain* **support technology** *we shall understand ways and means of implementing certain observed phenomena or certain conceived concepts.*

Example 12. **Railway Support Technology:** We give a rough sketch description of possible rail unit switch technologies.

(i) In "ye olde" days, rail switches were "thrown" by manual labour, i.e., by railway staff assigned to and positioned at switches.

(ii) With the advent of reasonably reliable mechanics, pulleys and levers[5] (and steel wires), switches were made to change state by means of "throwing" levers in a cabin tower located centrally at the station (with the lever then connected through wires etc., to the actual switch).

(iii) This partial mechanical technology then emerged into electromechanics, and cabin tower staff was "reduced" to pushing buttons.

(iv) Today, groups of switches, either from a station arrival point to a station track, or from a station track to a station departure point, are set and reset by means also of electronics, by what is known as interlocking (for example, so that two different routes cannot be open in a station if they cross one another). •

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functions, events and behaviours in one "fell swoop", one presents them in stages: first, around seven such (entities, functions, events and behaviours), then seven more, etc.

It must be stressed that Example 12 is just a rough sketch. In a proper narrative description the software (cum domain) engineer must describe, in detail, the subsystem of electronics, electromechanics and the human operator interface (buttons, lights, sounds, etc.).

An aspect of supporting technology includes recording the state-behaviour in response to external stimuli. We give an example.

Example 13. **Probabilistic Rail Switch Unit State Transitions:** Figure 2 indicates a way of formalising this aspect of a supporting technology. Figure 2 intends to model the probabilistic (erroneous and correct) behaviour of a switch when subjected to settings (to switched (s) state) and resettings (to direct (d) state). A switch may go to the switched state from the direct state when subjected to a switch setting s with probability psd. •

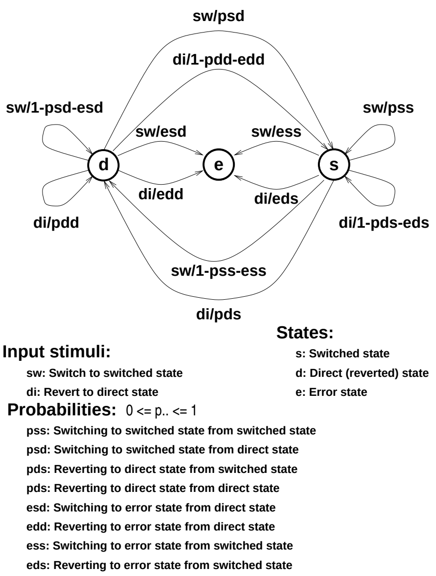

Fig. 2. Probabilistic state switching

Another example shows another aspect of support technology: Namely that the technology must guarantee certain of its own behaviours, so that software designed to interface with this technology, together with the technology, meets dependability requirements.

Example 14. **Railway Optical Gates:** Train traffic (itf:iTF), intrinsically, is a total function over some time interval, from time (t:T) to continuously positioned (p:P) trains (tn:TN).

Conventional optical gates sample, at regular intervals, the intrinsic train traffic. The result is a sampled traffic (stf:sTF). Hence the collection of all optical gates, for any given railway, is a partial function from intrinsic to sampled train traffics (stf).

We need to express quality criteria that any optical gate technology should satisfy — relative to a necessary and sufficient description of a closeness predicate. The following axiom does that:

> For all intrinsic traffics, itf, and for all optical gate technologies, og, the following must hold: Let stf be the traffic sampled by the optical gates. For all time points, t, in the sampled traffic, those time points must also be in the intrinsic traffic, and, for all trains, tn, in the intrinsic traffic at that time, the train must be observed by the optical gates, and the actual position of the train and the sampled position must somehow be checkable to be close, or identical to one another.

Since units change state with time, n:N, the railway net, needs to be part of any model of traffic.

Railway Optical Gate Technology Requirements

type
  T, TN
  $P = U^*$
  NetTraffic == net:N trf:(TN $\overrightarrow{m}$ P)
  iTF = T $\rightarrow$ NetTraffic
  sTF = T $\overrightarrow{m}$ NetTraffic
  oG = iTF $\tilde{\rightarrow}$ sTF
value
  [close] c: NetTraffic $\times$ TN $\times$ NetTraffic $\tilde{\rightarrow}$ **Bool**
axiom
  $\forall$ itt:iTF, og:OG • **let** stt = og(itt) **in**
    $\forall$ t:T • t $\in$ **dom** stt •
      t $\in$ $\mathcal{D}$ itt $\wedge$ $\forall$ Tn:TN • tn $\in$ **dom** trf(itt(t))
        $\Rightarrow$ tn $\in$ **dom** trf(stt(t)) $\wedge$ c(itt(t),tn,stt(t)) **end**

$\mathcal{D}$ is not an RSL operator. It is a mathematical way of expressing the definition set of a general function. Hence it is not a computable function. Checkability is an issue of testing the optical gates when delivered for conformance to the closeness predicate, i.e., to the axiom. •

On Modelling Support Technologies

Support technologies in their relation to the domain in which they reside tyically reflect real-time embeddedness. As such the techniques and languages for modelling support technologies resemble those for modelling event and process intensity, while temporal notions are brought into focus. Hence typical modelling notations include event-based languages (like Petri nets [46, 54, 56, 57, 58] or CSP [39, 63, 66, 40]), respectively process-based specification languages (like MSCs [41, 42, 43], LSCs [19, 36, 47], statecharts [32, 33, 35, 37, 34], or CSP [39, 63, 66, 40]), as well as temporal languages (like the Duration Calculus [77, 78] and Temporal Logic of Actions, TLA+, [48, 49, 51]).

4.3 Management and Organisation

Example 15. **Train Monitoring, I:** In China, as an example, rescheduling of trains occurs at stations and involves telephone negotiations with neighbouring stations ("up and down the lines"). Such rescheduling negotiations, by phone, imply reasonably strict management and organisation (M&O). This kind of M&O reflects the geographical layout of the rail net. •

- *By domain* **management** *we shall understand such people (such decisions) (i) who (which) determine, formulate and thus set standards (cf. rules and regulations, Sect. 4.4) concerning strategic, tactical and operational decisions; (ii) who ensure that these decisions are passed on to (lower) levels of management, and to floor staff; (iii) who make sure that such orders, as they were, are indeed carried out; (iv) who handle undesirable deviations in the carrying out of these orders cum decisions; and (v) who "backstop" complaints from lower management levels and from floor staff.*
- *By domain* **organisation** *we shall understand the structuring of management and nonmanagement staff levels; the allocation of strategic, tactical and operational concerns to within management and nonmanagement staff levels; and hence the "lines of command": who does what, and who reports to whom, administratively and functionally.*

Example 16. **Railway Management and Organisation: Train Monitoring, II:** We single out a rather special case of railway management and organisation. Certain (lowest-level operational and station-located) supervisors are responsible for the day-to-day timely progress of trains within a station and along its incoming and outgoing lines, and according to given timetables. These supervisors and their immediate (middle-level) managers (see below for regional managers) set guidelines (for local station and incoming and outgoing lines) for the monitoring of train traffic, and for controlling trains that are either ahead of or behind their schedules. By an incoming and an outgoing line we mean part of a line between two stations, the remaining part being handled by neighbouring station management. Once it has been decided, by such a manager, that a train is not following its schedule, based on information monitored by nonmanagement staff, then that manager directs that staff: (i) to suggest a new schedule for the train in question, as well as for possibly affected other trains, (ii) to negotiate the new schedule with appropriate neighbouring stations, until a proper reschedule can be decided upon, by the managers at respective stations, (iii) and to enact that new schedule.[6] A (middle-level operations) manager for regional traffic, i.e., train traffic involving several stations and lines, resolves possible disputes and conflicts. •

The above, albeit rough-sketch description, illustrated the following management and organisation issues: There is a set of lowest-level (as here: train traffic scheduling and rescheduling) supervisors and their staff. They are organised into one such group (as here: per station). There is a middle-level (as here: regional train traffic scheduling and rescheduling) manager (possibly with some small staff), organised with one such per suitable (as here: railway) region. The guidelines issued jointly by local and regional (...) supervisors and managers imply an organisational structuring of lines of information provision and command.

Conceptual Analysis, First Part

People staff enterprises, the components of infrastructures with which we are concerned, i.e., for which we develop software. The larger these enterprises — these infrastructure components — the more need there is for management and organisation. The role of management is roughly, for our purposes, twofold: first, to perform strategic, tactical and operational work, to set strategic, tactical and operational policies — and to see to it that they are followed. The role of management is, second, to react to adverse conditions, that is, to unforeseen situations, and to decide how they should be handled, i.e., conflict resolution.

Policy setting should help nonmanagement staff operate normal situations — those for which no management interference is thus needed. And management "backstops" problems: management takes these problems off the shoulders of nonmanagement staff.

To help management and staff know who's in charge wrt. policy setting and problem handling, a clear conception of the overall organisation is needed. Organisation defines lines of communication within management and staff, and between these. Whenever management and staff has to turn to others for assistance they usually, in a reasonably well-functioning enterprise, follow the command line: the paths of organigrams — the usually hierarchical box and arrow/line diagrams.

Methodological Consequences

The *management and organisation* model of a domain is a partial specification; hence all the usual abstraction and modelling principles, techniques and

[6] That enactment may possibly imply the movement of several trains incident upon several stations: the one at which the manager is located, as well as possibly at neighbouring stations.

tools apply. More specifically, management is a set of predicates, observer and generator functions which either parameterise other, the operations functions, that is, determine their behaviour, or yield results that become arguments to these other functions

Organisation is thus a set of constraints on communication behaviours. Hierarchical, rather than linear, and matrix structured organisations can also be modelled as sets (of recursively invoked sets) of equations.

Conceptual Analysis, Second Part

To relate classical organigrams to formal descriptions we first show such an organigram (Fig. 3), and then we show schematic processes which — for a rather simple scenario — model managers and the managed!

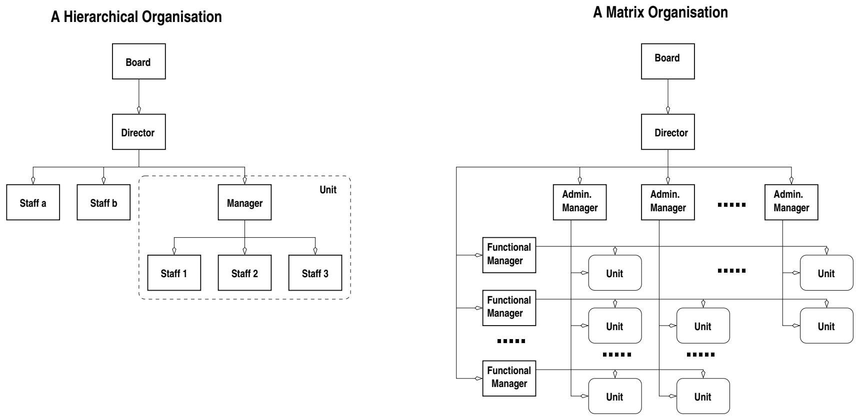

Fig. 3. Organisational structures

Based on such a diagram, and modelling only one neighbouring group of a manager and the staff working for that manager we get a system in which one manager, mgr, and many staff, stf, coexist or work concurrently, i.e., in parallel. The mgr operates in a context and a state modelled by ψ. Each staff, stf(i) operates in a context and a state modelled by sσ(i).

Conceptual Model of a Manager-Staff Relation, I

type
  Msg, Ψ, Σ, Sx
  SΣ = Sx $\overrightarrow{m}$ Σ
channel
  { ms[i]:Msg | i:Sx }
value
  sσ:SΣ, ψ:Ψ

$$\text{sys: } \mathbf{Unit} \rightarrow \mathbf{Unit}$$
$$\text{sys}() \equiv \| \{ \text{st(i)(s}\sigma\text{(i))} \mid \text{i:Sx} \} \| \text{mg}(\psi)$$

In this system the manager, mgr, (1) either broadcasts messages, m, to all staff via message channel ms[i]. The manager's concoction, m_out(ψ), of the message, msg, has changed the manager state. Or (2) is willing to receive messages, msg, from whichever staff i the manager sends a message. Receipt of the message changes, m_in(i,m)(ψ), the manager state. In both cases the manager resumes work as from the new state. The manager chooses — in this model — which of thetwo things (1 or 2) to do by a so-called nondeterministic internal choice ($\lceil\rceil$).

Conceptual Model of a Manager-Staff Relation, II

$$\text{mg: } \Psi \rightarrow \mathbf{in},\mathbf{out}\ \{\text{ms[i]}|\text{i:Sx}\}\ \mathbf{Unit}$$
$$\text{mg}(\psi) \equiv$$
(1) $(\mathbf{let}\ (\psi',\text{m})=\text{m\_out}(\psi)\ \mathbf{in}\ \|\{\text{ms[i]!m}|\text{i:Sx}\};\text{mg}(\psi')\mathbf{end})$
$\lceil\rceil$
(2) $(\mathbf{let}\ \ \psi'=\lceil\rceil\{\mathbf{let}\ \text{m=ms[i]?}\ \mathbf{in}\ \text{m\_in(i,m)}(\psi)\ \mathbf{end}|\text{i:Sx}\}\ \mathbf{in}\ \text{mg}(\psi')\ \mathbf{end})$

$$\text{m\_out: } \Psi \rightarrow \Psi \times \text{MSG},$$
$$\text{m\_in: Sx} \times \text{MSG} \rightarrow \Psi \rightarrow \Psi$$

And in this system, staff i, stf(i), (1) either is willing to receive a message, msg, from the manager, and then to change, st_in(msg)(σ), state accordingly, or (2) to concoct, st_out(σ), a message, msg (thus changing state) for the manager, and send it ms[i]!msg. In both cases the staff resumes work as from the new state. The staff member chooses — in this model — which of thetwo "things" (1 or 2) to do by a nondeterministic internal choice ($\lceil\rceil$).

Conceptual Model of a Manager-Staff Relation, III

$$\text{st: i:Sx} \rightarrow \Sigma \rightarrow \mathbf{in},\mathbf{out}\ \text{ms[i]}\ \mathbf{Unit}$$
$$\text{st(i)}(\sigma) \equiv$$
(1) $(\mathbf{let}\ \text{m} = \text{ms[i]?}\ \mathbf{in}\ \text{st(i)(stf\_in(m)}(\sigma))\ \mathbf{end})$
$\lceil\rceil$
(2) $(\mathbf{let}\ (\sigma',\text{m}) = \text{st\_out}(\sigma)\ \mathbf{in}\ \text{ms[i]!m; st(i)}(\sigma')\ \mathbf{end})$

$$\text{st\_in: MSG} \rightarrow \Sigma \rightarrow \Sigma,$$
$$\text{st\_out: } \Sigma \rightarrow \Sigma \times \text{MSG}$$

Both manager and staff processes recurse (i.e., iterate) over possibly changing states. The management process nondeterministically, external choice, "alternates" between "broadcast"-issuing orders to staff and receiving individual messages from staff. Staff processes likewise nondeterministically, external choice, alternate between receiving orders from management and issuing individual messages to management.

The conceptual example also illustrates modelling stakeholder behaviours as interacting (here CSP-like [39, 63, 66, 40]) processes.

On Modelling Management and Organisation

Management and organisation basically spans entity, function, event and behaviour intensities and thus typically require the full spectrum of modelling techniques and notations — summarised in the two "On Modelling ..." paragraphs at the end of the two previous sections.

4.4 Rules and Regulations

- *By a domain* **rule** *we shall understand some text (in the domain) which prescribes how people or equipment are expected to behave when dispatching their duty, respectively when performing their function.*
- *By a domain* **regulation** *we shall understand some text (in the domain) which prescribes what remedial actions are to be taken when it is decided that a rule has not been followed according to its intention.*

Example 17. **Trains at Stations:**

- Rule: In China the arrival and departure of trains at, respectively from, railway stations is subject to the following rule:

  In any three-minute interval at most one train may either arrive to or depart from a railway station.
- Regulation: *If it is discovered that the above rule is not obeyed*, then there is some regulation which prescribes administrative or legal management and/or staff action, as well as some correction to the railway traffic.

•

Example 18. **Trains Along Lines:**

- Rule: In many countries railway lines (between stations) are segmented into blocks or sectors. The purpose is to stipulate that if two or more trains are moving along the line, then:

  There must be at least one free sector (i.e., without a train) between any two trains along a line.
- Regulation: *If it is discovered that the above rule is not obeyed*, then there is some regulation which prescribes administrative or legal management and/or staff action, as well as some correction to the railway traffic.

•

A Meta-characterisation of Rules and Regulations

At a metalevel, i.e., explaining the general framework for describing the syntax and semantics of the human-oriented domain languages for expressing rules and regulations, we can say the following: There are, abstractly speaking, usually three kinds of languages involved wrt. (i.e., when expressing) rules and regulations (respectively when invoking actions that are subject to rules and regulations). Two languages, Rules and Reg, exist for describing rules, respectively regulations; and one, Stimulus, exists for describing the form of the [always current] domain action stimuli.

A syntactic stimulus, sy_sti, denotes a function, se_sti:STI: $\Theta \to \Theta$, from any configuration to a next configuration, where configurations are those of the system being subjected to stimulations. A syntactic rule, sy_rul:Rule, stands for, i.e., has as its semantics, its meaning, rul:RUL, a predicate over current and next configurations, $(\Theta \times \Theta) \to \mathbf{Bool}$, where these next configurations have been brought about, i.e., caused, by the stimuli. These stimuli express: If the predicate holds then the stimulus will result in a valid next configuration.

type
  Stimulus, Rule, Θ
  STI = $\Theta \to \Theta$
  RUL = $(\Theta \times \Theta) \to \mathbf{Bool}$
value
  meaning: Stimulus $\to$ STI
  meaning: Rule $\to$ RUL

  valid: Stimulus $\times$ Rule $\to \Theta \to \mathbf{Bool}$
  valid(sy_sti,sy_rul)(θ) $\equiv$ meaning(sy_rul)(θ,(meaning(sy_sti))(θ))

  valid: Stimulus $\times$ RUL $\to \Theta \to \mathbf{Bool}$
  valid(sy_sti,se_rul)(θ) $\equiv$ se_rul(θ,(meaning(sy_sti))(θ))

A syntactic regulation, sy_reg:Reg (related to a specific rule), stands for, i.e., has as its semantics, its meaning, a semantic regulation, se_reg:REG, which is a pair. This pair consists of a predicate, pre_reg:Pre_REG, where Pre_REG = $(\Theta \times \Theta) \to \mathbf{Bool}$, and a domain configuration-changing function, act_reg:Act_REG, where Act_REG = $\Theta \to \Theta$, that is, both involving current and next domain configurations. The two kinds of functions express: If the predicate holds, then the action can be applied.

The predicate is almost the inverse of the rules functions. The action function serves to undo the stimulus function.

type
  Reg
  Rul_and_Reg = Rule $\times$ Reg
  REG = Pre_REG $\times$ Act_REG

Pre_REG = $\Theta \times \Theta \to$ **Bool**
Act_REG = $\Theta \to \Theta$
value
interpret: Reg $\to$ REG

The idea is now the following: Any action of the system, i.e., the application of any stimulus, may be an action in accordance with the rules, or it may not. Rules therefore express whether stimuli are valid or not in the current configuration. And regulations therefore express whether they should be applied, and, if so, with what effort.

More specifically, there is usually, in any current system configuration, given a set of pairs of rules and regulations. Let (sy_rul,sy_reg) be any such pair. Let sy_sti be any possible stimulus. And let θ be the current configuration. Let the stimulus, sy_sti, applied in that configuration result in a next configuration, θ', where θ' = (meaning(sy_sti))(θ). Let θ' violate the rule, $\sim$valid(sy_sti,sy_rul)(θ), then if predicate part, pre_reg, of the meaning of the regulation, sy_reg, holds in that violating next configuration, pre_reg(θ,(meaning(sy_sti))(θ)), then the action part, act_reg, of the meaning of the regulation, sy_reg, must be applied, act_reg(θ), to remedy the situation.

axiom
$\forall$ (sy_rul,sy_reg):Rul_and_Regs •
  let se_rul = meaning(sy_rul),
      (pre_reg,act_reg) = meaning(sy_reg) **in**
  $\forall$ sy_sti:Stimulus, θ:Θ •
    $\sim$valid(sy_sti,se_rul)(θ)
      $\Rightarrow$ pre_reg(θ,(meaning(sy_sti))(θ))
        $\Rightarrow$ $\exists$ nθ:Θ • act_reg(θ)=nθ $\wedge$ se_rul(θ,nθ)
  end

It may be that the regulation predicate fails to detect applicability of regulations actions. That is, the interpretation of a rule differs, in that respect, from the interpretation of a regulation. Such is life in the domain, i.e., in actual reality

On Modelling Rules and Regulations

Usually rules (as well as regulations) are expressed in terms of domain entities, including those grouped into "the state", functions, eventws, and behaviours. Thus the full spectrum of modelling techniques and notations may be needed. Since rules usually express properties one often uses some combination of axioms and wellformedness predicates. Properties sometimes include temporality and hence temporal notations (like Duration Calculus [77, 78] or Temporal Logic of Actions [48, 49, 51]) are used. And since regulations usually express state (restoration) changes one often uses state changing notations (such as found in B [1], RSL [26], VDM-SL [10, 11, 23], and Z [68, 69, 75, 38]). In some cases it may be relevant to model using some constraint satisfaction notation [2] or some Fuzzy Logic notations [61].

4.5 Scripts and Licensing Languages

- *By a domain* **script** *we shall understand the structured, almost, if not outright, formally expressed, wording of a rule or a regulation that has legally binding power, that is, which may be contested in a court of law.*

Example 19. **A Casually Described Bank Script:** We deviate, momentarily, from our line of railway examples, to exemplify one from banking. Our formulation amounts to just a (casual) rough sketch. It is followed by a series of four large examples. Each of these elaborate on the theme of (bank) scripts.

The problem area is that of how repayments of mortgage loans are to be calculated. At any one time a mortgage loan has a balance, a most recent previous date of repayment, an interest rate and a handling fee. When a repayment occurs, then the following calculations shall take place: (i) the interest on the balance of the loan since the most recent repayment, (ii) the handling fee, normally considered fixed, (iii) the effective repayment — being the difference between the repayment and the sum of the interest and the handling fee — and the new balance, being the difference between the old balance and the effective repayment.

We assume repayments to occur from a designated account, say a demand/deposit account. We assume that bank to have designated fee and interest income accounts.

(i) The interest is subtracted from the mortgage holder's demand/deposit account and added to the bank's interest (income) account. (ii) The handling fee is subtracted from the mortgage holder's demand/deposit account and added to the bank's fee (income) account. (iii) The effective repayment is subtracted from the mortgage holder's demand/deposit account and also from the mortgage balance. Finally, one must also describe deviations such as overdue repayments, too large, or too small repayments, and so on. •

Example 20. **A Formally Described Bank Script:**

First we must informally and formally define the bank state:

There are clients (c:C), account numbers (a:A), mortgage numbers (m:M), account yields (ay:AY) and mortgage interest rates (mi:MI). The bank registers, by client, all accounts (ρ:A_Register) and all mortgages (μ:M_Register). To each account number there is a balance (α:Accounts). To each mortgage number there is a loan (ℓ:Loans). To each loan is attached the last date that interest was paid on the loan.

type
  C, A, M
  AY′ = **Real**, AY = {| ay:AY′ • 0<ay≤10 |}

$\mathrm{MI}' = \mathbf{Real}$, $\mathrm{MI} = \{|\ \mathrm{mi}:\mathrm{MI}' \bullet 0<\mathrm{mi}\leq 10\ |\}$
$\mathrm{Bank}' = \mathrm{A\_Register} \times \mathrm{Accounts} \times \mathrm{M\_Register} \times \mathrm{Loans}$
$\mathrm{Bank} = \{|\ \beta:\mathrm{Bank}' \bullet \mathrm{wf\_Bank}(\beta)|\}$
$\mathrm{A\_Register} = \mathrm{C} \xrightarrow[m]{} \mathrm{A}\textbf{-set}$
$\mathrm{Accounts} = \mathrm{A} \xrightarrow[m]{} \mathrm{Balance}$
$\mathrm{M\_Register} = \mathrm{C} \xrightarrow[m]{} \mathrm{M}\textbf{-set}$
$\mathrm{Loans} = \mathrm{M} \xrightarrow[m]{} (\mathrm{Loan} \times \mathrm{Date})$
$\mathrm{Loan},\mathrm{Balance} = \mathrm{P}$
$\mathrm{P} = \mathbf{Nat}$

Then we must define well-formedness of the bank state:

value
ay:AY, mi:MI

wf_Bank: Bank $\rightarrow$ **Bool**
$\mathrm{wf\_Bank}(\rho,\alpha,\mu,\ell) \equiv \cup\ \mathbf{rng}\ \rho = \mathbf{dom}\ \alpha \wedge \cup\ \mathbf{rng}\ \mu = \mathbf{dom}\ \ell$
axiom
ai<mi

Operations on banks are denoted by the commands of the bank script language. First the syntax:

type
Cmd = OpA | CloA | Dep | Wdr | OpM | CloM | Pay
OpA == mkOA(c:C)
CloA == mkCA(c:C,a:A)
Dep == mkD(c:C,a:A,p:P)
Wdr == mkW(c:C,a:A,p:P)
OpM == mkOM(c:C,p:P)
Pay == mkPM(c:C,a:A,m:M,p:P)
CloM == mkCM(c:C,m:M,p:P)
Reply = A | M | P | OkNok
OkNok == ok | notok

And then the semantics:

$\mathrm{int\_Cmd}(\mathrm{mkPM}(\mathrm{c},\mathrm{a},\mathrm{m},\mathrm{p},\mathrm{d}'))(\rho,\alpha,\mu,\ell) \equiv$
  let $(\mathrm{b},\mathrm{d}) = \ell(\mathrm{m})$ **in**
  if $\alpha(\mathrm{a})\geq\mathrm{p}$
    then
      let $\mathrm{i} = \mathrm{interest}(\mathrm{mi},\mathrm{b},\mathrm{d}'-\mathrm{d})$,
        $\ell' = \ell \dagger [\,\mathrm{m}\mapsto\ell(\mathrm{m})-(\mathrm{p}-\mathrm{i})\,]$
        $\alpha' = \alpha \dagger [\,\mathrm{a}\mapsto\alpha(\mathrm{a})-\mathrm{p},a_i\mapsto\alpha(a_i)+\mathrm{i}\,]$ **in**
      $((\rho,\alpha',\mu,\ell'),\mathrm{ok})$ **end**
    else

$((\rho,\alpha',\mu,\ell),\text{nok})$
end end
pre $\text{c} \in \textbf{dom}\ \mu \wedge \text{m} \in \mu(\text{c})$

•

The idea about scripts is that they can somehow be objectively enforced: that they can be precisely understood and consistently carried out by all stakeholders, eventually leading to computerisation. But they are, at all times, part of the domain.

Licensing Languages

A special form of scripts are increasingly appearing in some domains, notably the domain of electronic, or digital media, where these licenses express that the licensor permits the licensee to render (i.e., play) works of proprietary nature CD ROM-lile music, DVD-like movies, ec. while obligating the licensee to pay the licensor on behalf of the owners of these, usually artistic works. We refer to [30, 55, 64, 12, 4, 17, 76] for papers and reports on license languages.

On Modelling Scripts

Scripts (as are licenses) are like programs (respectively like prescriptions for program executions). Hence the full variety of techniques and notations for modelling programming (or specification) languages apply [20, 31, 60, 65, 73, 74]. Chapters 6–9 of Vol. 2 of [6, 7, 8] cover pragmatics, semantics and syntax techniques for defining languages.

4.6 Human Behaviour

- *By domain* **human behaviour** *we shall understand any of a quality spectrum of carrying out assigned work: from (i) careful, diligent and accurate, via (ii) sloppy dispatch, and (iii) delinquent work, to (iv) outright criminal pursuit.*

Example 21. **Banking — or Programming — Staff Behaviour:** Let us assume a bank clerk, “in ye olde” days, when calculating, say mortgage repayments (cf. Example 19).

We would characterise such a clerk as being *diligent*, diligent human behaviourhuman behaviour!diligentetc., if that person carefully follows the mortgage calculation rules, and checks and double-checks that calculations “tally up”, or lets others do so. We would characterise a clerk as being *sloppy* sloppy human behaviourhuman behaviour!sloppyif that person occasionally forgets the checks alluded to above. We would characterise a clerk as being *delinquent* delinquent human behaviourhuman behaviour!delinquentif that person systematically forgets these checks. And we would call such a person a *criminal*

criminal human behaviourhuman behaviour!criminalif that person intentionally miscalculates in such a way that the bank (and/or the mortgage client) is cheated out of funds which, instead, may be diverted to the cheater.

Let us, instead of a bank clerk, assume a software programmer charged with implementing an automatic routine for effecting mortgage repayments (cf. Example 20).

We would characterise the programmer as being *diligent*diligent human behaviour human behaviour!diligentif that person carefully follows the mortgage calculation rules, and throughout the development verifies and tests that the calculations are correct with respect to the rules. We would characterise the programmer as being *sloppy*sloppy human behaviour human behaviour!sloppyif that person forgets certain checks and tests when otherwise correcting the computing program under development. We would characterise the programmer as being *delinquent*delinquent human behaviour human behaviour!delinquentif that person systematically forgets these checks and tests. And we would characterise the programmer as being a *criminal*criminal human behaviour human behaviour!criminalif that person intentionally provides a program which miscalculates the mortgage interest, etc., in such a way that the bank (and/or the mortgage client) is cheated out of funds. •

Example 22. **A Human Behaviour Mortgage Calculation:**

Example 20 gave a semantics to the mortgage calculation request (i.e., command) as would a diligent bank clerk be expected to perform it. To express, that is, to model, how sloppy, delinquent, or outright criminal persons (staff?) could behave we must modify the int_Cmd(mkPM(c,a,m,p,d′))(ρ,α,μ,ℓ) definition.

```
int_Cmd(mkPM(c,a,m,p,d′))(ρ,α,μ,ℓ) ≡
    let (b,d) = ℓ(m) in
    if q(α(a),p) /∗ α(a)≤p∨α(a)=p∨α(a)≤p∨... ∗/
       then
          let i = f₁(interest(mi,b,d′−d)),
              ℓ′ = ℓ † [ m↦f₂(ℓ(m)−(p−i)) ]
              α′ = α † [ a↦f₃(α(a)−p),aᵢ↦f₄(α(aᵢ)+i),
                            a“staff”↦f“staff”(α(aᵢ)+i) ] in
          ((ρ,α′,μ,ℓ′),ok) end
       else
          ((ρ,α′,μ,ℓ),nok)
    end end
    pre c ∈ dom μ ∧ m ∈ μ(c)

q: P × P ⥲ Bool
f₁,f₂,f₃,f₄,f“staff”: P ⥲ P
```

•

The predicate q and the functions f_1, f_2, f_3, f_4 and $f_{\text{“staff”}}$ of Example 22 are deliberately left undefined. They are being defined by the “staffer” when performing (incl., programming) the mortgage calculation routine.

The point of Example 22 is that one must first define the mortgage calculation script precisely as one would like to see the diligent staff (prgrammer) to perform (incl., correctly program) it before one can “pinpoint” all the places where lack of diligence may “set in”. The invocations of q, f_1, f_2, f_3, f_4 and $f_{\text{“staff”}}$ designate those places.

The point of Example 22 is also that we must first domain-define, “to the best of our ability” all the places where human behaviour may play other than a desirable role. If we cannot, then we cannot claim that some requirements aim at countering undesriable human behaviour.

A Meta-characterisation of Human Behaviour

Commensurate with the above, humans interpret rules and regulations differently, and not always “consistently” — in the sense of repeatedly applying the same interpretations.

Our final specification pattern is therefore:

type
  Action = $\Theta \xrightarrow{\sim} \Theta$-**infset**
value
  hum_int: Rule $\to \Theta \to$ RUL-**infset**
  action: Stimulus $\to \Theta \to \Theta$
  hum_beha: Stimulus $\times$ Rules $\to$ Action $\to \Theta \xrightarrow{\sim} \Theta$-**infset**
  hum_beha(sy_sti,sy_rul)(α)(θ) **as** θset
    post
      θset = $\alpha(\theta)$ $\wedge$ action(sy_sti)(θ) $\in$ θset
      $\wedge$ $\forall$ θ':Θ•θ' $\in$ θset $\Rightarrow$
        $\exists$ se_rul:RUL•se_rul $\in$ hum_int(sy_rul)(θ)$\Rightarrow$se_rul(θ,θ')

The above is, necessarily, sketchy: There is a possibly infinite variety of ways of interpreting some rules. A human, in carrying out an action, interprets applicable rules and chooses one which that person believes suits some (professional, sloppy, delinquent or criminal) intent. “Suits” means that it satisfies the intent, i.e., yields **true** on the pre/post-configuration pair, when the action is performed — whether as intended by the ones who issued the rules and regulations or not. We do not cover the case of whether an appropriate regulation is applied or not

The above-stated axioms express how it is in the domain, not how we would like it to be. For that we have to establish requirements.

On Modelling Human Behaviour

To model human behaviour is, "initially", much like modelling management and organsiation. But only 'initially'. The most significant human behaviour modelling aspects is then that of modelling non-determinism and looseness, even ambiguity. So a specification language which allows specifying non-determinism and looseness (like CafeOBJ [25, 24, 21] and RSL [26]) is to be preferred.

4.7 Completion

Domain acquisition resulted in typically up to thousands of units of domain descriptions. Domain analysis subsequently also serves to classify which facet any one of these description units primarily characterise. But some such "compartmentalisations" may be difficult, and may be deferred till the step of "completion". It may then be, "at the end of the day", that is, after all of the above facets have been modelled that some description units are left as not having been described, not deliberately, but "circumstantially". It then behooves the domain engineer to fit these "dangling" description units into suitable parts of the domain description. This "slotting in" may be simple, and all is fine. Or it may be difficult. Such difficulty may be a sign that the chosen model, the chosen description, in its selection of entities, functions, events and behaviours to model — in choosing these over other possible selections of phenomena and concepts is not appropriate. Another attempt must be made. Another selection, another abstraction of entities, functions, etc., may need be chosen. Usually however, after having chosen the abstractions of the intrinsic phenomena and concepts, one can start checking whether "dangling" description units can be fitted in "with ease".

4.8 Integrating Formal Descriptions

We have seen that to model the full spectrum of domain facets one needs not one, but several specification languages. No single specification language suffices. It seems highly unlikely and it appears not to be desirable to obtain a single, "universal" specification language capable of "equally" elegantly, suitably abstractly modelling all aspects of a domain. Hence one must conclude that the full modelling of domains shall deploy several formal notations. The issues are then the following which combinations of notations to select, and how to make sure that the combined specification denotes something meaningful. The ongoing series of "Integrating Formal Methods" conferences [3, 29, 15, 13, 62] is a good source for technques, compositions and meaning.

5 From Domain Models to Requirements

One rôle for $\mathcal{D}$omain descriptions is to serve as a basis for constructing , $\mathcal{R}$equirements prescriptions. The purpose of constructing $\mathcal{R}$equirements prescriptions is to specify properties (not implementation) of a $\mathcal{M}$achine. The $\mathcal{M}$achine is the hardware (equiment) and the software that together implements the requirements. The implementation relations is

$$\mathcal{D}, \mathcal{M} \models \mathcal{R}$$

The $\mathcal{M}$achine is proven to implement the $\mathcal{R}$equirements in the context of [assumptions about] the $\mathcal{D}$omain. That is, proofs of correctness of the $\mathcal{M}$achine wrt. the $\mathcal{R}$equirements often refer to properties of the $\mathcal{D}$omain.

5.1 Domain Requirements

First, in a concrete sense, you copy the domain descriptiona and call it a requirements prescriptions. Then that requirements prescription is subjected to a number of operations: (i) removal (projection away) of all those aspects not needed in the requiriments; (ii) instantiation of remain aspects to the specifics of the client's domain; (iii) making determinate what is unnessarily or undesirably non-deterministic in the evolving requirements prescription; (iv) extending ot with concepts not feasible in the domain; and (v) fitting these requirements to those of related domains (say monitoring & control of public administration procedures). The result is called a domain requirements.

5.2 Interface Requirements

From the domain requirements one then constructs the interface requirements: First one identifies all phenomena and concepts, entities, functions, event and behaviours shared with the environment of the machine (hardware + software) being requirements specified. Then one requirements prescribe how each shared phenomenon and concept is being inistialised and updated: entity initialisation and refreshment, function initialisation and refreshment (interactive monitoring and control of computations), and the physiological man-machine and machine-machine implements.

5.3 Machine Requirements

Finally one deals with machine requirements performance, dependability, maintainability, portability, etc., where dependability addresses such issues as availability, accessability, reliability, safety, security, etc.

6 Why Domain Engineering?

6.1 Two Reasons for Domain Engineering

We believe that one can identify two almost diametrically opposed reasons for the pursuit of domain deswcriptions. One is utilitarian, concrete, commercial and engineering goal-oriented. It claims that domain engineering will lead to better software, and to development processes that can be better monitored and colled. and the other is science-oriented. It claims that establishing domain theories is a necessity, that it must be done, whither we develop software or not.

We basically take the latter, the science, view, while, of course, noting the former, the engineering consequences. We will briefly look at these.

6.2 An Engineering Reason for Domain Modelling

In a recent e-mail, in response, undoubtedly to my steadfast, perhaps conceived as stubborn insistance, on domain engineering, Tony Hoare summed up his reaction to domain engineering as follows, and I quote[7]:

"There are many unique contributions that can be made by domain modelling.

1. The models describe all aspects of the real world that are relevant for any good software design in the area. They describe possible places to define the system boundary for any particular project.
2. They make explicit the preconditions about the real world that have to be made in any embedded software design, especially one that is going to be formally proved.
3. They describe the whole range of possible designs for the software, and the whole range of technologies available for its realisation.
4. They provide a framework for a full analysis of requirements, which is wholly independent of the technology of implementation.
5. They enumerate and analyse the decisions that must be taken earlier or later any design project, and identify those that are independent and those that conflict. Late discovery of feature interactions can be avoided."

All of these issues are dealt with, one-by-one, and in depth, in Vol. 3 of my three volume book [8].

6.3 On a Science of Domains

Domain Theories

Although not brought out in this chapter the concept of domain theories must now be mentioned.

- *By a* **domain theory** *we shall understand a domain description together with lemmas, propositions and theorems that can be proved about the description — and hence can be claimed to hold in the domain.*

[7] E-Mail to Dines Bjørner, CC to Robin Milner et al., July 19, 2006

To create a domain theory the specification language must possess a proof system. It appears that the essence of possible theorems of — that is, laws about — domains can be found in laws of physics. For a delightful view of the law-based nature of physics — and hence possibly also of man-made universes we refer to Richard Feynmann's Lectures on Physics [22].

A Scientific Reason for Domain Modelling

So, inasmuch as the above-listed issues of Sect. 6.2, so aptly expressed in Tony's mastery, also of concepts (through his delightful mastery of words), are of course of utmost engineering importance, it is really, in our mind, the science issues that are foremost: We must first and foremost understand. There is no excuse for not trying to first understand. Whether that understanding can be "translated" into engineering tools and techniques is then another matter. But then, of course, it is nice that clear and elegant understanding also leads to better tools and hence better engineering. It usually does.

7 Conclusion

7.1 Summary

We have introduced the scientific and engineering concept of domain theories and domain engineering; and we have brought but a mere sample of the principles, techniques and tools that can be used in creating domain descriptions.

7.2 Grand Challenges of Informatics

To establish a reasonably trustworthy and believable theory of a domain, say the transportation, or just the railway domain, may take years, possibly 10–15! Similarly for domains such as the financial service industry, the market (of consumers and producers, retailers, wholesaler, distribution cum supply chain), healthcare, and so forth.

The current author urges younger scientists to get going! It is about time.

7.3 Acknowledgements

The author (and editors) thank Springer for allowing this chapter — which is basically a rewriting and "condensing-editing" of Chaps. 5 and 11 of [8] to appear also in the present book. The author thanks Paul Boca and Jonathan P. Bowen for inviting him in the first place to give basically this chapter as a BCS FACS seminar in the summer of 2005.[8] The author thanks Tony Hoare

[8] At that time the book [6, 7, 8] was in the hands of the editors of Springer, i.e., not yet published.

for permission to quote his characterisation of domain engineering (Sect. 6.2); and the author thanks Prof. Kokichi Futatsugi, of JAIST, for inviting him to JAIST for a full year, 2006, to further study domain engineering and its implications, including arranging for time to compose this chapter. The author finally thanks his home institute (IMM/DTU) for giving him a year's sabbatical.

8 Bibliographical Notes

To create domain descriptions, or requirements prescriptions, or software designs, properly, at least such as this author sees it, is a joy to behold. The beauty of carefully selected and balanced abstractions, their interplay with other such, the relations between phases, stages and steps, and many more conceptual constructions make software engineering possibly the most challenging intellectual pursuit today. For this and more consult [6, 7, 8].

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