

I217: Functional Programming

1. Sorts, Operators, Terms & Equations

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Roadmap

- Preliminaries
 - Programming paradigms
 - Functional programming & CafeOBJ
 - Sets, Sequences, Functions
- Some examples
- Sorts, Operators, Variables, Terms, Equations
- Some commands and comments
- Exercises

Programming Paradigms

- Imperative (procedural) programming
 - Pascal, C, C++, Standard ML, Oz, Ruby, Python
- Logic programming
 - Prolog, Oz, GHC, KL1
- Object-oriented programming
 - Smalltalk, C++, Java, Self, Oz, Scala, Ruby, ABCL/1, ConcurrentSmalltalk, MultithreadSmalltalk, Python
- *Functional programming*
 - Miranda, Haskell, Standard ML, Oz, Scala, Maude, CafeOBJ

Functional Programming & CafeOBJ

- Functional Programming
 - Programs are functions in the *mathematical* sense.
 - To execute programs is to evaluate expressions (i.e. function applications).
 - *No destructive assignment.*
 - More amenable to *program verification.*
- CafeOBJ
 - An executable *specification* language.
 - Not a real programming language, but can be used as an educational programming language.
 - CafeOBJ programs (specifications) can be verified, namely that we can prove that programs written in CafeOBJ enjoy (desired) properties with the CafeOBJ system.

Some Example in CafeOBJ

```

open NAT .
op gcd : Nat Nat -> Nat .
var X : Nat .
var NzY : NzNat .
eq gcd(X,0) = X .
eq gcd(X,NzY) = gcd(NzY,X rem NzY) .
--
red gcd(0,0) . -- compute the gcd of 0 & 0
red gcd(2,0) . -- compute the gcd of 2 & 0
red gcd(0,2) . -- compute the gcd of 0 & 2
red gcd(24,36) . -- compute the gcd of 24 & 36
red gcd(2015,31031) . -- compute the gcd of 2015 & 31031
close

```

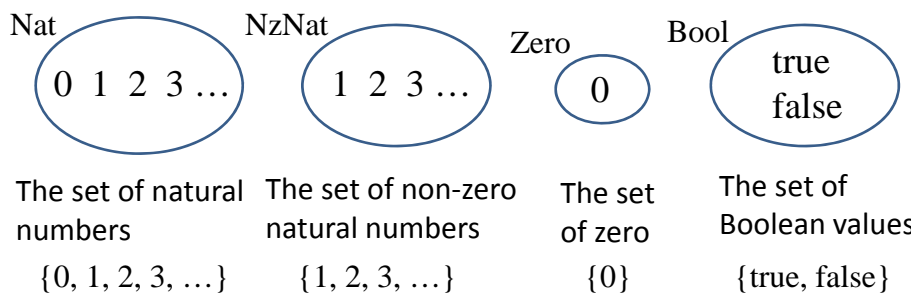
```

op _rem_ : Nat NzNat -> Nat .
eq X rem NzY
  = remainder of dividing X by NzY .

```

Sets

A set is a collection of similar things such that (1) duplications are not allowed and (2) the enumeration order is irrelevant.



Because of (1) & (2), the following three sets are the same:

{0, 1, 2, 3} {3, 2, 0, 1} {1, 3, 2, 0, 2, 1}

Sequences

A sequence of length n (≥ 0) is a collection that consists of n elements such that the enumeration order is relevant and the elements are not necessarily similar.

$(110, \text{true}, 119)$ (117) (which is the same as 117)

$()$ (the empty sequence & may be omitted)

The set of the sequences of length n such that the i th ($0 \leq i \leq n$) elements are those of a set S_i is represented as

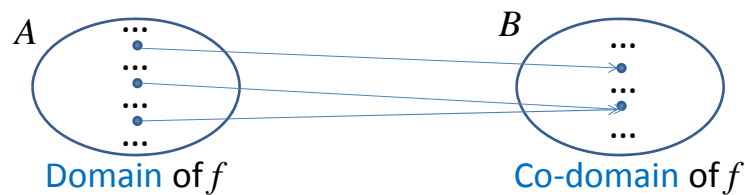
$S_0 \times \dots \times S_n$

The above three sequences are elements of the following sets respectively:

$\text{Nat} \times \text{Bool} \times \text{Nat}$ Nat $\{()\}$ (which may be omitted)

Functions

A function f from a set A to a set B maps each element of A to an exactly one element of B .



gcd is a function from $\text{Nat} \times \text{Nat}$ to Nat .

gcd maps $(0,0)$, $(2,0)$, $(0,2)$, $(24,36)$ and $(2015,31031)$ to 0, 2, 2, 12 and 403, respectively.

Constants, such as 0, can be regarded as functions from $\{()\}$ to a set. 0 is a function from $\{()\}$ to Nat .

Factorial

```

open NAT .
op fact : Nat -> Nat .
var NzX : NzNat .
eq fact(0) = 1 .
eq fact(NzX) = NzX * fact(p NzX) .
--
red fact(0) .
red fact(1) .
red fact(10) .
red fact(100) .
red fact(1000) .
red fact(10000) .
-- red fact(100000) . -- stack overflow
close

```

```

op p_ : NzNat -> Nat .
eq p NzX = the previous number of NzX .

```

```

op *_ : Nat Nat -> Nat .
vars X Y : Nat .
eq X * Y = multiplication of X & Y .

```

Odd-Even Divide & Conquer Factorial

```

open NAT .
op cond : Bool Nat Nat -> Nat .
op g : Nat Nat -> Nat .
op oedc-fact : Nat -> Nat .
vars X Y : Nat . var NzX : NzNat .
-- cond
eq cond(true,X,Y) = X .
eq cond(false,X,Y) = Y .
-- g
eq g(X,Y) = cond(X > Y, g(X,2 * Y) * g(sd(X,Y),2 * Y), X) .
-- oedc-fact
eq oedc-fact(0) = 1 .
eq oedc-fact(NzX) = g(NzX,1) .
--
red oedc-fact(10000) . -- can be computed
close

```

```

op _>_ : Nat Nat -> Nat .
eq X > Y = true if so
           false otherwise .

```

```

op sd : Nat Nat -> Nat .
eq sd(X,Y)
  = symmetric difference between X & Y .

```

Fibonacci

```

open NAT .
  op fib : Nat -> Nat .
  op sfib : Nat -> Nat .
  var NzX : NzNat .
  -- fib
  eq fib(0) = 0 .
  eq fib(NzX) = sfib(p NzX) .
  -- sfib
  eq sfib(0) = 1 .
  eq sfib(NzX) = fib(NzX) + fib(p NzX) .
  --
  red fib(10) .
  red fib(20) .
  red fib(30) . -- can be computed, although it takes time.
close

```

```

op _+_ : Nat Nat -> Nat .
vars X Y : Nat .
eq X + Y = addition of X & Y .

```

Ackerman

```

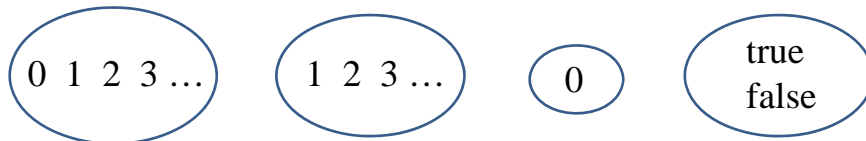
open NAT .
  op ack : Nat Nat -> Nat .
  var Y : Nat .
  vars NzX NzY : NzNat .
  eq ack(0,Y) = Y + 1 .
  eq ack(NzX,0) = ack(p NzX,1) .
  eq ack(NzX,NzY) = ack(p NzX,ack(NzX,p NzY)) .
  --
  red ack(0,0) .
  red ack(1,1) .
  red ack(2,2) .
  red ack(3,2) .
  red ack(3,3) .
close

```

Sorts

Interpreted as sets and correspond to types in programming languages.

Nat, **NzNat**, **Zero** and **Bool** are sorts that are interpreted as the following sets respectively:



May be used as the sets as which the sorts are interpreted.

Declared by enclosing them with [and], such as

[NatList]

Operators

Interpreted as functions and declared as follows:

$$\mathbf{op} f : S_1 \dots S_n \rightarrow S . \quad \text{where } n \geq 0$$

sorts
← ↓ →
arity
rank

co-arity

f is interpreted as a function from $S_1 \times \dots \times S_n$ to S .

When $n = 0$, f is called a constant of S .

$\mathbf{op} \text{ cond} : \text{Bool Nat Nat} \rightarrow \text{Nat} .$ are interpreted as functions from
 $\mathbf{op} \text{ g} : \text{Nat Nat} \rightarrow \text{Nat} .$ $\text{Bool} \times \text{Nat} \times \text{Nat}$ to Nat , $\text{Nat} \times \text{Nat}$ to Nat
 $\mathbf{op} \text{ oedc-fact} : \text{Nat} \rightarrow \text{Nat} .$ and Nat to Nat .

May be used as the functions as which they are interpreted.

Variables

Declared as follows:

var $V : S$.



Multiple variables of a same sort can be declared as follows:

vars $V_1 \dots V_n : S$.

vars $X Y : \text{Nat}$.

var $\text{NzX} : \text{NzNat}$.

X & Y are variables of Nat , and NzX is a variable of NzNat .

Terms

Constructed from operators and variables.

Have sorts.

$\text{gcd}(X,0)$	a term of sort Nat
$\text{gcd}(X,\text{NzY})$	a term of sort Nat
$\text{gcd}(\text{NzY},X \text{ rem } \text{NzY})$	a term of sort Nat
X	a term of sort Nat
0	a term of sort Zero
NzY	a term of sort NzNat
$X \text{ rem } \text{NzY}$	a term of sort Nat

Note that a term of sort Zero is also a term of sort Nat and a term of sort NzNat is also a term of sort Nat .

Terms

Inductively defined as follows:

- (1) A variable of sort S is a term of sort S .
- (2) For an operator $f : S_1 \dots S_n \rightarrow S$, if t_1, \dots, t_n are term of sorts S_1, \dots, S_n , then $f(t_1, \dots, t_n)$ is a term of sort S .

Note that when $n = 0$, f itself is a term of sort S (called a constant of sort S).

Note that terms of sort `Zero` or `NzNat` are also terms of sort `Nat`.

Terms

```

open NAT .
op gcd : Nat Nat -> Nat .
var X : Nat .
var NzY : NzNat .
...

```

`X` is a term of sort `Nat` because `X` is a variable of sort `Nat`.

`0` is a term of sort `Nat` because `0` is a term of sort `Zero`.

`gcd(X,0)` is a term of sort `Nat` because `gcd` is an operator whose rank is `Nat Nat -> Nat`, `X` is a term of sort `Nat` and `0` is a term of sort `Nat`.

Equations

eq $\text{gcd}(X,0) = X$.

says that for all natural numbers X , $\text{gcd}(X,0)$ equals X .

$$\text{gcd}(0,0) = 0$$

$$\text{gcd}(1,0) = 1$$

$$\text{gcd}(2,0) = 2$$

eq $\text{gcd}(X,\text{Nz}Y) = \text{gcd}(\text{Nz}Y, X \text{ rem } \text{Nz}Y)$.

says that for all natural numbers X & all non-zero natural numbers $\text{Nz}Y$, $\text{gcd}(X,\text{Nz}Y)$ equals $\text{gcd}(\text{Nz}Y, X \text{ rem } \text{Nz}Y)$.

$$\text{gcd}(3,1) = \text{gcd}(1, 3 \text{ rem } 1) .$$

$$\text{gcd}(31031,2015) = \text{gcd}(2015, 31031 \text{ rem } 2015) .$$

Equations

Declared as follows:

eq $\text{LeftTerm} = \text{RightTerm}$.

where LeftTerm and RightTerm are terms of a same sort.

If variables X_1, X_2, \dots of sorts S_1, S_2, \dots occur in the equation, then the equation says that for all X_1 of S_1 , all X_2 of S_2 , ... LeftTerm equals RightTerm .

Mix-fix Operators

The operators used in $3 + 4$, $p\ 1$ and $10\ !$ are called infix, prefix and postfix operators, which are declared as follows:

```

op _+_ : Nat Nat -> Nat .
op p_ : NzNat -> Nat .
op !_ : Nat -> Nat .

```

Moreover, the operator (called a *mix-fix operator*) used in `if X > Y then { g(X,2 * Y) * g(sd(X,Y),2 * Y) } else { X }` can be used and declared as follows:

```

op if_then { _ } else { _ } : Bool Nat Nat -> Nat .

```

Some Commands of CafeOBJ

Command `in`

Programs written in CafeOBJ are saved as text files whose names have the extension `.cafe`, such as `gcd.cafe` and `fact.cafe`. The command takes a file name and loads the programs in it.

```
CafeOBJ> in gcd.cafe
```

The extension can be omitted.

```
CafeOBJ> in gcd
```

Some Commands of CafeOBJ

Command `open` & `close`

`open` takes a module (a unit of programs in CafeOBJ) and makes it available.

```
CafeOBJ> open NAT .
```

where `NAT` is the built-in module in which natural numbers and functions over them are described and the prompt becomes `%NAT>` after opening `NAT`. Note that a period is needed.

`close` makes the currently open module close and the prompt back to `CafeOBJ>`.

```
%NAT> close
```

Some Commands of CafeOBJ

Command `red`

`red` takes a term and computes it with the equations available.

```
%NAT> red 3 + 4 .
```

Note that a period is needed.

Its result is `(7):NzNat`, where the sort `NzNat` is the important part of the result.

Some Commands of CafeOBJ

Command `full-reset`

It fully resets the system.

```
CafeOBJ> full-reset
```

Command `?`

It displays a list of commands available.

```
CafeOBJ> ?
```

Comments

A comment starts with the first occurrence `--` on a line and continues to the end of the line.

```
-- This is a comment.
```

A segment that starts with a double quotation mark and ends with a double quotation mark is also a comment.

```
"This is another way  
to write a comment."
```

When you want to use a double quotation mark in the segment, a backslash should be used in front of it.

```
"You can use \" double quotation  
marks\" in a comment."
```

Exercises

1. Type each piece of programs in the slides as one file, feed it into CafeOBJ, and do some testing. Note that the extension of a file name in which CafeOBJ programs are written is .cafe, such as fact.cafe.
2. Explain in which way fact(5) is computed.
3. Explain in which way oedc-fact(5) is computed.
4. Write two versions of programs computing the summation $0+1+2+\dots+n$ for a given number n , where one corresponds to fact and the other corresponds to oedc-fact and do some testing for both versions.

Exercises

5. Write a program in CafeOBJ that corresponds to the following and do some testing

$$\text{ext-fib}(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ 2 & \text{if } n = 2 \\ \text{ext-fib}(n-1) + \text{ext-fib}(n-2) & \text{otherwise} \\ \quad + \text{ext-fib}(n-3) & \end{cases}$$

The program should be such that sd is not used in it and no warning message on what are called error sorts is displayed when feeding it into the CafeOBJ system.

Exercises

6. Revise the program in which the factorial function is defined such that instead of `fact` the following postfix operator is used:

op `_!` : `Nat -> Nat` .

7. Revise the program in which the odd-even divide & conquer factorial function is defined such that instead of `oedc-fact` and `cond` the following postfix and mix-fix operators are used:

op `_!` : `Nat -> Nat` .

op `if_then {_} else {_}` : `Bool Nat Nat -> Nat` .