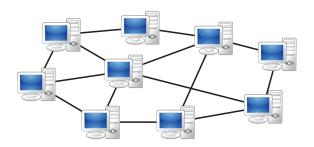
Lesson 10. Introduction to Mobile Robots 1628E – Information Processing Theory

Giovanni Viglietta
johnny@jaist.ac.jp

JAIST - January 22, 2020

Distributed systems: overview

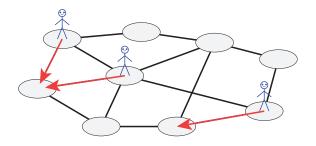
The traditional setting is a <u>network of computers</u>, each of which has its own local memory. Neighboring computers can communicate by sending messages.



A typical problem in this setting is electing a unique leader in a network of anonymous computers (i.e., where all computers are identical and execute the same algorithm).

Distributed systems: overview

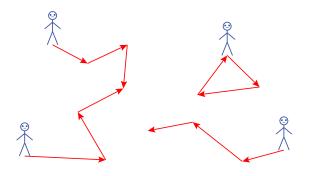
The standard model evolved into a more dynamic one: now the computers can <u>move</u> within the network (they are "robots") and write messages on shared whiteboards at each node.



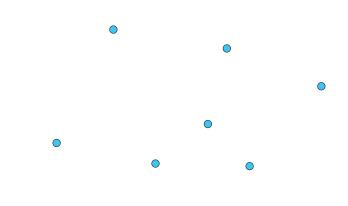
A typical problem in this setting is gathering in the same node. Again, robots are anonymous, i.e., they are all identical and execute the same algorithm.

Distributed systems: overview

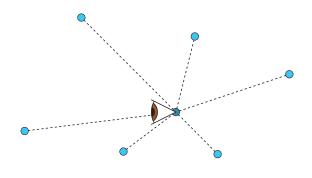
A further evolution of the model saw robots no longer constrained to a fixed network, but able to freely move in space.

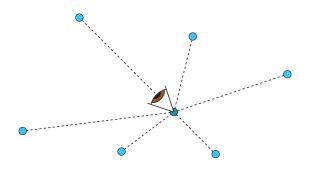


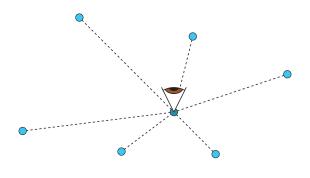
Each robot is now a computational unit that can sense the positions of other robots and compute its next destination.

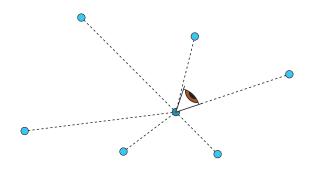


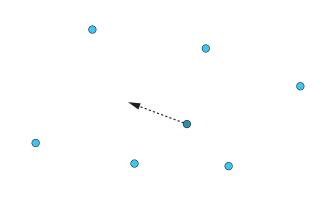
We consider a *swarm* of anonymous robots in the Euclidean plane.



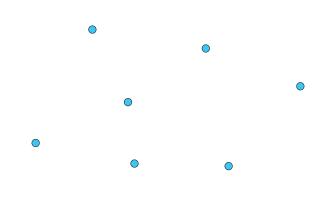




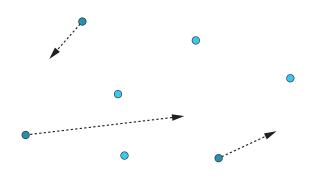


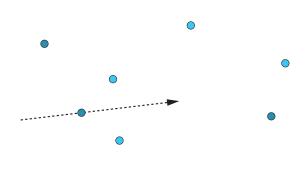


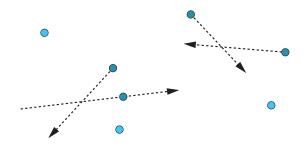
...And move according to a $\underline{\mathsf{deterministic}}$ algorithm.

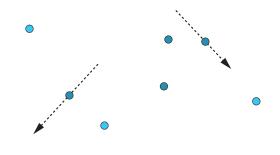


...And move according to a <u>deterministic</u> algorithm.





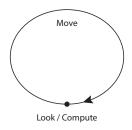




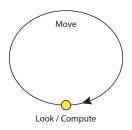
Model definition

Robots are:

- Dimensionless (robots are modeled as geometric points)
- Anonymous (no unique identifiers)
- Homogeneous (the same algorithm is executed by all robots)
- Deterministic (the algorithm has no randomization)
- Autonomous (no centralized control)
- Oblivious (no memory of past events and observations)
- Silent (no explicit way of communicating)
- Disoriented (robots do not share a common reference frame)
 - No common unit distance
 - No common compass
 - No common notion of clockwise direction

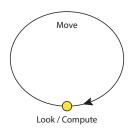


Each robot perpetually repeats a $\underline{\mathsf{Look}}/\mathsf{Compute}/\mathsf{Move}$ cycle.



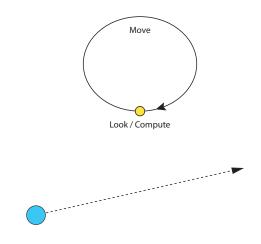


Each robot perpetually repeats a Look/Compute/Move cycle.

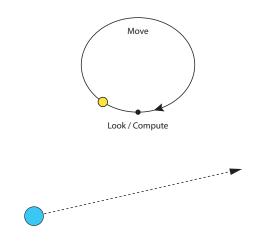


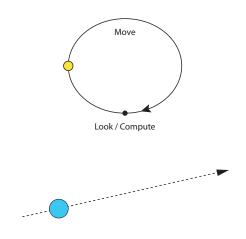


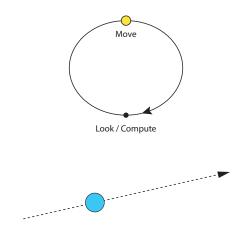
In a Look phase, a snapshot is taken of all visible robots.

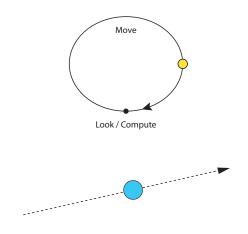


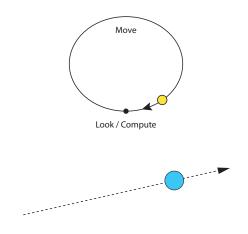
A <u>destination</u> is computed based only on the last snapshot.

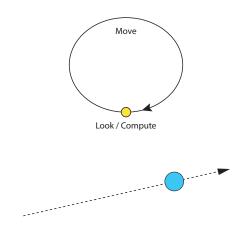




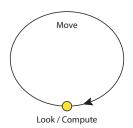






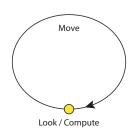


The robot may unpredictably stop before reaching the destination...



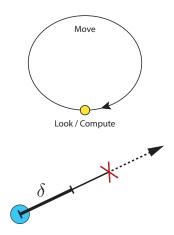


...and execute a new Look/Compute phase.

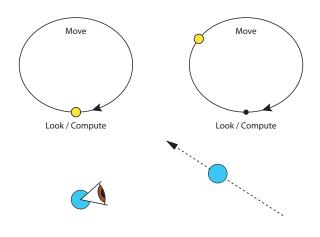




...and execute a new Look/Compute phase.

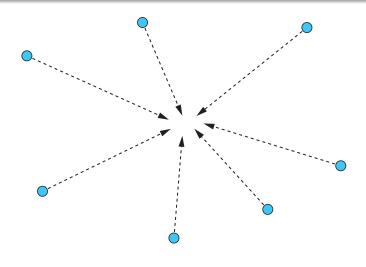


At each cycle, a robot is guaranteed to move by at least δ .



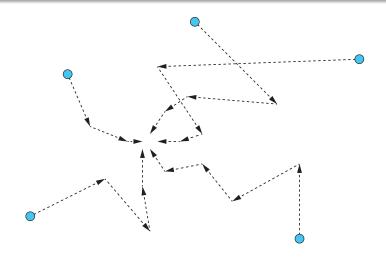
Different robots execute independent cycles, asynchronously.

Gathering-like problems



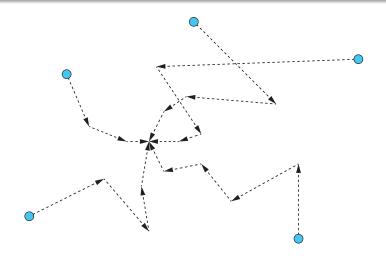
Perhaps the most studied class of problems: Design an algorithm that makes all robots "gather".

Convergence problem



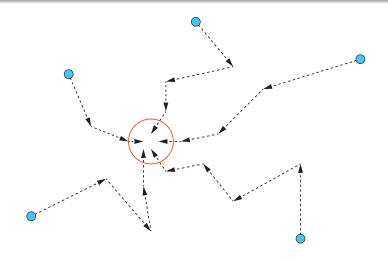
Convergence: Make all robots converge to the same point (possibly never actually reaching it, and possibly colliding).

Gathering problem



Gathering: Make all robots reach the same point in finite time.

Near-Gathering problem



Near-Gathering: Make all robots reach a small-enough area, avoiding collisions.

Case study: two robots

Suppose the swarm consists of only two robots.

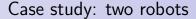
Case study: two robots

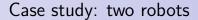
>

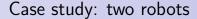
Strategy: Move to the midpoint.



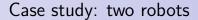
It solves the Gathering problem if robots are fully synchronous...



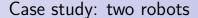












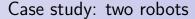








Modified strategy: Reduce the distance by a fraction of 1/3.



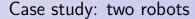
It solves Near-Gathering if robots are fully synchronous...

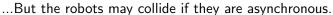


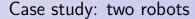
It solves Near-Gathering if robots are fully synchronous...



It solves Near-Gathering if robots are fully synchronous...





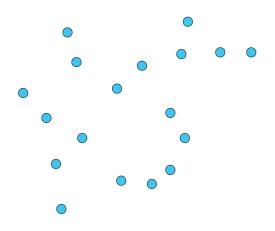


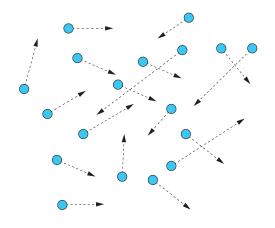


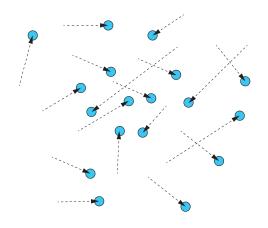


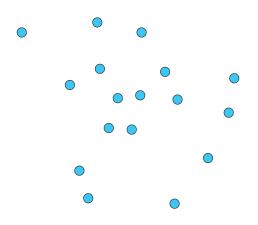


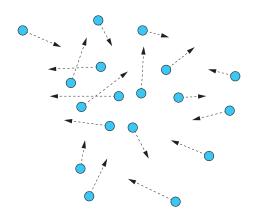




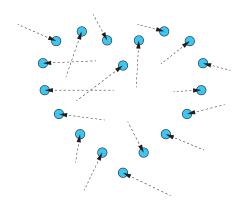


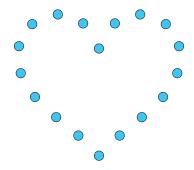


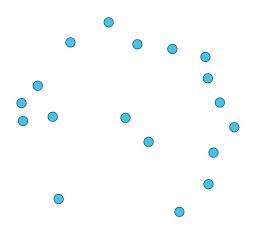


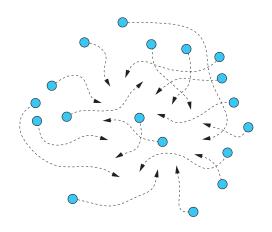


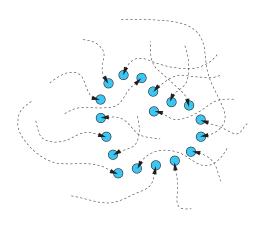
Pattern Formation: form a pattern (given as input) from <u>any</u> initial configuration

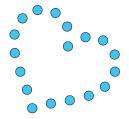


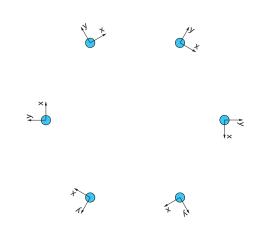




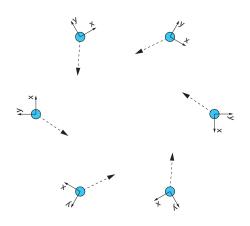




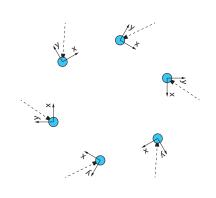




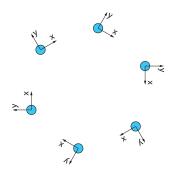
Let the initial configuration be rotationally symmetric



All robots have the same view and compute symmetric destinations



If they are all activated synchronously, they remain symmetric

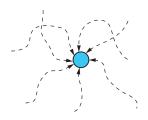


Hence Pattern Formation is unsolvable if the pattern is asymmetric

No pattern is formable from every possible initial configuration, except:

• Single point (aka Gathering problem)

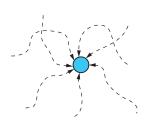
→ Solved (Cieliebak-Flocchini-Prencipe-Santoro, 2012)

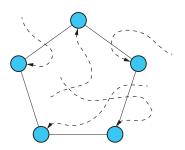


No pattern is formable from every possible initial configuration, except:

• Single point (aka Gathering problem)

⇒ Solved (Cieliebak-Flocchini-Prencipe-Santoro, 2012)

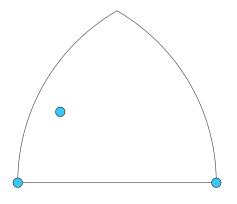




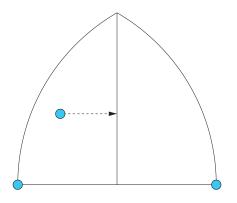
• Regular polygon (aka <u>Uniform Circle Formation</u> problem)

⇒ Today's lesson

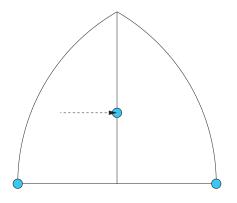




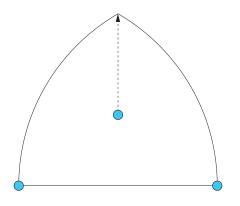
Identify the longest edge



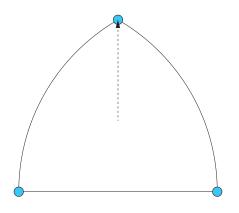
Move the third vertex parallel to the longest edge...



...until the triangle is isosceles

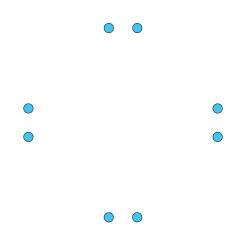


Then move the apex to the final position

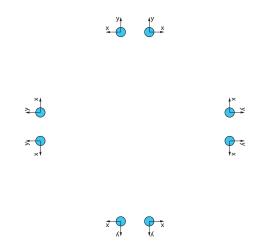


Then move the apex to the final position

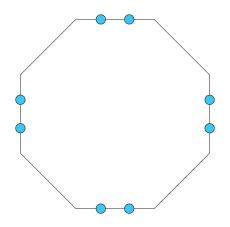
Correctness: only one robot is ever allowed to move



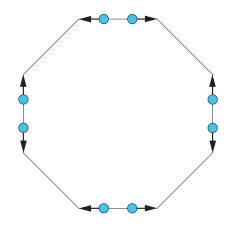
In a Biangular configuration all robots may have the same view



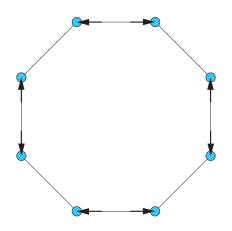
In a Biangular configuration all robots may have the same view



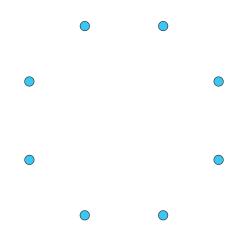
Solution: identify a "supporting polygon"



Each robot moves to the closest vertex



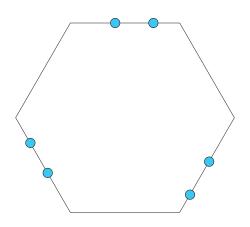
During the motion, the supporting polygon remains fixed (if $n>4\)$



During the motion, the supporting polygon remains fixed (if $n>4\)$

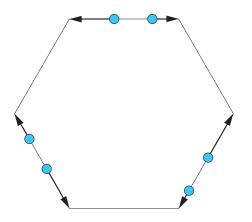
Pre-regular configurations

The possible (asynchronous) evolutions of a Biangular configuration are called **Pre-regular** configurations

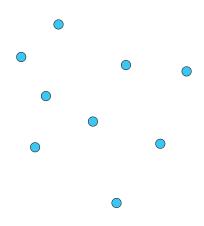


Pre-regular configurations

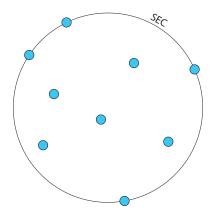
The possible (asynchronous) evolutions of a Biangular configuration are called **Pre-regular** configurations



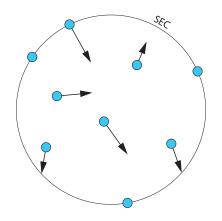
For consistency, they must be resolved in the same fashion



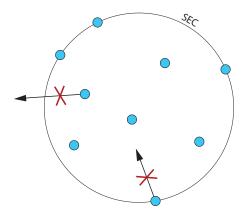
There always exists a unique SEC, which is easily computable



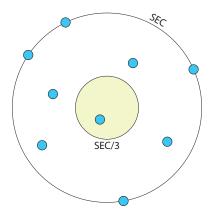
There always exists a unique SEC, which is easily computable



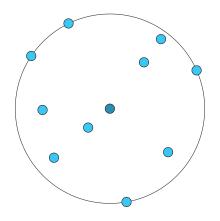
Strategy: always preserve the SEC, until a Pre-regular is formed



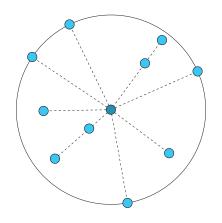
Strategy: always preserve the SEC, until a Pre-regular is formed



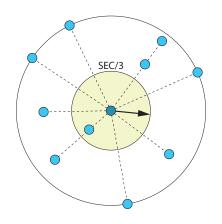
Another circle plays an important role: SEC/3



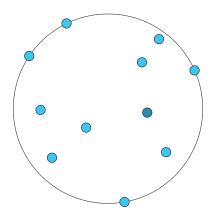
If there is a robot at the center of the SEC



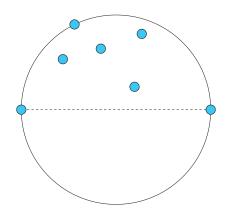
Identify the occupied radii



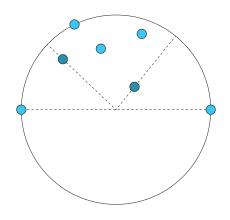
Move to an unoccupied radius all the way to $\ensuremath{\mathsf{SEC}}/3$



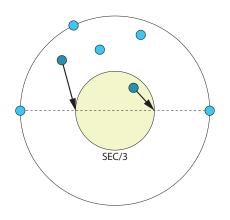
No Central configuration will ever be formed again



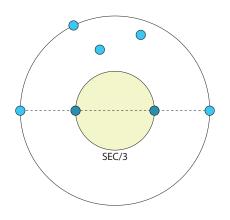
If all the robots lie in one half of the SEC



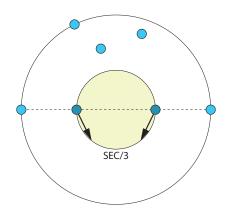
Identify the robots "angularly closest" to the diameter



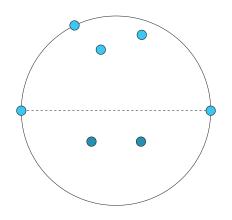
Move them to the diameter and make sure they are in $\ensuremath{\mathsf{SEC/3}}$



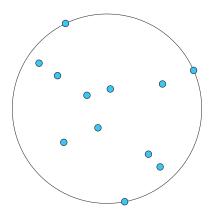
Move them to the diameter and make sure they are in $\ensuremath{\mathsf{SEC/3}}$



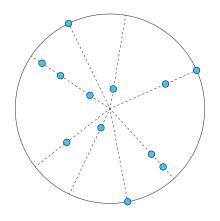
Make them cross the diameter but remain in SEC/3



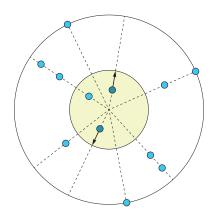
No half-disk configuration will ever be formed again



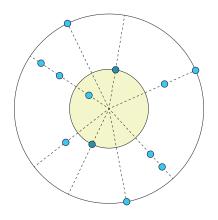
If there are co-radial robots



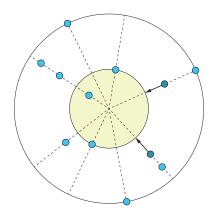
Identify the occupied radii



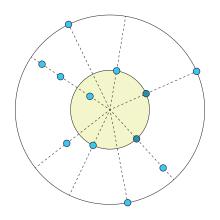
Move radially to SEC/3 the non-co-radial robots that lie inside it



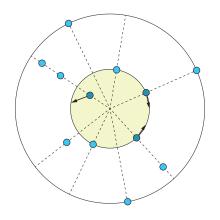
Move radially to SEC/3 the non-co-radial robots that lie inside it



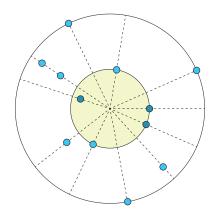
Move to SEC/3 the innermost co-radial robots that lie outside



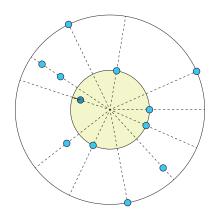
Move to SEC/3 the innermost co-radial robots that lie outside



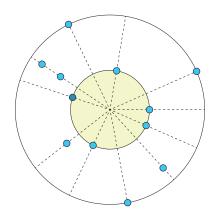
Move the innermost co-radial robots laterally by a small angle



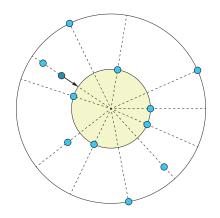
If moves are small enough, no new co-radialities are formed



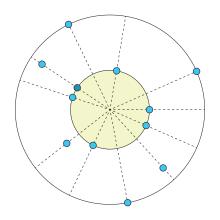
Move radially to SEC/3 the non-co-radial robots that lie inside it



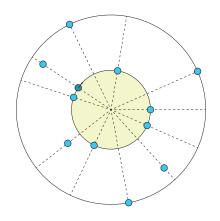
Move radially to SEC/3 the non-co-radial robots that lie inside it



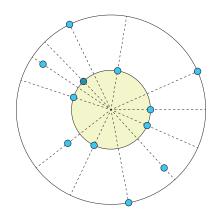
Move to SEC/3 the innermost co-radial robots that lie outside



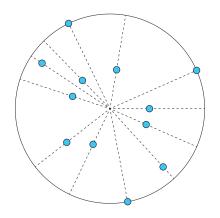
Move to SEC/3 the innermost co-radial robots that lie outside



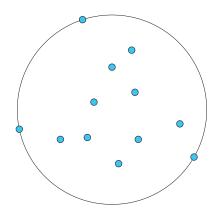
Move the innermost co-radial robots laterally by a small angle $\,$



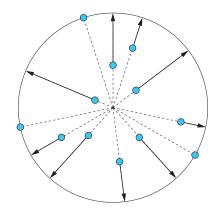
Move the innermost co-radial robots laterally by a small angle $\,$



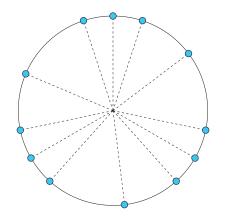
No Co-radial configuration will ever be formed again



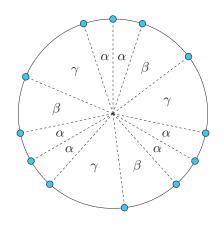
If we are not in one of the previous special cases



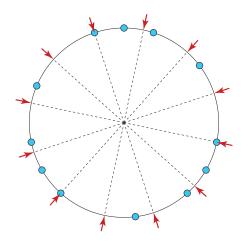
All robots move radially to SEC



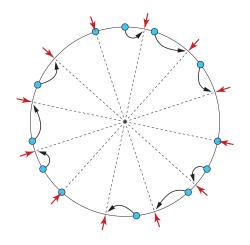
All robots move radially to SEC



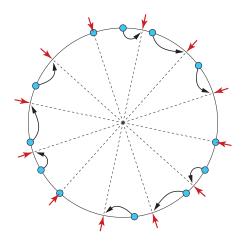
Consider the angle sequence



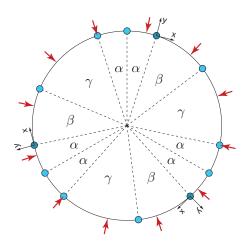
To solve Uniform Circle Formation, all angles must become equal



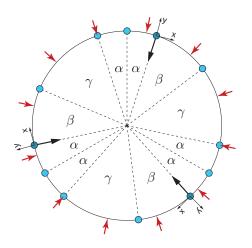
Identify a target for each robot



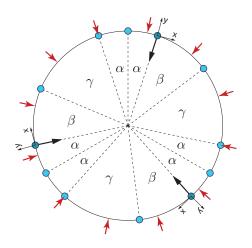
If all robots move together, they may forget their targets!



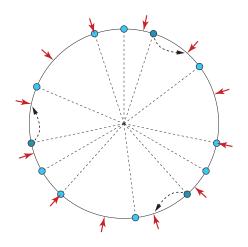
Robots that "see" the same angle sequence are analogous



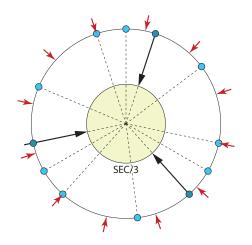
The scheduler can force analogous robots to move together



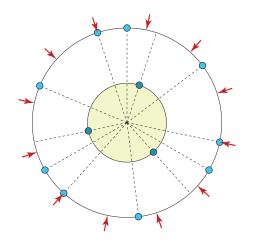
Hence the algorithm will move one analogy class at a time



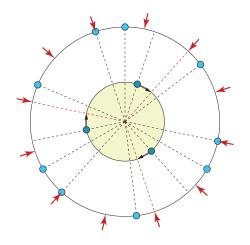
Strategy: choose analogous robots that can "see" their targets



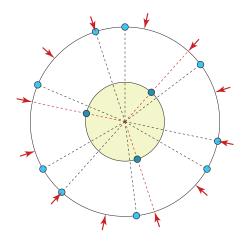
Move these walkers radially to $\ensuremath{\mathsf{SEC}}/3$



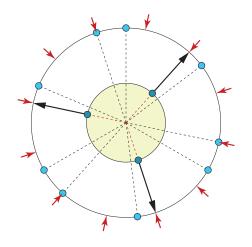
Move these walkers radially to SEC/3



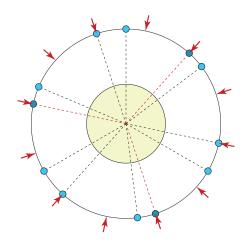
Move them laterally to their targets



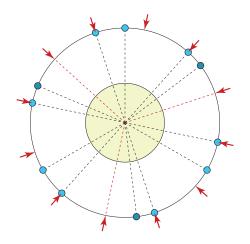
Move them laterally to their targets

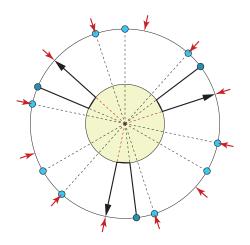


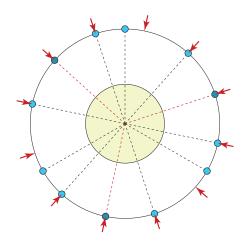
Move them back to SEC

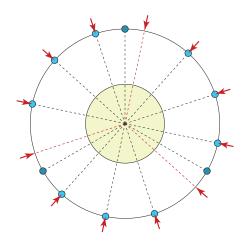


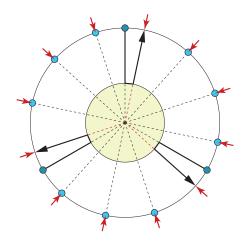
Move them back to SEC

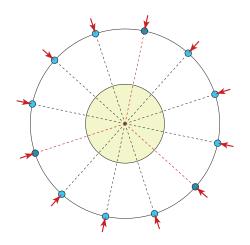


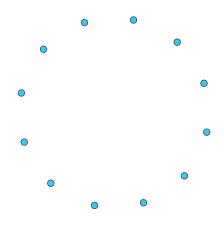




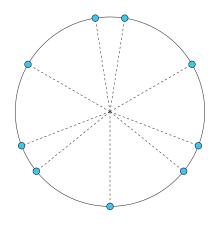






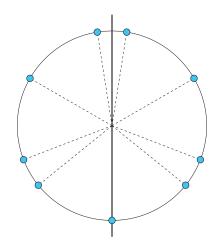


How to identify the target set



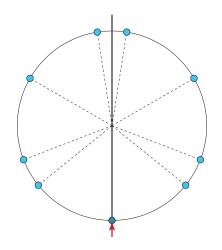
Suppose the configuration has an axis of symmetry $% \left(x\right) =\left(x\right) +\left(x\right)$

How to identify the target set

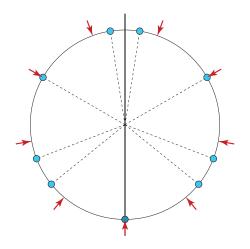


Suppose the configuration has an axis of symmetry $% \left(x\right) =\left(x\right) +\left(x\right)$

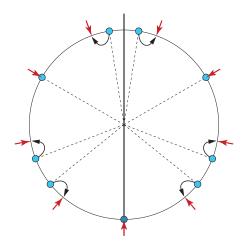
How to identify the target set



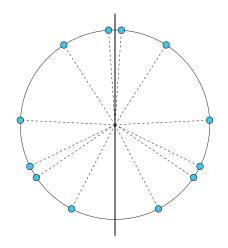
If a robot lies on the axis, it is on its target by definition



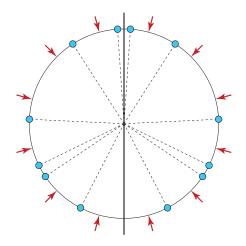
The other targets are determined accordingly



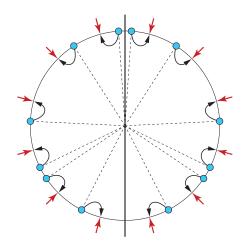
The other targets are determined accordingly



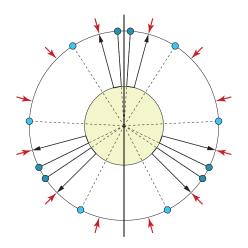
Suppose that no robot lies on the axis of symmetry



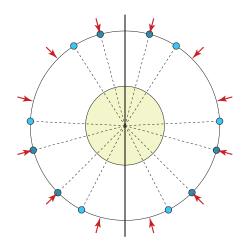
Then no target lies on the axis of symmetry, either



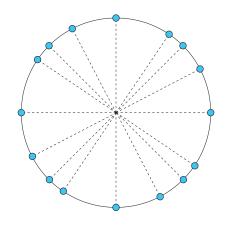
Then no target lies on the axis of symmetry, either



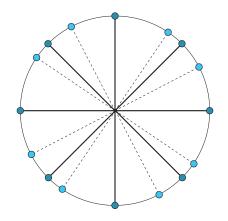
When an analogy class moves, the axis of symmetry is preserved



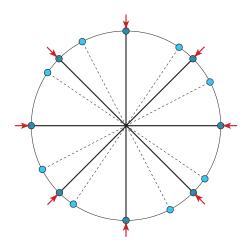
Hence also the targets are preserved $% \left(t_{0}\right) =\left(t_{0}\right) +\left(t_{0}\right) =\left(t_{0}\right) +\left(t_{0}\right) +\left$



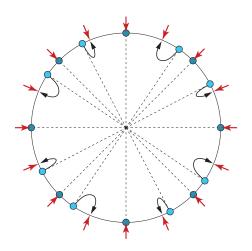
Suppose the configuration has no axis of symmetry $% \left(x\right) =\left(x\right) +\left(x\right)$



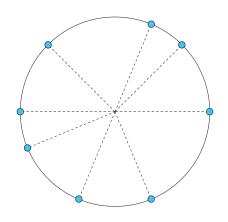
Robots having the "correct" angular distance are $\boldsymbol{concordant}$



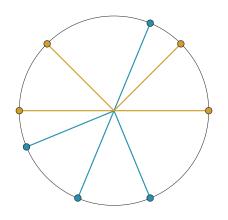
The targets are determined by the largest concordance class



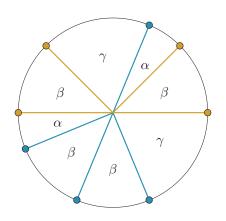
The targets are determined by the largest concordance class



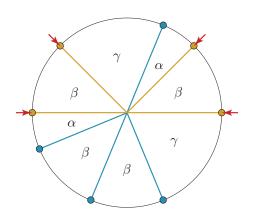
If there is more than one largest concordance class...



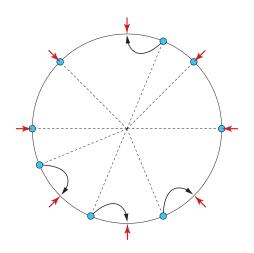
If there is more than one largest concordance class...



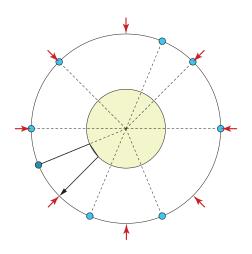
 $\ldots they \ are \ all \ "non-equivalent", so one can always be chosen$



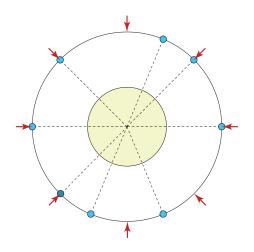
 \ldots they are all "non-equivalent", so one can always be chosen



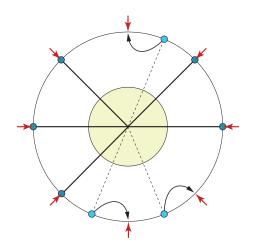
 \ldots they are all "non-equivalent", so one can always be chosen



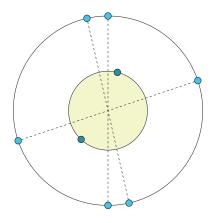
After a move, the chosen concordance class becomes the largest



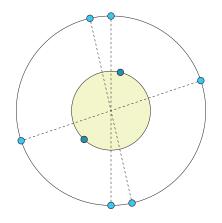
After a move, the chosen concordance class becomes the largest



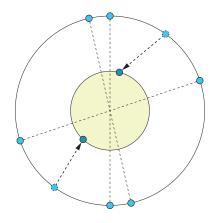
So the targets are preserved, until an axis of symmetry is created



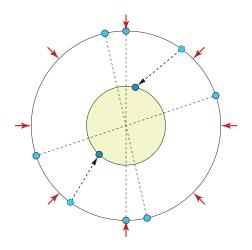
Target locations depend on the walkers' initial positions on SEC



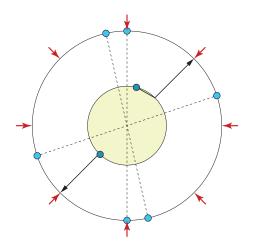
Once in SEC/3, they can only "guess" the initial positions



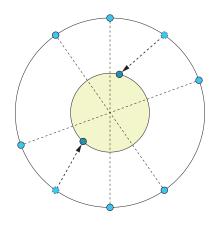
Reasonable guess: we were equidistant from our adjacent robots



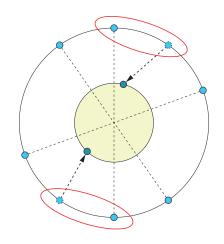
Now they can reconstruct a consistent set of targets...



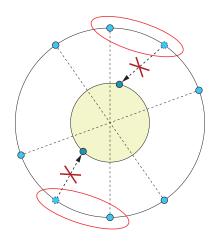
...and move to the appropriate ones



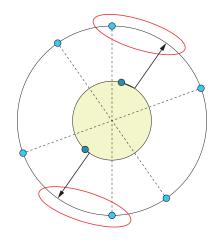
Problem: the guessed locations may not form an analogy class...



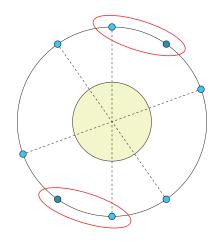
...they may be a proper subset of one!



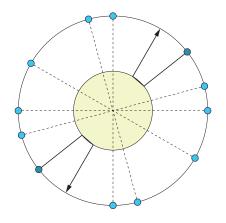
Wrong guess: the walkers are supposed to form an analogy class $% \left(1\right) =\left(1\right) \left(1$



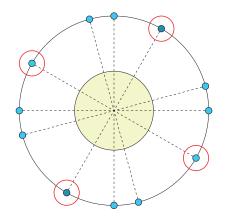
Solution: move to the guessed positions!



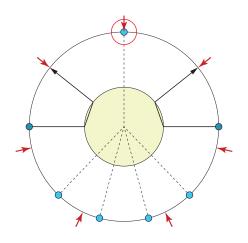
This causes two analogy classes to merge



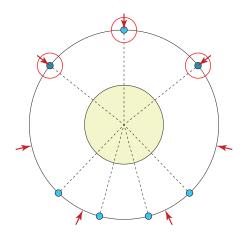
When the walkers complete their journey, two things can happen



Either two analogy classes merge (and the targets may change)...

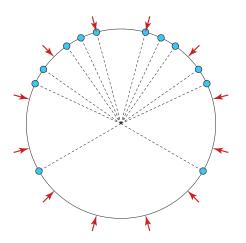


...or more robots reach their targets



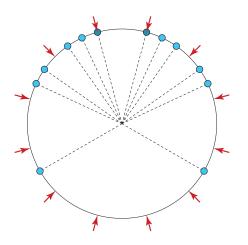
Eventually all robots will reach their targets

Locked configurations



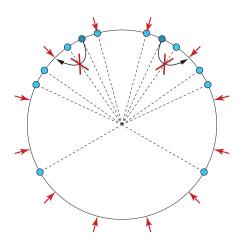
Sometimes no analogy class is able to move to reach its targets

Locked configurations

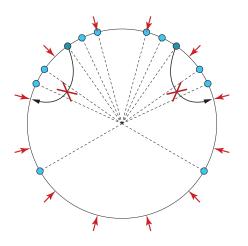


 $\label{eq:entropy} \mbox{Either because it is already there...}$

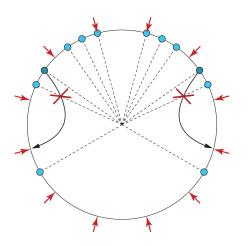
Locked configurations



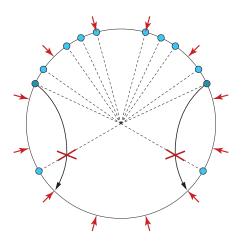
...or because it would form a Co-radial configuration...



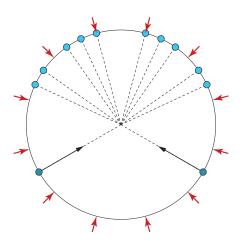
...or because it would form a Co-radial configuration...



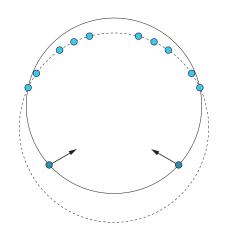
...or because it would form a Co-radial configuration...



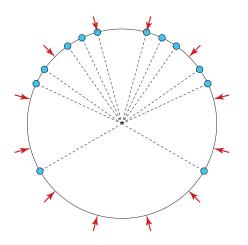
...or because it would form a Co-radial configuration...



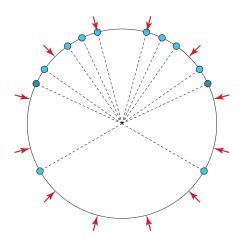
...or because it would alter the SEC



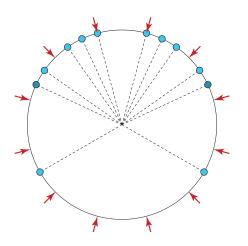
...or because it would alter the SEC



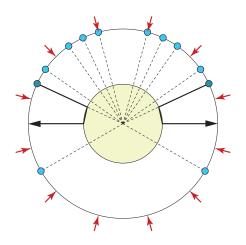
In this case the configuration is locked



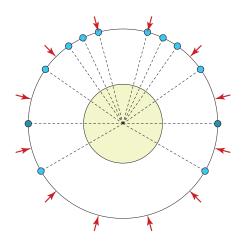
Strategy: identify an unlocking analogy class



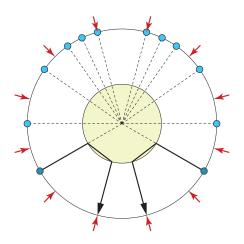
It exists because the configuration is not Half-disk!



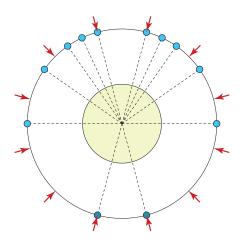
The unlocking class makes a preliminary move...



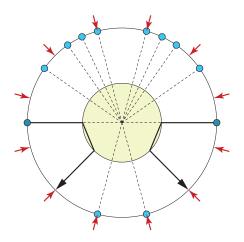
The unlocking class makes a preliminary move...



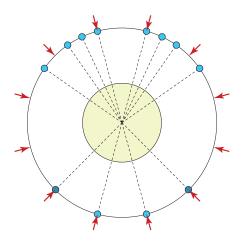
...and the previously unmovable class becomes free to move



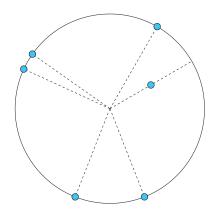
...and the previously unmovable class becomes free to move

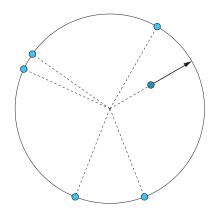


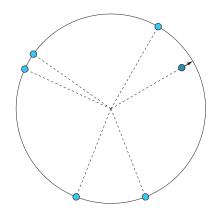
The unlocking step does not make robots lose their progress

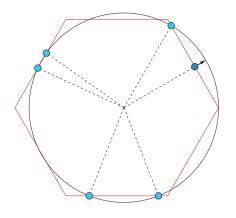


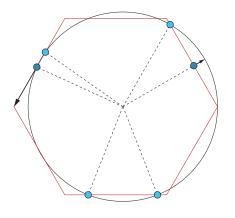
And after the unlocking step, steady progress is made again



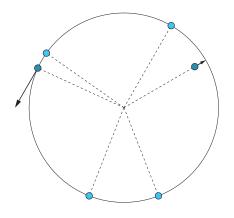




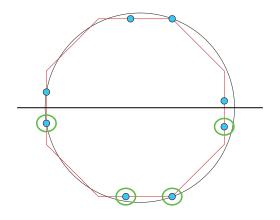




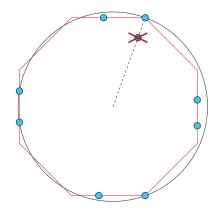
Due to asynchronicity, the behavior may be inconsistent



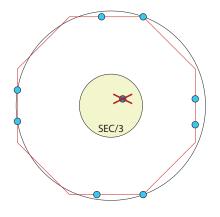
Due to asynchronicity, the behavior may be inconsistent



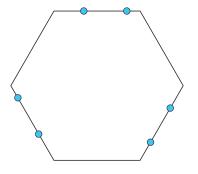
Lemma: no Pre-regular configuration is Half-disk



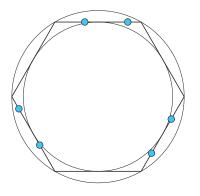
Lemma: no Pre-regular configuration is Co-radial



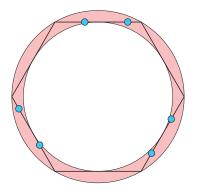
Lemma: no Pre-regular configuration has robots in SEC/3



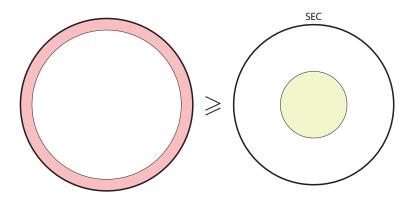
Proof: suppose the configuration is Pre-regular



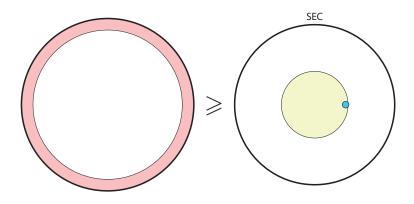
All robots lie in a thin-enough annulus...



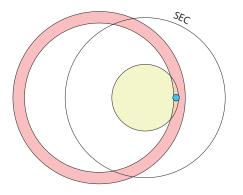
All robots lie in a thin-enough annulus...



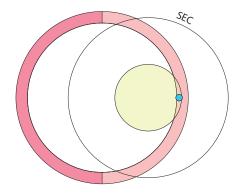
 \ldots whose outer circle is at least as large as the SEC



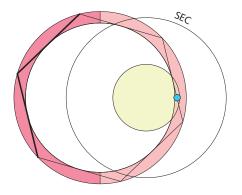
Suppose there is a robot in SEC/3



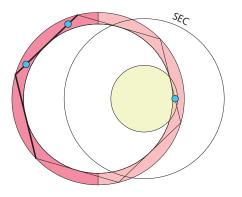
Then the annulus and SEC/3 must overlap



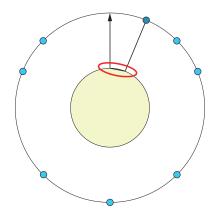
Hence a half-annulus lies outside SEC...



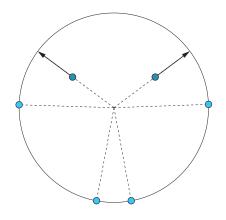
...but it contains two consecutive edges of the polygon...



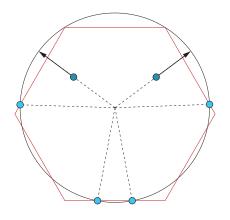
...and one of them must contain robots: contradiction!



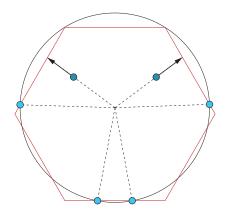
Hence all non-radial moves are safe, as they happen in $\ensuremath{\mathsf{SEC/3}}$



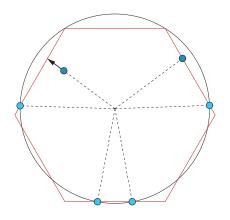
Radial moves are not safe, even if only one analogy class moves



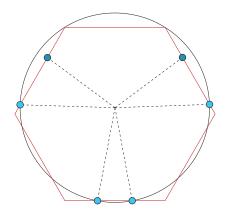
Radial moves are not safe, even if only one analogy class moves



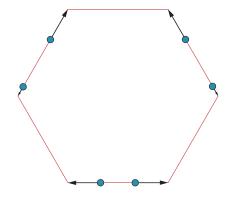
Strategy: add **critical points** corresponding to Pre-regulars...



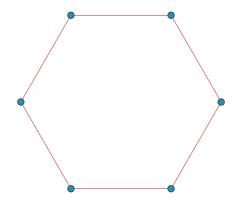
...stop at the next critical point and wait for each other



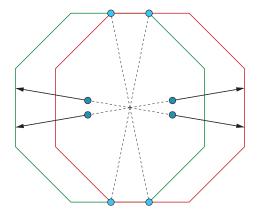
Whenever a Pre-regular is formed, all robots are stopped



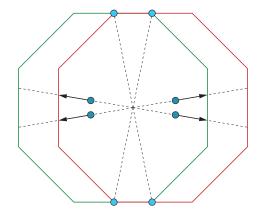
So they correctly transition

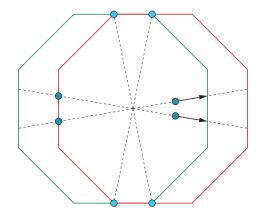


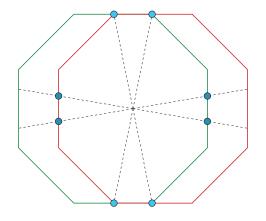
So they correctly transition

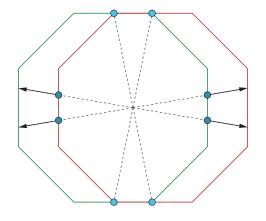


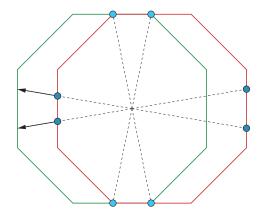
Corresponding critical points may lie on different Pre-regulars

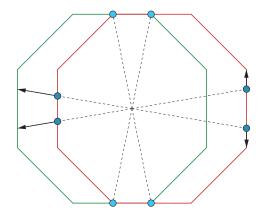


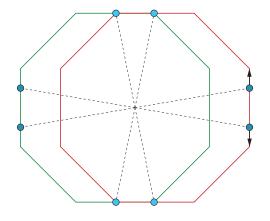


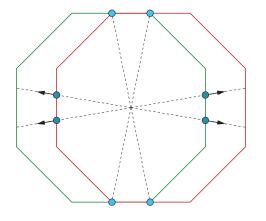




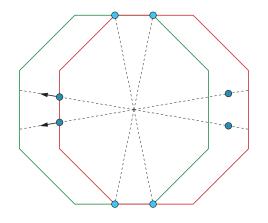


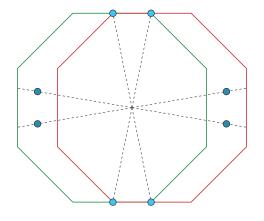


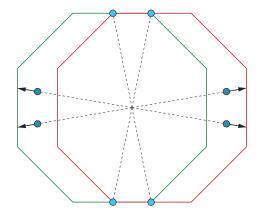


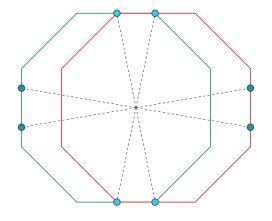


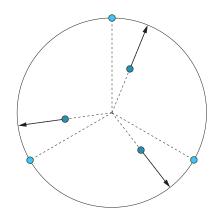
Strategy: add extra critical points between Pre-regulars

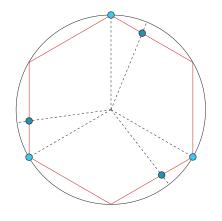


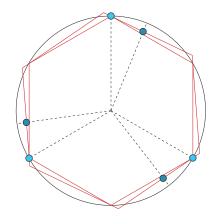


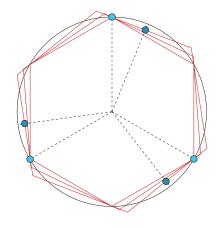


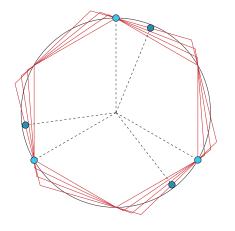


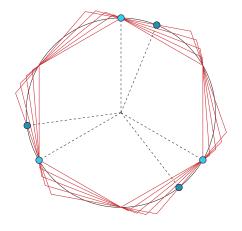




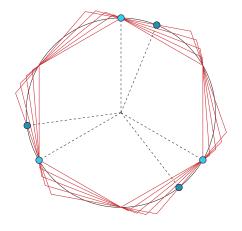




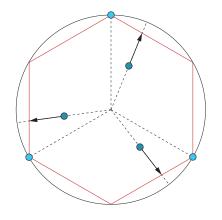




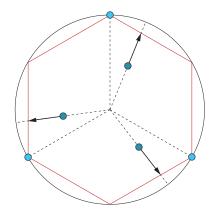
Formable Pre-regulars may be infinitely many!

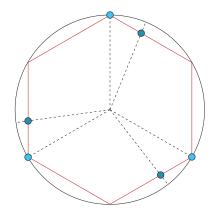


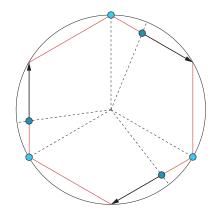
But we cannot have infinitely many critical points

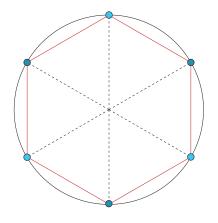


Lemma: a unique Pre-regular is formable before the others



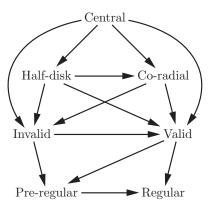






Algorithm flow

As the algorithm proceeds, the current configuration transitions according to this diagram:



There are no loops in the diagram, and the only sink is the Regular configuration. So, in finite time a Regular configuration is reached.

Summary and concluding remarks

The only solvable Pattern Formation problems are:

- Gathering problem (for $n \neq 2$ robots)
- Uniform Circle Formation problem (today we saw a solution for $n \neq 4$ robots: the case with 4 robots is solvable with an ad-hoc algorithm)

This is true even if

- robots are fully synchronous
- robots have a common notion of "clockwise"
- robots always reach their destination
- ⇒ For Pattern Formation problems, these features are computationally irrelevant!

References



P. Flocchini, G. Prencipe, and N. Santoro Distributed Computing by Oblivious Mobile Robots Morgan & Claypool, 2012



G. Viglietta"Uniform Circle Formation"In Distributed Computing by Mobile Entities, Springer, 2019

Assignment 2

In this assignment we study the *Line Formation* problem for asynchronous robots: the goal is to reach a configuration where all robots are on the same straight line. Robots should never collide, and they should permanently stop as soon as they form a line.

We limit our analysis to a swarm of only <u>3 robots</u>:

- Prove that no distributed algorithm can solve the Line Formation problem from an initial configuration in which the robots form an equilateral triangle.
- 2. Give a distributed algorithm that solves the Line Formation problem from any initial configuration that is not an equilateral triangle (assuming that no two robots are initially in the same location), and prove the algorithm's correctness.