

Topic 26

Jackson's Description Principles

- The **prerequisite** for following this (part of the) lecture is that you have followed the lectures on Description Principles and On Defining and On Definitions.
- The **aims** are
 - ★ to introduce Jackson's concepts of designations, definitions and refutable assertions, and
 - ★ to provide formal tools for the expression of manifestations of these concepts.

- The **objective** is
 - ★ to ensure that the developer becomes a professional specifier.
- The **treatment** is systematic to formal.

We build on ideas eloquently expressed in:

Software Requirements & Specifications
a lexicon of practice, principles and prejudices
Michael Jackson
ACM Press and Addison-Wesley Publishing Company
ISBN 0-201-87712-0, 1995; xvi+228 Pages

- Since all we do is
 - ★ construct, analyse and compare descriptions
 - ★ we shall analyse
 - ★ the concept and constituents of descriptions.⁸

⁸We shall here use the term description as also covering the terms prescription and specification.

- A description is about
 - ★ manifest individuals, i.e., phenomena and concepts.
 - ★ Some of these represent, or are intended to represent, facts;
 - ★ others represent mental constructions, i.e., concepts.
 - ★ A description includes *designations*, definitions and refutable descriptions.
 - ★ A description can either be formal or informal.
 - ★ A description sets a scope and a span, and a description
 - ★ expresses moods.
- By an individual we mean a physically manifest phenomenon.

Phenomena, Facts and Individuals

- Phenomena are what appears to exist.
Domain phenomena are built up from facts about individuals.
- A fact is a simple truth about the world:
It is the smallest unit of observation.
- Large and complex observations and truths can be broken down into assertions about several facts.

- A fact involves one or more *individuals*. Anything can be an individual. Each individual is identical to itself but distinct from all other individuals.
- If x is identical to y , then x and y are the same individual: $x = y$.
- If we say that x is similar or equivalent to y then x and y are distinct individuals that share some property, characteristic, attribute or quality.

Designations

Characterisation 7.134 • By a *designation description* (for short: *designation*) we syntactically mean a textual triple:

- ★ a *designation term*,
- ★ a *designation recognition rule*,
- ★ and a *designation identification*

Characterisation 7.135 A *designation term* is a simple, i.e., atomic name.

Characterisation 7.136 A *designation recognition rule* is a text which purports to *designate* something.

Characterisation 7.137 • The *designation identification* links the *designation recognition rule* to that something (mentioned in the previous characterisation)

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 694, Topic: 26, Foli: 9
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

• To create a *designation* we thus write down two items:

- ★ a *designation description*
- ★ and a *designation identification*.

• The *designation description* is usually of the form:

- ★ *Designated term*: **dt**.
- ★ *Recognition rule*: **dt**'s satisfy the following informally stated properties: ... — where all other words, i.e., terms, are assumed to be well-known!

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 695, Topic: 26, Foli: 10
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

• That is: *recognition rules* are expressed only in terms that are otherwise well understood, i.e., part of the folklore.

• If a recognition rule thus contains a term that is elsewhere *defined*,

- ★ then it is not a recognition rule
- ★ but becomes a *definition*.

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 696, Topic: 26, Foli: 11
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

Example 7.79 Rail Units, I: We give the first in a series of examples relating to rail nets.

- *Designated term*: **rail_unit**, or just **U**.
- *Recognition rule*: A **rail_unit** is a composition of an even number of parallel positioned rails (long, narrow, profiled iron bars) separated such that one can always identify pairs of rails of the composition that are at a specified distance (the rail gauge), and otherwise held together by a set of ties.

When we write **rail_unit** we mean the extensional meaning (the **type**) of that term. ■

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 697, Topic: 26, Foli: 12
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

Example 7.80 Rail Units, II: We give the second in a series of examples relating to rail nets.

- *Designated term*: **linear_rail_unit**, or just **U**, for which a predicate, to be assumed, holds: **is_linear_rail_unit(u)** for all **u** in the extension **U**, i.e., **u:U**.
- *Recognition rule*: A **linear_rail_unit** is a pair of parallel positioned rails (long, narrow, profiled iron bars) separated at a specified distance (the rail gauge) and held together by a set of ties. (The predicate **is_linear_rail_unit(u)** “arises” from the recognition rule.)

Observe that the designation term and the recognition rule dealt with a concrete concept for which the immediate concretisation points to a set of a phenomena. What is being designated and to be recognised is any one member of this set. ■

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 698, Topic: 26, Foli: 13
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

- The latter example is a specialisation of the former.
- The former example expressed: ... *composition of an even number of parallel positioned ‘rails’* ...
- thus allowing for, say, two pairs as in a switch or in a simple crossover rail unit.
- A *designation identification* is usually of the form:
 - ★ “That ‘thing’ (*u*) there is a *U*.”
 - ★ So is that ‘thing’ (*u'*) over there!” —
- thus physically pointing out a *designation set*, that is, instances of values that together become part of a type.

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

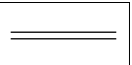
SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Designations		April 5, 2006, 09:18	Page 699, Topic: 26, Foli: 14
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

Some Observations


- In general a *designation description* can be formalised;
- but one cannot formalise the *identification* — as it relates a formal world to an inherently informal world.
- As illustrated in the two previous examples,
- the latter of which was a specialisation of the former,
- a *designation* may extensionally denote a class (of things) which can be subdivided into subclasses.

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

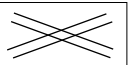
SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3.1 Some Observations		April 5, 2006, 09:18	Page 700, Topic: 26, Foli: 15
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark



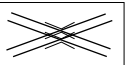
Linear Unit



Switchable Unit



Simple Crossover



Switchable Crossover

Figure 7.12: Roughly drawn rail units

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3.1 Some Observations		April 5, 2006, 09:18	Page 701, Topic: 26, Foli: 16
home:/db/voll8/3a7/3a7			Richard Porenskov Plank, DK-2800 Kong Lyngby, Denmark

Example 7.81 Rail Units, III: We give the third in a series of examples relating to rail nets.

- *Designated term*: **rail_unit**.
- *Recognition rule*: — as for the above, previous **rail_unit** example.
- *Designated term*: **linear_rail_unit**.
- *Recognition rule*: — as for the above, previous **linear_rail_unit** example, and additionally: Thus a **linear_rail_unit** has two “ends” and a single (two-way) link between these ends. Any other **rail_unit** may be **connected** to either of these ends (and if so, then the end is called a **connector**).

4435 2726, Fax: +45 4450 5874 © Olsen Egeboer, Fønsting 11, DK-2800 Høvd, Denmark E-mail: dd@immi.dtu.dk, sp@immi.dtu.dk, d@immi.dtu.dk, URL: www.immi.dtu.dk DTU

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.1 Some Observations		Page 702, Topic: 26, Foli: 17	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- To support textually expressed recognition rules it is often useful to deploy graphic means as in Fig. 7.12.
- Even photographic means could be used.
- And then one usually could show several photos of variations of the same kind of designated individuals.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Some Observations		Page 703, Topic: 26, Foli: 18	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- **Designated term:** `switchable_rail_unit`.
- **Recognition rule:** — as for `rail_unit` just above, and additionally: A `switchable_rail_unit` has three **connectors** which “define” (allow, permit) two two-way links through the `switchable_rail_unit`.
- **Designated term:** `simple_crossover_rail_unit`.
- **Recognition rule:** — as for `rail_unit` just above, and additionally: A `simple_crossover_rail_unit` has four **connectors** which “define” (allow, permit) two two-way links through the `simple_crossover_rail_unit`.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.1 Some Observations		Page 704, Topic: 26, Foli: 19	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- **Designated term:** `switchable_crossover_rail_unit`.
- **Recognition rule:** — as for `rail_unit` just above, and additionally: A `switchable_crossover_rail_unit` has four **connectors** which, together with the switching ability of the unit, “define” (allow, permit) four two-way links through the `switchable_crossover_rail_unit`.
- **Designated term:** `connector`.
- **Recognition rule:** A `connector` is that which allows two `rail_units` to be connected, “end to end”.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Some Observations		Page 705, Topic: 26, Foli: 20	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- Further recognition rules deal with **connectors** and **rail_units**.
- A **connector** is any “end” of a **rail_unit** to which other **rail_units** may be ‘connected’.
- Any one **connector** is shared by at most two **rail_units**.
- A **rail_unit** is mutually exclusive either a **linear_rail_unit** or a **switchable_rail_unit** or a **simple_crossover_rail_unit** or a **switchable_crossover_rail_unit**.

See axioms [2,3] in a formalisation given later. ■

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.1 Some Observations		Page 706, Topic: 26, Foli: 21	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- As we shall see in a later lecture,
- doubt may be raised as to whether some of the text parts of the above recognition rules express assertions.
- (Designations state facts, not assertions.)
- For example: “has two ends” (or three, or four).
- These parts are part of recognition rules,
- but what about:
 - ★ “any one connector is shared by at most two rail units”?
- Or the part right after it:
 - ★ “... mutually exclusive ...”?

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Some Observations		Page 707, Topic: 26, Foli: 22	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- For the last (the mutual exclusion) we can claim it to be a fact,
- hence it is part of a recognition rule.
- For the “at most two”, the case is a bit more complicated.
- And then,
 - ★ when we formalise the whole thing, as we shall see below,
 - ★ and when such a formalisation is compared to that of a refutable assertion,
 - ★ we shall see that, formally speaking, the difference is almost invisible.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.1 Some Observations		Page 708, Topic: 26, Foli: 23	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

- ★ Thus we must be prepared for the eventuality
 - ◇ that the pragmatics of making a distinction between designations, definitions and refutable assertions
 - ◇ can not be carried visibly over into formal models of the things being designated or defined, or for which assertions are expressed.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Formalisation		Page 709, Topic: 26, Foli: 24	Richard Pousgaard, DTU, DK-2800 Kongens Lyngby, Denmark
home:/db/voll8/3a37/3a37	April 5, 2006, 09:18		

Formalisation

- Which of the above alternative ways of designations are we going to formalise?
- Typically we formalise a *designation description* as shown in the next example.
- It allows for the general case of several alternatively expressible designations.

4435 2726, Fol: +45 4435 2874 © Olsen Egeboer, Fødsling 11, DK-2800 Høvd, Denmark E-mail: d@hmm.dtu.dk, l@ipm@ipm.dtu.dk, d@ipm@ipm.dtu.dk, www.ipm.dtu.dk/46

Example 7.82 Units:

- We choose to introduce just one type (i.e., sort), **U**, for rail units,
 - and we choose to let
 - * (recognition rule-oriented) observer functions
 - * and suitable axioms
- help separate rail units into a partition of its more specialised rail units.

Let **W** denote a concept of undirected “ways” through a unit (Fig. 7.13).

type
U, C, W

value
 obs.Cs: $U \rightarrow C\text{-set}$
 obs.Ws: $U \rightarrow W\text{-set}$
 is.linear, is.switch, is.simpl.cross, is.switch.cross: $U \rightarrow \mathbf{Bool}$

axiom

[1] $\forall u:U, \exists c,c',c'',c''':C \cdot$
 $\mathbf{card}\{c,c',c'',c'''\}=4 \wedge \text{obs.Cs}(u) \subseteq \{c,c',c'',c'''\} \Rightarrow$
 $\text{is.linear}(u) \Rightarrow \text{obs.Cs}(u) \subseteq \{c,c'\} \wedge \mathbf{card}\ \text{obs.Ws}(u)=1 \vee$
 $\text{is.switch}(u) \Rightarrow \text{obs.Cs}(u) = \{c,c',c''\} \wedge \mathbf{card}\ \text{obs.Ws}(u)=2 \vee$
 $\text{is.simpl.cross}(u) \Rightarrow \text{obs.Cs}(u) = \{c,c',c'',c'''\} \wedge \mathbf{card}\ \text{obs.Ws}(u)=2 \vee$
 $\text{is.switch.cross}(u) \Rightarrow \text{obs.Cs}(u) = \{c,c',c'',c'''\} \wedge \mathbf{card}\ \text{obs.Ws}(u)=4,$

[2] $\forall c:C \cdot \mathbf{card}\{u \mid u:U \cdot c \in \text{obs.Cs}(u)\} \leq 2,$

[3] **let** $\text{lus} = \{u:U \mid \text{is.linear}(u)\}, \text{sus} = \{u:U \mid \text{is.switch}(u)\},$
 $\text{cus} = \{u:U \mid \text{is.simpl.cross}(u)\}, \text{scus} = \{u:U \mid \text{is.switch.cross}(u)\}$
in $\text{lus} \cap \text{sus} = \text{lus} \cap \text{cus} = \text{lus} \cap \text{scus} = \{\} \wedge$
 $\text{sus} \cap \text{cus} = \text{sus} \cap \text{scus} = \text{cus} \cap \text{scus} = \{\} \mathbf{end}$

- **U** denotes the possibly infinite set of all *designations* satisfying the *rail_unit recognition rule*.
- **C** stands for the possibly infinite set of all *designations* satisfying the *connector recognition rule*.
- **W** denotes the concept of *way* (or *link*).
- The predicates
 - * *is.linear, is.switch, is.simpl.cross* and *is.switch.cross*
 further constrain the *rail_unit recognition rule*,
- as indicated by the **axiom**.

Observe how the formalisation fits with the narration. ■

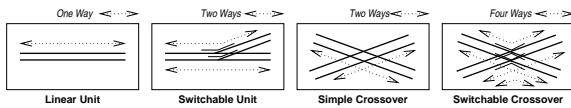


Figure 7.13: Undirected ways through units

Observer Functions and Identification

- The *observer functions*, eg. *obs.Cs* and *obs.Ls*, are the closest counterpart to the undefined terms of *recognition rules* that we have.
- These *observer functions* cannot be defined.
- They are postulated. They “arise” as the result of *identification*.

- Given an actual *universe of discourse* — to which the above *designations* are said to apply —
 - * one can now “define” these *observer functions*
 - * by “walking out and into” the *universe of discourse* and
 - * by providing the *identification*.
- To get a better grasp of possible relationships between recognition rules and observer functions let us give a further example.

Example 7.83 Rail Net: Lines and Stations:

- A *railway net* consists of [one or more] lines and [two or more] stations.
- A *line* is a linear sequence of one or more linear rail units.
- A *station* is any composition (connection) of rail units. [A line connects exactly two distinct stations.]
- The above (excluding the bracketed parts) may be claimed to be a definition, and we shall soon turn to a treatment of definitions.
- But we claim that each of the (unbracketed parts of the) above notions can be designated.
- The *italic text* above may not look like *recognition rules*, but they are!

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.3 Observer Functions and Identification		Page 718, Topic: 26, Foli: 33	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

- One can indeed “walk out, into” the railway system domain and, with a sweeping hand, express:
 - ★ That “thing” there is a linear rail unit. That one, next to it, is likewise.
 - ★ That particular sequence of those two linear units I just pointed to forms (part of) a line.
 - ★ This “thing” here is a rail unit of a station.
 - ★ It is a crossover.
 - ★ Here is the connector that separates a line from a station.
 - ★ No rail unit is both a rail unit of a line and of a station, or of two otherwise distinct lines or stations.

■

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Observer Functions and Identification		Page 719, Topic: 26, Foli: 34	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

type
N, L, S, U, C, L

value
obs_Ls: N → L-set, obs_Ss: N → S-set,
obs_Us: (N|L|S) → U-set

axiom
 $\forall n:N, l,l':L, s,s':S \cdot$
 $\text{card } \text{obs_Ls}(n) \geq 1 \wedge \text{card } \text{obs_Ss}(n) \geq 2 \wedge$
 $\{l,l'\} \subseteq \text{obs_Ls}(n) \wedge \{s,s'\} \subseteq \text{obs_Ss}(n) \Rightarrow$
 $l \neq l' \Rightarrow \text{obs_Us}(l) \cap \text{obs_Us}(l') = \{\} \wedge$
 $s \neq s' \Rightarrow \text{obs_Us}(s) \cap \text{obs_Us}(s') = \{\} \wedge$
 $\text{obs_Us}(l) \cap \text{obs_Us}(s) = \{\} \wedge$
 exactly_two_distinct_stations(l,obs_Ss(n))

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4 Mathematical and Computing Entities		Page 720, Topic: 26, Foli: 35	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

Mathematical and Computing Entities

- From a formal point of view, i.e., from the points of view of mathematics and computer and computing science, which are the kinds of designations?
- Which kinds of mathematical entities are they?
- Or, in the jargon of computing: Which types of computer and computing science entities are they?

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
1 Mathematical Entities		Page 721, Topic: 26, Foli: 36	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

Mathematical Entities

- When viewed mathematically, are the designations scalars, like numbers (integers, reals (or complex), rationals, transcendental), or truth values, etc.?
- Or are they composite, like sets, Cartesians, lists, maps or functions (or algebras, etc.)?
- And if the latter, then which are their component elements?

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.2 Computing Entities		Page 722, Topic: 26, Foli: 37	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

Computing Entities

- When viewed computer and computing science wise, are the designations values (of, for example, the above-mentioned mathematical kind (where models are like algebras))?
- Or are they types of these, i.e., types as we know them in computer science: simple lattices, Scott domains or otherwise?
- Or are they events — whatever they are?
- If lists, then what do these lists model: behaviours (of processes, in terms of traces of actions and/or events, etc.), or other?
- In any case, when we model computing entities (values, types, events, (process) behaviours, semantic algebras), then we model them in terms of the above-mentioned mathematical entities.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Awareness		Page 723, Topic: 26, Foli: 38	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

Awareness

- We do not intend to give a full answer to the question:
- which kind (type) of formal entities can designations be modelled by?
- We only advise that the practicing, professional software engineer be reasonably well-versed in these matters:
 - ★ modelling designations formally in terms of mathematical entities — perhaps couched in the computing jargon of, for example,
 - ◊ types and values,
 - ◊ events and (process) behaviours,
 - ◊ semantic algebras or other.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.4 Some Guidelines		Page 724, Topic: 26, Foli: 39	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

Some Guidelines

- In the following we shall try to stay clear of the deeper problems of conceptual modelling as currently studied in computer and computing science
- But we must — since it cannot be avoided (if we are to cover any ground at all) — postulate some possibilities of concept modelling (since that is what it, in essence, is all about).
- We do so in order to identify some designation (etc.) principles and techniques.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
4 Some Guidelines		Page 725, Topic: 26, Foli: 40	Richard Preece, Peter, Ole-Anders Kjaer Lyngby, Denmark
home:/db/vollf/3sk7/3sk7	April 5, 2006, 09:18		

- Some *designations* are specific, “one of a kind” things (“that rail unit there”), or they are specific events, or specific behaviours (etc.).
- Or *designations* are types or models (i.e., algebras) over these.
- Some designations are those of components of contexts and states.
- Contexts are entities whose properties — whose attributes, whose values — remain static over time.
- States are such whose values change with time — they are dynamic.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.4 Some Guidelines	April 5, 2006, 09:18	Page 726, Topic: 26, Foli: 43
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

- Examples of contexts are: *road* and *rail nets* (when viewed topologically and over time periods that are not too large), similarly for airline *timetables*.
- Examples of states are: *road*, rail and air traffic, hence *trucks*, trains and aircraft, and the *hubs* (where they meet and passengers embark and disembark, or where *freight receipt*, *transfer* and *delivery* can take place).
- A *bill of lading* is somewhere “in between”: The *route* along which a *freight item* should be conveyed is, we assume, static, but the *bill of lading* is (probably) marked (“updated”) to show the (believed) current (or last recorded) position of the *freight item* for which it is the *bill of lading*.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
4 Some Guidelines	April 5, 2006, 09:18	Page 727, Topic: 26, Foli: 42
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

- But, in any case, all these examples can be modelled as typed, usually composite values, fixed or variable.
- These values we consider inert:
- They do not change by themselves.
- Some external action has to change them.
- Designations thus could, alternatively, be the, or a, specific action (active phenomenon, *function*, operation, task, procedure) of *inscribing a freight item for conveyance with a logistics firm, loading it onto the truck, unloading it*, etc.
- In this case we refer to the designations as function values.
- Or designations could be the *events of truck departure* from, resp. *arrival at a hub*.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.4 Some Guidelines	April 5, 2006, 09:18	Page 728, Topic: 26, Foli: 43
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

- Or designations could be part or the entire (*process*) behaviour *from inscription to unloading* described in terms of the specific sequence of inert and dynamic phenomena.
- In this case we refer to the designations as *behaviours* (a certain kind of function values).

• • •

- The above explication presents just one kind of conceptual modelling approach.
- It is “slanted” towards **CSP** and **RAISE**.

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
4 Some Guidelines	April 5, 2006, 09:18	Page 729, Topic: 26, Foli: 44
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

Principles 7.31 Conceptual Frameworks: When setting out on a first description, identify which conceptual modelling framework you intend to work within. ■

Techniques 24 Framework Model: Among the conceptual modelling frameworks that can be believably offered are:

- B, VDM, Z
- RAISE
- CafeOBJ, CASL

We list three groups. The first are solely model-oriented, the last is solely property-oriented (in the algebraic semantics style). **RAISE** basically offers both styles. ■

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.5 A “Large” Example	April 5, 2006, 09:18	Page 730, Topic: 26, Foli: 45
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

A “Large” Example

Example 7.84 Transport:

type
Fre, BoL, Trk, Hub

value
m,n: Nat
obs,BoL: Fre → BoL

axiom
m>0 ∧ n>0

type
HIIdx = { | 1..m | }, TIDx = { | 1..n | }

channel
{ ht[h,t]: (Fre × BoL) · h:HIIdx, t:TIDx }

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
5.A “Large” Example	April 5, 2006, 09:18	Page 731, Topic: 26, Foli: 46
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

value

truck: t:TIDx → Trk → **Unit** → **in,out** {ht[h,t]:HIIdx} **Unit**
truck(t)(tr) ≡ (...; **let** h = ... **in** truck(t)(unload(tr)(h,t)) **end**) [] (...)
hub: h:HIIdx → Hub → **in,out** {ht[h,t]:TIDx} **Unit**
hub(h)(hu) ≡ [] {hub(j)(add(ht[h,t]?)(hu))|t:TIDx} [] (...)
load: Fre × BoL → t:TIDx × h:HIIdx → **in,out** ht[h,t] Trk
unload: Trk → (h:HIIdx × t:TIDx) → **out** {ht[i,t]:HIIdx} Trk
unload(tr)(h,t) ≡
let (f,b):(Fre × BoL)·ii(tr,f,b,h) **in** ht[h,t]!(f,b);rem(f,b)(tr) **end**
ii: tr:Trk × f:Fre × b:BoL × h:HIIdx → **Bool** /* b of f of tr mentions h */
rem: Fre × BoL → Trk → Trk
add: Fre × BoL → Hub → Hub

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
2.4.6 Formal Modelling	April 5, 2006, 09:18	Page 732, Topic: 26, Foli: 47
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

Formal Modelling

- It is not the intention of the present section to suggest neither principles nor techniques for exactly how to formally model designations.
- Such principles and techniques were and are the subject of earlier lectures and later lectures.
- Instead it is the aim of the current section to point out the following principle:

Principle 7.32 Choice of Description Style: In considering which things to designate and to describe, let the informal description style be governed by the chosen, most suitable formal description style. ■

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Discussion: Designations	April 5, 2006, 09:18	Page 733, Topic: 26, Foli: 48
home:/db/vollf/3a7/3a7		Richard Preece, Peter Dink, 00-260 Kgl. Lyngby, Denmark

Discussion: Designations

Principle 7.33 Type Versus Value (Instantiation) Modelling:

- In building up designations —
- within the conceptual paradigm (i.e., framework), say, of **RAISE** —
- one must decide
 - ★ whether one is trying to designate the **type** of all those things that are similar to a few designations, i.e., satisfy their common recognition rule — as for **Fre**, **BoL**, **Trk** and **Hub** above —
 - ★ or one is trying to define only a specific **value** (one particular thing) — as for **hub**, **truck**, **load**, **unload**, **ii**, **rem** or **add** above

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3 Explicit Definitions	April 5, 2006, 09:18	Page 734, Topic: 26, Foli: 49
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Explicit Definitions

- First, designations represent one form of definition.
- In any description all terms that are special to the *universe of discourse* being described must be *defined*
 - ★ either through *designations*
 - ★ or by *explicit definitions*.

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Definitions: "The Narrow Bridge"	April 5, 2006, 09:18	Page 735, Topic: 26, Foli: 50
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Definitions: "The Narrow Bridge"

- The example above contained only designatable quantities, or so we claimed.
- But that may not always be possible, or convenient.
- Sometimes it is easier to *define a term* by giving it a *definition*, but (notably) using already *defined* (including *designated*) **terms**.

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3.1 Definitions: "The Narrow Bridge"	April 5, 2006, 09:18	Page 736, Topic: 26, Foli: 51
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Example 7.85 Rail Lines:

- A *railway line* is a *linear, acyclic, sequence of linear rail units*, that is, a *sequence where adjacent elements of the sequence are rail units that share exactly one connector*.

type

U, C

L1' = U-set

L1 = { | us:L1' · wf.L1(us) | }

L2 = U*

L2 = { | us:L2 · wf.L2(us) | }

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Definitions: "The Narrow Bridge"	April 5, 2006, 09:18	Page 737, Topic: 26, Foli: 52
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

value

obs.Cs: U → C-set

wf.L1: U-set → Bool

wf.L1(us) ≡ ∃ ul:U* · len ul = card us ∧ elems ul = us ∧ wf.L2(ul)

wf.L2: U* → Bool

wf.L2(ul) ≡

∀ u:U · u ∈ elems ul ⇒ is_linear(u) ∧

∀ i:Nat · {i,i+1} ⊆ inds ul ⇒

card(obs.Cs(ul[i]) ∩ obs.Cs(ul[i+1]))=1 ∧

len ul > 1 ⇒ obs.Cs(hd ul) ∩ obs.Cs(ul[len ul]) = { }

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3.1 Definitions: "The Narrow Bridge"	April 5, 2006, 09:18	Page 738, Topic: 26, Foli: 53
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

- Here the terms **linear rail unit** and **connector** are already designated terms;
- and the concept of a **line** is defined in terms of those designations.
- Recall an earlier axiom which stated that for any connector there are at most two units sharing that connector.
- That axiom is assumed above.

■

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Definitions: "The Narrow Bridge"	April 5, 2006, 09:18	Page 739, Topic: 26, Foli: 54
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Techniques 25 Definitions must be expressed in terms of designated and/or defined terms, and ultimately definitions must be based on designations. ■

Principles 7.34 Although it may seem utterly obvious we urge the developer to *develop alternative definitions*. ■

- Here we developed alternatives L1 and L2.
- Although one may claim that lines could be designations, since, in a sense, they are tangible, we have here defined the concept of line.
- We follow the advice of Michael Jackson when we raise it to a principle:

■

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
3.2 Definition of Abstract, Intangible Concepts	April 5, 2006, 09:18	Page 740, Topic: 26, Foli: 55
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Principle 7.35 *The Narrow Bridge*: Seek as few designations as possible: Enough to define all the other possibly designatable as well as all the desirable abstract, nontangible concepts. ■

Definition of Abstract, Intangible Concepts

- We shall next define some abstract, intangible concepts.

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

SOFTWARE ENGINEERING: Domains, Requirements and Software Design	Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark
Definition of Abstract, Intangible Concepts	April 5, 2006, 09:18	Page 741, Topic: 26, Foli: 56
home:/db/voll8/3ak7/3ak7		Richard Preece, Peter, DK-2600 Kong Lyngby, Denmark

Example 7.86 *Path*:

- A *rail unit* defines a concept of *path*.
- A *path* (through a *rail unit*) represents the ability for a *train to move along that path, through the rail unit*.⁹
- (Notice that we neither, at present, designate nor define what we mean by **train** or **move**.)
- We *define* (i.e., we *model*) a *path* as a pair of *connectors*.
- (Based on the definition of **path** we then proceed to define further intangible, i.e., not physically manifest, concepts.)

+45 452 3726 Fax: +45 4450 8874 © 2006 Sigmund Fostang 11, DK-2600 Lyngby, Denmark E-mail: dff@imm.dtu.dk, sp@imm.dtu.dk, dff@imm.dtu.dk, URL: www.imm.dtu.dk/34

⁹When we earlier spoke of a concept of link, that could, in a sense, be seen as an abstract concept, much like a pair of opposite direction paths.

- With any *rail unit* we can associate a *state*: a possibly empty set of *paths*.
- And, finally, we can associate with any *rail unit* its *state space*: the set of all the *states* that a *rail unit* may “be in” over its lifetime as a *rail unit*.

type
 $U, C, P = C \times C$
 $\Sigma = P\text{-set}, \Omega = \Sigma\text{-set}$
axiom
 $\forall (c, c'): P \cdot c \neq c'$
value
 $\text{obs}.\Sigma: U \rightarrow \Sigma, \text{obs}.\Omega: U \rightarrow \Omega$

We shall soon look at possible relations between designations and definitions of the railway system. ■

How Much, How Little to Define?

- The question is now: how much to *define*?
- For the *universes of discourse* of *requirements* and *software design* there is a *utility* concept: We know what we aim at, so we *define* at least that; but why more than that?
- For the *universes of discourse* of *application domains* we do not, as a matter of principle, know what we are aiming at, so here we go for more: as long as some interesting ideas are being *defined*, then why not?

- It is only, we believe, in “roaming around” and experimenting with *definitions*, and, later, with *refutable assertions*, that we may discover the *most interesting, most relevant, most fitting* domain concepts.
- It is in this exploration, and based on *definitions* that we can expect to build theories of specific domains.

Principle 7.36 Exploring Theory Bases: In constructing *domain models*, i.e., *descriptions of domains*, *designate* according to the “*narrow bridge*” principle, and then *define* as many abstract concepts as long as their *definition* (and those of attendant *refutable assertions* help) “reveals” further concepts. ■

A last example may illustrate the previous point.

Example 7.87 Routes:

- (i) A *route* is a sequence of connected *rail units*.
- (ii) A *route* is an *open route* if the *states* of all *rail units* of the *route* are such that there are *paths* in those *states* that also *connect* (in one direction or another).
- (iii) Given a *unit* within a *station* define as *reachable* from the *station unit* all those *lines* that are incident upon and (thus) emanate from that *station*, and for which there is a *route* between the *unit* and the *line*.

- (iv) Given any two *rail units*, an ‘*origin*’ and a possible ‘*destination*’, of a *railway network* define a more general notion of *reachability*, one that is satisfied if there are *open routes* from the ‘*origin*’ to the ‘*destination*’.
- And so forth.
- The present example illustrates the definition of a number of (viz.: four) concepts without a priori making sure that these concepts will serve a useful purpose. ■

Discussion: Definitions

- Choosing an appropriate balance between a reasonably small set of designations and an appropriate (large, “rich”) set of definitions is an art!
- No definite advice can be given, other than perhaps that given above, which, when followed, helps ensure an “appropriate balance”.
- We have now covered two of the informal and formal description principles and their formalisation techniques.
- From the examples so far we can already now conclude that one cannot see from the formulas whether a description formalises a designation or a definition!

- That is, the *pragmatics* of distinguishing between designation and definitions is lost in the formulas.
- Thus it is important with some *syntax* to start the informal and formal texts in order to alert the reader to which is what!

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
4.1 Designation and Definition Assertions		Page 750, Topic: 26, Foli: 65	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Refutable Assertions

Characterisation 7.138 By a *refutable assertion* we mean a claim that may be shown to be wrong.

Designation and Definition Assertions

- In Example 7.83, where we described railway nets of lines and stations, we put in some bracketed sentence parts.
- Those parts did not designate manifest things.
- They expressed suitable, desirable or other relations of designations (and, later, of definitions).
- We shall refer to such constraints as *refutable assertions*.

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
Designation and Definition Assertions		Page 751, Topic: 26, Foli: 66	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Example 7.88 Rail Net Assertion:

- That a *line connects exactly two distinct stations* is one such example of a refutable assertion.
- It implies that there must be *at least two stations in a railway net* — another refutable assertion — and *at least one line* — yet another refutable assertion.
- We asserted elsewhere that *every connector is shared by at most two (otherwise distinct) rail units; that linear rail units have two connectors and define one link; that switchable rail units have three connectors (and define two links); etc.*
- These sentence parts can all be considered refutable assertions.

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
4.1 Designation and Definition Assertions		Page 752, Topic: 26, Foli: 67	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Techniques 26 *Refutable assertions* must be expressed in terms of either designated or defined terms, or both kinds of terms.

- Designations, in a sense, are facts.
- And facts cannot be refuted.
- Definitions are definitions, and as such must be accepted.
- But definitions may seem general and may thus need to be characterised (“tied down”) through some form of constraining assertions.

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
Designation and Definition Assertions		Page 753, Topic: 26, Foli: 68	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Example 7.89 Unit States:

- We continue the rail unit path, state and state space example given earlier.
- *The paths of a rail unit, i.e., the paths of any of its states, must mention only connectors of that unit.*

axiom

$$\forall u:U, c:C, \sigma:\Sigma \cdot$$

$$\text{obs.}\Sigma(u) \in \text{obs.}\Omega(u) \wedge$$

$$\sigma \in \text{obs.}\Omega(u) \Rightarrow \forall (c, c'): C \times C \cdot (c, c') \in \sigma \Rightarrow \{c, c'\} \subseteq \text{obs.}C_s(u)$$

- The above assumes an earlier axiom expressing the fact that the connectors of a path are always distinct.

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
4.1 Designation and Definition Assertions		Page 754, Topic: 26, Foli: 69	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

- This is a refutable assertion:
- One might think of cases where there could be connectors of a path in some unit's state that were not connectors of that unit, or could one?
- We basically will not know till someone one day shows us an understanding of a railway net that “favours” such an interpretation!

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
Analysis		Page 755, Topic: 26, Foli: 70	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Analysis

- What was claimed, above, as examples of refutable assertions, certainly were assertions.
- Whether they will be refuted remains to be seen, but they are potentially refutable.
- Take the example of *any connector* being “constrained” to be *shared by at most two rail units*.
- If one could show, in the future, that there were indeed connectors that were shared by three (or more) rail units, then we have indeed refuted the assertion.
- But is it likely?

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
4.2 Analysis		Page 756, Topic: 26, Foli: 71	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Example 7.90 Stations and Lines: *A station is connected to at least one line.*

- Refuting this assertion would amount to allowing railway nets with completely isolated stations.
- Is that likely?

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

SOFTWARE ENGINEERING: Domains, Requirements and Software Design		Volume 3	Department of Computer Science and Engineering Institute of Informatics and Mathematical Modelling Technical University of Denmark	
“Dangling” Assertions		Page 757, Topic: 26, Foli: 72	Richard Pousset Plank, DK-2800 Kong Lyngby, Denmark	
home:/db/voll8/3a7/3a7	April 5, 2006, 09:18			

Example 7.91 Stations and Open Routes: *For any station, there is at least one potentially open route from, and at least one such route to, some other station of the net.*

- If this assertion is refuted then there will be at least one station from, or into which one cannot “travel”!
- Is that likely?

“Dangling” Assertions

- It is much too easy to write down texts that “weave a web” of pseudodesignations, pseudodefinitions and pseudoassertions.

+45 452 3726, Fax: +45 4493 8874 © Olsen Egeboer, Frederik 11, DK-2800 Høvd, Denmark E-mail: dahl@im.mim.tu.dk, spang@iprnet.com, dave@iprnet.dk, URS, www.imm.dtu.dk

Example 7.92 Freight Transport Bill of Lading:

- Suppose the following text was all we had: *A bill of lading lists the transport hubs where the referenced freight items first loaded onto a conveyor are being transferred from one conveyor to a next, and are finally unloaded.*
- A transport according to a bill of lading does not visit a hub more than at most once.
- Then what were we to mean?

Suppose further that we went straight ahead and formalised the above italic text:

type

Hub, HIdx, Fre, Con
 $BoL' = (F \times (HIdx \times HIdx^* \times HIdx))$
 $BoL = \{ \} b:BoL' \cdot wf_BoL(b) \{ \}$
 $Transport = Fre \times (Con \times Hub)^*$

value

$wf_BoL: BoL' \rightarrow Bool$
 $wf_BoL(f,(o,hl,d)) \equiv card(\{o,d\} \cup elems hl) = 2 + len\ bol$

$obs_Fn: Fre \rightarrow FIdx$
 $load,unload: Fre \times BoL' \rightarrow Con \rightarrow Con$
 $transfer: Fre \times BoL' \rightarrow (Con \times Con) \rightarrow (Con \times Con)$
 $M: BoL' \rightarrow Transport\text{-}in\text{-}fset$

- What is the problem?
- The problem is that we have taken the liberty of mentioning such possibly designatable phenomena as **Hub**, **Hub index**, **Freight**, **Freight index**, **Conveyor**, **load**, **unload** and **transfer**, and given signature to the function values without having attempted the slightest bit of designations!
- We have likewise defined a **Meaning** function which defines the meaning of a bill of lading as a possibly infinite set of defined **Transports**.
- The assertion *does not visit a hub more than at most once* is thus pretty meaningless.

Techniques 27 Dangling Assertions: Make sure all terms of an assertion are either designated, defined or otherwise bound in the assertion.

Example 7.93 Transport Subtypes: With respect to the previous example we should thus also have properly subtyped the *Transport* type, etc.

Discussion: Refutable Assertions

- Assertions are like axioms. They are self-evident truths — until refuted by observations made in the referenced universe of discourse.
- Assertions are thus not facts, nor are they theorems.
- Theorems are predicate expressions that can be proven to follow from designations, definitions and (the axioms of) refutable assertions.
- Theorems (lemmas, propositions) together with designations, definitions, the axioms of assertions and other theorems (lemmas, propositions) form a theory of the described universe of discourse.

Discussion: Description Principles

- We have related the description principle concepts of designations, definitions and refutable assertions to (albeit sketchy examples of) formal specifications.
- We have, most recently above, seen how formal specifications,
 - ★ cf. the latest formalisation example above,
 - ★ without proper, necessarily informal descriptions that adhere to Jackson's principles,
- can too easily “drug” us into believing that the model identification principle can be skipped.
- We ourselves readily admit to having, far too often for anybody's good, fallen into that trap!