
Correspondence problem under isometric transformation as nested linear assignment problems

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Abstract

1 Three-dimensional motion tracking has enjoyed various applications in computer
2 graphics and other fields. One of the technical issues common in optical motion
3 capture systems is time-consuming manual data correction for partially missing
4 points and mis-tracking of objects. There has been several existing algorithms for
5 recovering a point-to-point correspondence between two non-aligned point-sets at
6 two distinct time frames. The many of these algorithms explicitly estimate trans-
7 formation from one set of points to the other, but they often result in sub-optimal
8 solution. Unlike this existing approach, we propose to exploit the distance matrix
9 that is invariant under isometric transformation. This distance-based approach
10 takes a form of quadratic assignment problem, that is efficiently approximated
11 by a nested form of linear assignment problems. In numerical experiments, the
12 proposed nested Hungarian algorithm found the one-to-one correspondence be-
13 tween point-sets with missing points and mis-tracking more accurately than the
14 alternative algorithms.

15 1 Motion tracking by optical motion capture systems

16 Three-dimensional motion tracking has enjoyed various applications in the field of computer graph-
17 ics and biomechanics. In computer graphics applications, such as video games, films, and virtual
18 reality, the motion dataset of human subjects is used to produce realistic motion of the virtual char-
19 acter. In biomechanics, sports science, and their related research fields, recording human or animal
20 movements provides the primary data source for the basic research on kinematic characteristics of
21 their movements.

22 The motion capture system records the shape/posture of a subject at each time point as a collection
23 of points in the three-dimensional coordinates, (x, y, z) , each represents the location of a marker at-
24 tached on the subject. Among various motion capture technologies, optical motion capture systems,
25 composed of multiple infra-red ray cameras and reflective markers, is of a standard choice to capture
26 fine-grained motion patterns due to its high frame rate (over 240 frames per second). In principle,
27 however, it has the technical limitation such that markers would be frequently dropped from the
28 sample due to occluding objects. Since infra-red rays cannot pass through the subject's body, the
29 optical systems cannot detect reflective markers being occluded by the subject's body or something
30 occluders. The other problem is that the markers are often mis-tracked or switched over, when two
31 markers pass similar spatial and temporal point.

32 These two types of errors, dropping and mis-tracking, cause considerable cost in data cleaning, often
33 made manually, before any data analysis of the data collected by the optical motion capture system.
34 In practice, data cleaning, interpolation of the missing markers and identification of markers, is often
35 manually performed and eyeball checked. It is time-consuming (e.g., 40 markers captured for 10
36 minutes in 240 Hz produce 5,760,000 points), and it is often a bottle-neck of the work flow in any

37 use of motion data. The goal of this study is to propose an automated preprocessing algorithm for
38 this class of problems.

39 In 1970s, the researchers of computer vision have formulated the motion tracking as the so-called
40 *correspondence problem* [10]. Since then, the correspondence problem is one of active research
41 topics in computer vision. The correspondence problem is to find a one-to-one correspondence
42 between the two given sets of points (point-sets). In the motion tracking case, each point-set includes
43 multiple points, where each marker at a time frame is represented as a point in the 3D coordinates,
44 and motion tracking is to find a one-to-one correspondence between points across the two point-sets
45 at distinct time frames.

46 A naive algorithm for the correspondence problem would be to match a nearest-neighbor point in
47 a point-set to each point in the other one, under the assumptions that the two point-sets have the
48 same number of points, and the point-set has points quite close to one of them in the other point-set.
49 In practice, this algorithm may work well, if this assumption holds (that is likely, if two point-sets
50 are the consecutive samples at high frame rate). With some markers missing in either/ both of the
51 point-sets, however, this algorithm can fail to find a one-to-one correspondence.

52 In practice, however, a number of markers keep missing for seconds or even longer interval, for
53 example, when the subject’s limbs occlude the markers attached on his/her torso. With such long-
54 term missing data points, the naive algorithm mentioned above cannot work properly. A class of
55 algorithms, e.g., [1], is motivated to find a partial point-to-point correspondence under a certain
56 assumption on the global consistency of the object. For instance, the subject’s body is supposed to
57 be one or more rigid bodies, and a one-to-one correspondence between two point-sets is searched on
58 the basis of point-to-point proximity under a certain isometric (distance-preserving) transformation.

59 In this paper, we present our new algorithm to find a partial correspondence between two point-
60 sets under the assumption that the object is a rigid body. Our algorithm is motivated to exploit the
61 characteristics of the point-set, invariant under isometric transformation. Specifically, our algorithm
62 uses the $n \times n$ and $m \times m$ distance matrices of the point-set with n points and the other point-set with
63 m points, respectively. These distance matrices are invariant under any isometric transformation up
64 to a permutation, and thus, the correspondence problem is reduced to find a permutation for each
65 of these distance matrices. This computational process is closely approximated by a set of linear
66 assignment problems [4]. As each linear assignment problem is solved efficiently by the well-known
67 Hungarian algorithm [7, 11], we call our algorithm *nested Hungarian algorithm*.

68 In Section 2, we briefly illustrate the correspondence problem. In Section 3, we introduce the special
69 case of correspondence problem of a rigid body, and present the nested Hungarian algorithm. In
70 Section 4, we evaluate the proposed algorithm in a series of numerical case studies by comparing it
71 with the existing algorithms.

72 **2 Correspondence for real-world 3D shapes as point-sets**

73 **2.1 The correspondence problem**

74 Suppose that there are n reflective markers, each attached on a body part of a subject, and the
75 markers are fixed on the same points on the body over time. We assume the subject’s body is one or
76 more rigid bodies, but not completely deformable matter.

77 Let $\mathcal{A} \subset \mathbb{R}^3$ and $\mathcal{B} \subset \mathbb{R}^3$ be two point-sets, and each has the n markers (as points) in a three-
78 dimensional Euclidean space. We identify each marker in the point-set by the indices $i = 1, \dots, n$,
79 and denote the i^{th} point in the point-set \mathcal{A} by $\mathbf{a}_i \in \mathcal{A}$ (and the point-set \mathcal{B} by $\mathbf{b}_i \in \mathcal{B}$). The point-sets
80 \mathcal{A} and \mathcal{B} represent the two set of markers recorded at two distinct time frames, but the identity of
81 each marker is not always preserved by the indices across the two point-sets.

82 For now, we suppose there is no missing marker: Every marker in the point-set \mathcal{A} has exactly
83 one corresponding marker in the point-set \mathcal{B} . We will relax this condition later. Denote by the
84 matrix with the vector \mathbf{a}_i at the i^{th} row by $A = (\mathbf{a}_i) \in \mathbb{R}^{n \times 3}$ (and with \mathbf{b}_i at the i^{th} row by
85 $B = (\mathbf{b}_i) \in \mathbb{R}^{n \times 3}$).

86 A one-to-one correspondence between the point-sets \mathcal{A} and \mathcal{B} is a bijective map $p : \{1, \dots, n\} \rightarrow$
87 $\{1, \dots, n\}$, that satisfies a certain condition. In a matrix form, such a bijective map is equivalent
88 to a permutation matrix $P \in \{0, 1\}^{n \times n}$, where the sum of all elements in every row and column is

89 1. For given A and B , the corresponding problem under an unknown isometric transformation is to
 90 find a pair of permutation P and isometric transform $f : \mathbb{R}^{n \times 3} \mapsto \mathbb{R}^{n \times 3}$ such that

$$Pf(A) \sim B, \quad (1)$$

91 where the operator $X \sim Y$ represents a class of similarity between matrices X and Y . Equation (1)
 92 implies that there are two inter-dependent problems in the ‘‘chicken-and-egg’’ relationship [1]. If a
 93 permutation P is given, one can find the isometric transformation f satisfying (1). If f is given, one
 94 can find the P . However, if all P and f are simultaneously unknown, it is not trivial to find P and f
 95 of a desired condition.

96 2.2 Transformation-based and distance-based approach

97 In past literature, there are at least two major classes of approach to the correspondence problem (1).
 98 The first approach is to explicitly estimate the isometric transformation f as well as P in an iterative
 99 manner, which we call *transformation-based* approach. The second one is to estimate only the
 100 permutation P by using the distance (dissimilarity) matrices constructed from A and B , which we
 101 call *distance-based* approach. The transformation-based one is computationally cheaper, but it may
 102 result in sub-optimal solution, due to its iterative procedure for dual-minimization of both P and f
 103 [12]. The distance-based approach essentially avoids the chicken-and-egg problem and need to find
 104 only permutation P , as the distance matrices are invariant under any isometric transformation. It is,
 105 however, computationally extremely costly, as its exact search is known as a NP-complete problem.

106 In this study, we propose a computationally feasible algorithm as an approximation of the distance-
 107 based approach.

108 3 Correspondence under isometric transformations

109 3.1 Distance-based formulation of the correspondence problem

110 The key observation of the distance-based approach is the distance matrix $D(A) =$
 111 $(d_{i,j}^A)_{i,j \in \{1, \dots, n\}} \in \mathbb{R}^{n \times n}$, where $d_{i,j}^A = \|A_i - A_j\|$ and $A_i \in \mathbb{R}^3$ is the i^{th} row of A , is invari-
 112 ant under isometric transformation f : i.e., $D(f(A)) = D(A)$. Thus, the correspondence problem
 113 (1) is reformulated to find a permutation $P \in \mathbb{R}^{n \times n}$ such that

$$PD(A)P^\top \sim D(B). \quad (2)$$

114 Introducing an error function, such as the summed squared error $E = \|PD(A)P^\top - D(B)\|$ where
 115 $\|X\| = \sum_{i,j} X_{i,j}^2$, it is a quadratic assignment problem to find the minimizer P of E , that is known
 116 as a NP-complete problem [3].

117 3.2 Nested linear assignment problems

118 Therefore, we reformulate the problem (2) by a hierarchical linear assignment problem. To illustrate
 119 our idea, consider the special case $PD(A)P^\top = D(B)$ with an exact correspondence between \mathcal{A}
 120 and \mathcal{B} . In this case, with such permutation P , there is some permutation $Q \in \mathbb{R}^{n \times n}$ independent of
 121 P satisfies

$$PD(A)Q^\top = D(B), \quad (3)$$

122 and $Q = P$ is unique, if there is no exchange of any two rows in A that preserves the distance matrix
 123 $D(A)$. This means that there is permutation Q such that

$$(d_{j,1}^A, \dots, d_{j,n}^A)^\top = Q(d_{i,1}^B, \dots, d_{i,n}^B)^\top, \quad (4)$$

124 for all i, j such that $P_{i,j} = 1$ or i^{th} point in \mathcal{B} corresponds with j^{th} point in \mathcal{A} . For each pair (i, j) ,
 125 Q in (4) is solved by a linear assignment problem

$$\hat{Q} := \arg \min_Q \text{tr}(D_{i,j}Q), \quad (5)$$

126 where $D_{i,j} := (\|d_{j,k}^A - d_{i,l}^B\|)_{k,l \in \{1, \dots, n\}} \in \mathbb{R}^{n \times n}$. Then the permutation P is given by another
 127 linear assignment problem

$$\hat{P} := \arg \min_P \text{tr}(CP), \quad (6)$$

ALGORITHM 1: Nested Hungarian algorithm

Input: Two point-sets $A \in \mathbb{R}^{n \times 3}$ and $B \in \mathbb{R}^{m \times 3}$

Output: Correspondence P

$D(A) = (d_i^A)$ is the distance matrix of A

$D(B) = (d_j^B)$ is the distance matrix of B

$C \in \mathbb{R}^{n \times m}$

for $i = 1$ **to** n **do**

for $j = 1$ **to** m **do**

 Compute $[\hat{Q}, c_{ij}]$ using Hungarian($D_{i,j}$)

 Set the (i, j) element of C to c_{ij}

end

end

Compute $[\hat{P}, c]$ using Hungarian(C)

128 where $C = (c_{i,j}) \in \mathbb{R}^{n \times n}$ with $c_{i,j} := \min_Q \text{tr}(D_{i,j}Q)$. This observation holds only if there is
129 no exact match $PD(A)P^\top = D(B)$, but it reduces the original quadratic assignment problem (2)
130 to the nested linear assignment problem of (5) and (6). As the linear assignment is solved by the
131 Hungarian algorithm calling the function $O(n^3)$ times [4], this nested linear assignment problem is
132 solved by $O(n^5)$ times of the function calls.

133 In summary, this procedure mentioned above is implemented by the nested Hungarian algorithm
134 described in the psudo-code (Algorithm 1) for the point-set \mathcal{A} with n points and \mathcal{B} with m points.
135 In the pseudo-code, the function Hungarian(\cdot) is an implementation of the Hungarian algorithm,
136 computing the minimal total cost c_{ij} and the correspondence \hat{Q} . Note that the nested **for** loops can
137 be computed independently. Thus, the construction of C in Algorithm 1 can be parallelized, and
138 then its computational complexity will be reduced to $O(n^3)$.

139 4 Related works

140 In the following section, we will report numerical studies comparing the nested Hungarian algorithm
141 with the representative existing algorithms. Here we briefly overview them. Some of the following
142 algorithms for shape/feature matching is not developed specifically for the correspondence problem
143 of interest, but for more general problems than it. Here we choose the existing algorithms according
144 to their applicability to the correspondence problem under isometric transformation.

145 As an algorithm taking the transformation-based approach, Basl [1] has proposed Iterative clos-
146 est point (ICP) algorithm that finds a locally minimal correspondence by iteratively searching the
147 nearest-neighbor and the least-square rotation and translation. However, this algorithm often results
148 in a local minimum, that is not sufficiently good in practice [12]. An extension of ICP algorithm [5]
149 for non-rigid shape matching approximates non-rigid deformable subjects as a patchwork of small
150 rigid segments.

151 As an algorithm taking the distance-based approach, Maciel [9] and Berg [3] have proposed an
152 algorithm to match image patch of a set of pixels, by formulating it as a constrained concave pro-
153 gramming problem. Leordeanu [8] has proposed the spectral matching algorithm, that is motivated
154 by the graph/network analysis techniques. It solves a Google’s Page rank-like problem on a graph
155 of points (as nodes) $\mathcal{A} \cup \mathcal{B}$ and then extracts a node-to-node correspondence from the visiting fre-
156 quency of the nodes. This spectral matching algorithm is computationally efficient and reported its
157 good performances in practice [12].

158 5 Experiments

159 We tested our algorithm by comparing it with existing algorithms. Our motivation here is how
160 accurately our algorithm finds the underlying one-to-one correspondence, rather than the running
161 speed of algorithms. Thus, we analyzed the accuracy of estimated correspondence P under several
162 perturbations, including measurement noise and isometric transformations (rotation and translation).
163 For the notational simplicity, we identify point-set \mathcal{A} and its matrix representation A in this section.

164 5.1 Implementation of existing algorithms

165 We compared our algorithms with the existing ones proposed in [1, 9, 3, 8]. We implemented
166 Iterative closet point algorithm [1] using the k-d tree [2] and the least-square algorithm [6]. For
167 the shape-in-image matching algorithms [9, 3, 8], instead of the pixel-based point-sets in high-
168 dimensional feature space, we just used the 3D coordinates in the Euclidean space, and their dissim-
169 ilarity matrix was constructed from distances within A , within B , and between A and B .

170 5.2 Tolerance against noise and invariance under isometric transformations

171 To test these algorithms, we define the accuracy by the ratio of the number of datasets for which each
172 algorithm found the exactly correct correspondence to the number of all datasets. Specifically, (1) a
173 point-set A of size $n \times 3$ are uniformly randomly generated within $[0, 1]^{n \times 3}$. Then, (2) generate the
174 opposite point-set by $B = AR^\Delta + \mathbf{1}_n T + E$ for a given matrix A , where $\mathbf{1}_n \in \mathbb{R}^{n \times 1}$ is the vector
175 with every element being 1, $R^\Delta \in \mathbb{R}^{3 \times 3}$ is randomly generated rotation matrix, $T \in [0, \tau]^{1 \times 3}$ is a
176 translation vector with uniformly random values, and $E \in \mathbb{R}^{n \times 3}$ is element-wise noise with each
177 element drawn from the normal distribution $N(0, \epsilon)$ with the variance ϵ^2 . The rotation matrix R
178 is generated via QR decomposition of uniform random matrix, and R^Δ is R to the power of Δ .
179 To evaluate the accuracy, each algorithm was tested for 100 randomly-generated A and B for each
180 combination of Δ , τ , and ϵ . Here we set the number of points in the point-set A to be $n = 10$.

181 Figure 1(a) shows the accuracy against the power $\Delta = 0, 0.1, \dots, 1.2$ to rotation given $\tau = \epsilon = 0$
182 and (b) the accuracy against translation $\tau = 0, 0.1, \dots, 1.2$ given $\Delta = \epsilon = 0$. Our algorithm
183 ‘Proposed’ shown as red markers produced the correct correspondence P exactly 100% for all the
184 rotations and translations. This result shows that the algorithm is robust for a wide range of isometric
185 transformations. This robustness comes from the nature of our algorithm based on the distance mat-
186 rix, is the invariant under isometric transformations. We found ‘Berg2005’ [9, 3] also performing
187 as good as ours against rotation and translation, whose result overlaps underneath of the ‘Proposed’
188 one in Figure 1 (a) and (b). Iterative closet point algorithm ‘Basl1992’ [1] showed lower accuracy
189 against larger rotations and translations, and the result shows its limitation under isometric transfor-
190 mations.

191 Figure 1(c) shows tolerance against measurement noise $\epsilon = 0, 0.01, \dots, 0.12$ given $\Delta = \tau = 0$
192 and (d) against noise ϵ combined with rotation $\Delta = 1$ and translation $\tau = 1$. In the conditions of
193 noise effect alone (Figure 1(c)), our proposed algorithm performed worse than ‘Berg2005’ [9, 3] and
194 ‘Basl1992’ [1] as the scale of noise increased. This is probably due to the two algorithms ‘Berg2005’
195 [9, 3] use the distances between points within each point-set A or B , as well as distances between
196 the point-sets A and B . In contrast, our proposed algorithm is based on the distances within each
197 point-set A or B . However, combining noise and transformations shown in Figure 1(d), our proposed
198 algorithm outperformed others, since our algorithm less affected by isometric transformations. Thus,
199 these analyses with noise and isometric transformations revealed both advantage and disadvantage
200 of our algorithm.

201 5.3 Tolerance against missing values: partial subset correspondence

202 Next, we tested our algorithm against missing values. Specifically, we generated two point-sets A
203 and B in the same procedure above. Then, after generated A and B , 7 rows (7 points) of A were
204 dropped as missing values. By this, all the corresponding points of A still exist in B but some points
205 of B has no corresponding points in A , i.e., the problem of injective partial correspondence. In this
206 analysis, the accuracy is defined by the ratio of the number of datasets for which each algorithm
207 could find the exactly correct correspondence between non-missing points between A and B to the
208 number of all datasets. Here we supposed that the all 3 remaining points of A are on the same rigid
209 body of the subject, but the other 7 omitted points of A can be on arbitrary rigid bodies other than
210 this one. In other words, this experiment was intended for the cases finding a partial (subset) corre-
211 spondence between a single rigid body A and multiple rigid bodies B , or finding a corresponding
212 subset of A in B .

213 Figure 2(a)–(d) show the results of partial correspondence under the same conditions as Figure 1(a)–
214 (d). In Figure 2(a) rotation only and Figure 2(b) translation only, our proposed algorithm achieved
215 100% of accuracy. But, for these conditions, ‘Berg2005’ [9, 3] greatly decreased its accuracy of
216 correct response, in contrast to the full correspondence case (see Figure 1(a) and (b)). In (c) noise

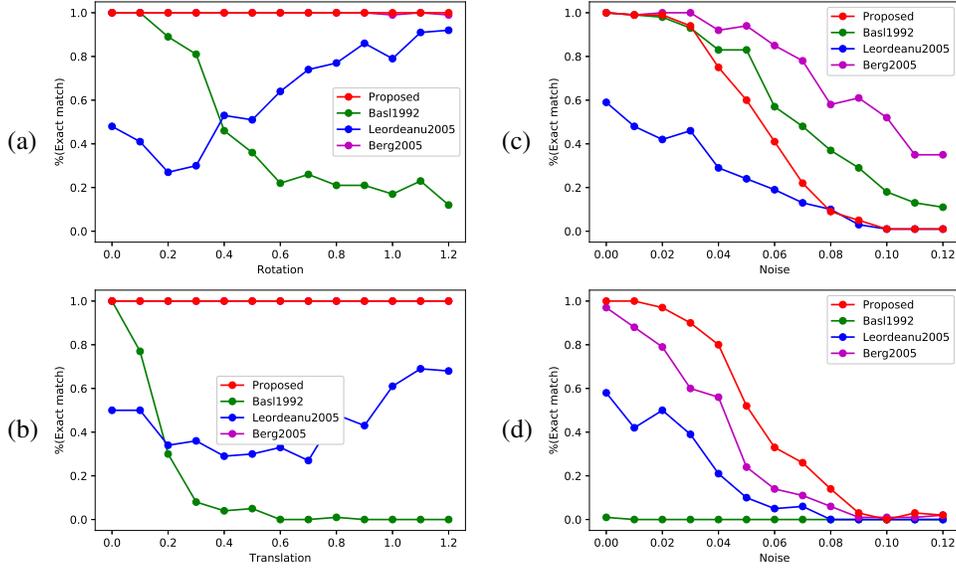


Figure 1: Rates of producing correct correspondence in several conditions. The number of points in \mathcal{A} and \mathcal{B} is both 10, and no missing value. Types of perturbations differ: (a) rotation only, (b) translation only, (c) noise only, (d) combination of noise, rotation and translation.

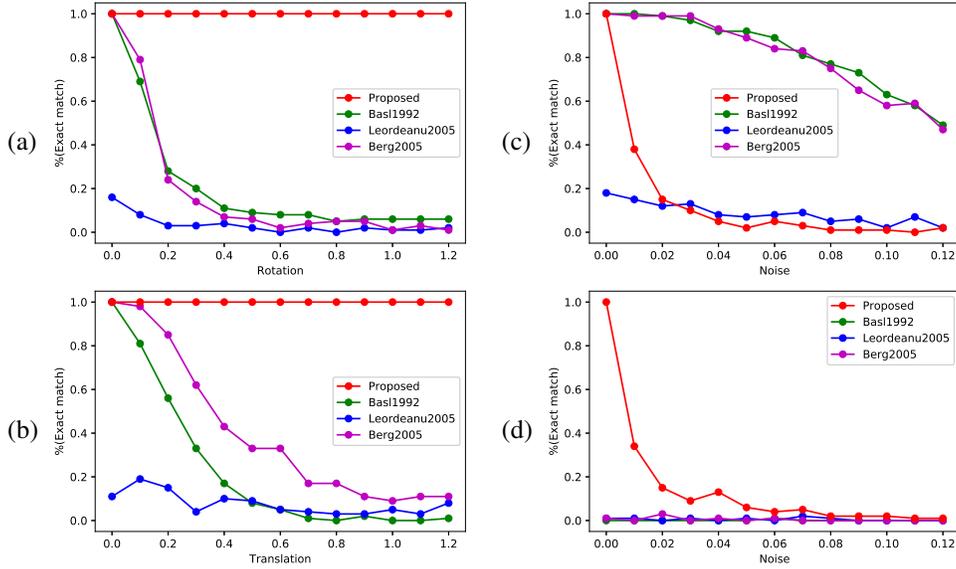


Figure 2: Rates of producing correct “partial” correspondence in several conditions. The number of points in \mathcal{A} is 3 and in \mathcal{B} is 10. There are 7 missing values. Types of perturbations differ: (a) rotation only, (b) translation only, (c) noise only, (d) combination of noise, rotation and translation.

217 only, our proposed algorithm hardly produced correct partial correspondence under largely noisy
 218 conditions, but ‘Berg2005’ [9, 3] and ‘Basl1992’ [1] kept higher accuracy overall. It is probably
 219 due to the use of distances between point-sets. However, again in (d) combination of noise and
 220 transformations, our proposed algorithm relatively worked better than other existing algorithms.
 221 Again, this analysis on the partial correspondence shows the advantage of using the invariant under
 222 isometric transformations.

223 6 Discussion

224 In applications to realistic data recorded by optical motion capture systems, the relationship between
225 the smallest distance between points and the scale of noise matters. The precision of a recent optical
226 motion capture system [13] achieves up to 0.5 mm. When using a standard marker placement,
227 the smallest distance between two markers is about 100 mm. In the experiments in the previous
228 section, the smallest distance between two points is in average about 0.165. Thus, the corresponding
229 precision (scale of noise) in the previous experiments is about 0.000825 ($= 0.5 \times 0.165/100$).
230 At this scale of noise $\epsilon = 0.000825$ with random Euclidean transformations ($\Delta = \tau = 1$), the
231 accuracy or the rate of exactly correct response over 97% in average was achieved by our proposed
232 algorithm for finding partial (subset) correspondence (as Figure 2(d)). However, at the same scale
233 of noise, the other algorithms achieved less than 2% of correct response. Thus, when to find full and
234 partial correspondence between two point-sets at two largely-apart time frames, the proposed nested
235 Hungarian algorithm alone is effectively applicable to a recorded data with optical motion capture
236 systems.

237 Our algorithm currently has limitation in finding a correspondence between two multiple rigid bod-
238 ies. Extension of our algorithm for two multiple rigid bodies is our on-going future work.

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