Scale-free networks on a geographical planar space for efficient ad hoc communication

Yukio Hayashi, and Jun Matsukubo

yhayashi@jaist.ac.jp

Japan Advanced Institute of Science and Technology

1. Introduction

Dynamic configuration of (backbone) networks

Problems for ad hoc communication such as in sensor networks or P2P systems



dissipation of wireless beam-power or line-cost for long-range links
interference by crossing links

a few nodes have sufficient power to be hubs

2. Efficient Routing

Planar triangulation: reasonable math. abstraction of ad hoc net. (each triangle forms a service region) Moreover, a memoryless, no defeat, and competitive online routing algorithm has been developed for networks on triangulation.



3-1. Scale-free Nets with Hubs

Existing a surprisingly common structure: SF net. the degree dist. exhibits $P(k) \sim k^{-\gamma}$, $2 < \gamma < 3$.

Social: acquaintance, world trading, actor-collabo., citation, language

Technological: Internet, WWW, email, power gridBiological: neural net, genom, metabolic pathway, foodweb

One of the fundamental generation mechanism has been proposed: Growth & Preferential Attachment Barabási and Albert, Physica A, 272, 1999

3-2. BA model: $P(k) \sim k^{-3}$

- A network grows from initial N_0 nodes with $m < N_0$ links among them.
- At every time step t, a new node is introduced, and is randomly connected to m previous nodes i with an attachment probability $\prod_{i}^{BA}(t) \sim k_i(t)$.



3-3. Properties of the SF Nets

Efficient & Economy with short paths as the hop-count O(ln N) and the low cost of links as few as possible in the connectivity (Solé et al., Advances in Complex Systems 5, 2002),
 Robust connectivity for random failures (Albert et al., Nature 406, 2000)

3-3. Properties of the SF Nets

- Efficient & Economy with short paths as the hop-count O(ln N) and the low cost of links as few as possible in the connectivity (Solé et al., Advances in Complex Systems 5, 2002),
 Robust connectivity for random failures (Albert et al., Nature 406, 2000)
- However, most of SF net models were irrelevant to a geographical space.

3-3. Properties of the SF Nets

- Efficient & Economy with short paths as the hop-count O(ln N) and the low cost of links as few as possible in the connectivity (Solé et al., Advances in Complex Systems 5, 2002),
 Robust connectivity for random failures (Albert et al., Nature 406, 2000)
- However, most of SF net models were irrelevant to a geographical space.
- Therefore, we consider geographical SF nets, especially as planner graphs without crossing links to avoid interference of wireless beams.

4-1. Rare Long-range Links



4-2. Geographical SF Nets

- Modulated BA: $\Pi_i \sim k_i \times l^{\alpha}$, rand. position (a of node
- SF on lattices: connect within $r = A \times k_i^{1/d}$
- Space-filling: subdivision of a region (heterogeneous dist. of nodes)



4-3. Planarity and Shortness

class	planarity of net	shortness of links
Modulated BA	×	\bigcirc
Manna'02,	∃ crossing links	with disadvantaged
Xulvi-Brunet'02	(not prohibited)	long-range links
SF on lattices	×	\bigtriangleup
ben-Avraham'03,	cross of regular	\exists long shortcuts
Warren'02	links and shortcuts	from hubs
Space-fi lling	\bigcirc	\bigtriangleup
Apollonian nets.	by subdivision	∃ long-range links
Doye'05, Zhou'04	of a selected region	in narrow triangles

4-4. Space-filling Packing



The dual by connecting the centers of touched circles is nothing but the (deterministic) Apollonian network. When a triangle is randomly selected for subdivision, it's called RAN.

4-5. Analysis of P(k) in RAN

Evolution eq. by iterative triangulation:

$$n(k+1, N+1) = \frac{k}{N_{\triangle}}n(k, N) + \left(1 - \frac{k+1}{N_{\triangle}}\right)n(k+1, N),$$

where N_{\triangle} denotes the number of triangles. In the $P(k) \approx n(k, N)/N$, it can be rewritten as

$$(N+1)P(k+1) = NkP(k)/N_{\triangle} + NP(k+1) - N(k+1)P(k+1)/N_{\triangle}.$$

By the continuous approx., we obtain $P(k) \sim k^{-\gamma_{RA}}$ with $\gamma_{RA} = (N_{\triangle} + N)/N \approx 3$ for large N.

Zhou et al., arXiv:cond-mat/0409414, 2004

5-1. Voronoi and Delaunay

Consider the combination of RA (by triangulation on a plane) and DT to avoid the long-range links



 \Rightarrow optimal triangulation in some criteria: maximim angle, minimax circumscribed circle, short path length close to the direct Euclidean dist., etc.

5-2. Delaunay-like SF Net

We propose RA+NN:

- Set an initial planar triangulation.
- Select a triangle at random and add a new node at the barycenter. Then, connect the new node to the three nodes of its triangle. By iteratively applying diagonal flips, connect it to the nearest node(s) within a radius.



6-1. Simulation

We compare the topological properties of the models: RA, DT, and our proposed RA+NN in the averaging of 100 realizations at size N = 10,000.



RA

DT

RA+NN(one)

\Rightarrow RA+NN has the intermediate structure.

6-2. Degree Dist. & Correlation



 \Rightarrow RA+NN has a power law dist. with an exp. cutoff, and a weak negative deg.-deg. correlation.



6-4. Distance for Movement



NOLTA 2005 - p.17/22

6-5. Min. Hops for Transfer



NOLTA 2005 - p.18/22

7. Summary

- We've briefly reviewed recent studies of geographical SF net models, and proposed a modified one to reduce long-range links.
- The Delaunay-like SF net without crossing links is dynamically constructed by the triangulation and diagonal flips based on local rules.
- Simulation results have shown that our proposed model has short path lengths and small num. of hops (DT is not opt. in minimizing them), which are suitable topological properties for efficient communications.

We'll further investigate the traffic properties and the fault-tolerance.

Appendix 1.

Average degree k(s) of the node inserted at time s, and the betweenness B(k) of the nodes with degree k



 \Rightarrow Old nodes of RA and RA+NNs tend to be hubs, and the traffic load of RA+NNs is the intermediate

Appendix 2.

Estimated parameters for the degree distribution in each model by a NMSE method

	model	estimated function	parameters
	RA	$P(k) \sim k^{-\gamma_{RA}}$	$\gamma_{RA}pprox 3$
	DT	$P(k) \sim \exp\left(-\frac{(\ln k - \mu)^2}{2\sigma^2}\right)$	$\mu = 1.7755, \sigma = 0.2383$
RA	+NN(one)	$P(k) \sim k^{-\gamma} \exp(-ak)$	$\gamma = 2.26,$
			a = 0.0647
RA	+NN(all)	$P(k) \sim k^{-\gamma} \exp(-ak)$	$\gamma = 1.7248,$
			a = 0.0979

Appendix 3.

Assortative and Disassortative correlations observed in social and technological/biological networks







Dis: between hub and peripheral nodes with low degrees