

## Jackson's Description Principles

- The **prerequisite** for studying this chapter is that you have read the two previous chapters.
- The **aims** are to introduce Jackson's concepts of designations, definitions and refutable assertions, and to provide formal tools for the expression of manifestations of these concepts.
- The **objective** is to ensure that the developer becomes a professional specifier.
- The **treatment** is systematic to formal.

We build on ideas eloquently expressed in [189] *Software Requirements & Specifications, a lexicon of practice, principles and prejudices*, by Michael Jackson.

Since all we do is construct, analyse and compare descriptions we shall analyse the concept and constituents of descriptions.<sup>1</sup> A description is about manifest individuals, i.e., phenomena and concepts. Some of these represent, or are intended to represent, facts; others represent mental constructions, i.e., concepts. A description includes *designations*, definitions and refutable descriptions. A description can either be formal or informal. A description sets a scope and a span, and a description expresses moods. By an individual we mean a physically manifest phenomenon. In other contexts we may refer to individuals instead by the term things.

### 7.1 Phenomena, Facts and Individuals

Phenomena are what appears to exist. Domain phenomena are built up from facts about individuals. A fact is a simple truth about the world: It is the smallest unit of observation. Large and complex observations and truths can

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<sup>1</sup> We shall here use the term description as also covering the terms prescription and specification.

be broken down into assertions about several facts. A fact involves one or more *individuals*. Anything can be an individual. Each individual is identical to itself but distinct from all other individuals. If  $x$  is identical to  $y$ , then  $x$  and  $y$  are the same individual:  $x = y$ . If we say that  $x$  is similar or equivalent to  $y$  then  $x$  and  $y$  are distinct individuals that share some property, characteristic, attribute or quality. We may use the operator  $\equiv$  or  $\approx$  to denote this relationship. Although it could be a potential misuse of terminology, we may sometimes say that a phenomenon is identified with (i.e., is bound to) some identifier.

## 7.2 Designations

The first subsections of this section, although they may use the term “universe of discourse”, are formulated to apply primarily to domains (to be described). However, they also apply, *inter alia*, with only minor textual adjustments, to such universes of discourse as requirements and software design. In this and the following two sections we shall implicitly be quoting from Michael Jackson [189].

**Characterisation.** By a *designation description* (for short: *designation*) we syntactically mean a textual triple: a *designation term*, a *designation recognition rule*, and a *designation identification*. ■

**Characterisation.** A *designation term* is a simple, i.e., atomic name. ■

**Characterisation.** A *designation recognition rule* is a text which purports to *designate* something. ■

**Characterisation.** The *designation identification* links the *designation recognition rule* to that something (mentioned in the previous characterisation). ■

To create a *designation* we thus write down two items: a *designation description* and a *designation identification*. The *designation description* is usually of the form:

- *Designated term*:  $dt$ .
- *Recognition rule*:  $dt$ 's satisfy the following informally stated properties:  
... — where all other words, i.e., terms, are assumed to be well-known!

That is: *recognition rules* are expressed only in terms that are otherwise well understood, i.e., part of the folklore. If a recognition rule thus contains a term that is elsewhere *defined*, then it is not a recognition rule but becomes a *definition*.

**Example 7.1** *Rail Units, I:* We give the first in a series of examples relating to rail nets.

- *Designated term:* `rail_unit`, or just `U`.
- *Recognition rule:* A `rail_unit` is a composition of an even number of parallel positioned rails (long, narrow, profiled iron bars) separated such that one can always identify pairs of rails of the composition that are at a specified distance (the rail gauge), and otherwise held together by a set of ties.

When we write `rail_unit` we mean the extensional meaning (the **type**) of that term. ■

When there is doubt as to whether a value of a type or a type is being referred to we shall always mean type — except where value of a type is specifically mentioned.

**Example 7.2** *Rail Units, II:* We give the second in a series of examples relating to rail nets.

- *Designated term:* `linear_rail_unit`, or just `U`, for which a predicate, to be assumed, holds: `is_linear_rail_unit(u)` for all `u` in the extension `U`, i.e., `u:U`.
- *Recognition rule:* A `linear_rail_unit` is a pair of parallel positioned rails (long, narrow, profiled iron bars) separated at a specified distance (the rail gauge) and held together by a set of ties. (The predicate `is_linear_rail_unit(u)` “arises” from the recognition rule.)

Observe that the designation term and the recognition rule dealt with a concrete concept for which the immediate concretisation points to a set of a phenomena. What is being designated and to be recognised is any one member of this set. ■

The latter example is a specialisation of the former. The former example expressed: ... *composition of an even number of parallel positioned ‘rails’* ... thus allowing for, say, two pairs as in a switch or in a simple crossover rail unit.

A *designation identification* is usually of the form: “That ‘thing’ ( $u$ ) there is a  $U$ . So is that ‘thing’ ( $u'$ ) over there!” — thus physically pointing out a *designation set*, that is, instances of values that together become part of a type.

### 7.2.1 Some Observations

In general a *designation description* can be formalised; but one cannot formalise the *identification* — as it relates a formal world to an inherently informal world. As illustrated in the two previous examples, the latter of which

was a specialisation of the former, a *designation* may extensionally denote a class (of things) which can be subdivided into subclasses.

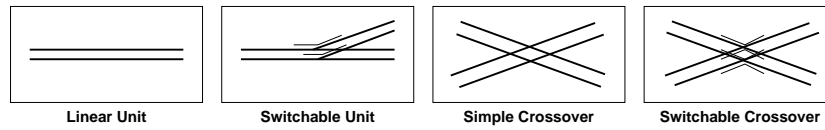


Fig. 7.1. Roughly drawn rail units

**Example 7.3** *Rail Units, III:* We give the third in a series of examples relating to rail nets.

- *Designated term:* `rail_unit`.
- *Recognition rule:* — as for the above, previous `rail_unit` example.
- *Designated term:* `linear_rail_unit`.
- *Recognition rule:* — as for the above, previous `linear_rail_unit` example, and additionally: Thus a `linear_rail_unit` has two “ends” and a single (two-way) link between these ends. Any other `rail_unit` may be *connected* to either of these ends (and if so, then the end is called a *connector*).

To support textually expressed recognition rules it is often useful to deploy graphic means as in Fig. 7.1. Even photographic means could be used. And then one usually could show several photos of variations of the same kind of designated individuals.

- *Designated term:* `switchable_rail_unit`.
- *Recognition rule:* — as for `rail_unit` just above, and additionally: A `switchable_rail_unit` has three *connectors* which “define” (allow, permit) two two-way links through the `switchable_rail_unit`.
- *Designated term:* `simple_crossover_rail_unit`.
- *Recognition rule:* — as for `rail_unit` just above, and additionally: A `simple_crossover_rail_unit` has four *connectors* which “define” (allow, permit) two two-way links through the `simple_crossover_rail_unit`.
- *Designated term:* `switchable_crossover_rail_unit`.
- *Recognition rule:* — as for `rail_unit` just above, and additionally: A `switchable_crossover_rail_unit` has four *connectors* which, together with the switching ability of the unit, “define” (allow, permit) four two-way links through the `switchable_crossover_rail_unit`.
- *Designated term:* `connector`.

- *Recognition rule:* A **connector** is that which allows two **rail\_units** to be connected, “end to end”.

Further recognition rules deal with **connectors** and **rail\_units**. A **connector** is any “end” of a **rail\_unit** to which other **rail\_units** may be ‘connected’. Any one **connector** is shared by at most two **rail\_units**. A **rail\_unit** is mutually exclusive either a **linear\_rail\_unit** or a **switchable\_rail\_unit** or a **simpl\_crossover\_rail\_unit** or a **switchable\_crossover\_rail\_unit**. See axioms [2,3] in a formalisation of the above, in Sect. 7.2.2. ■

As we shall see in Section 7.4, doubt may be raised as to whether some of the text parts of the above recognition rules express assertions. (Designations state facts, not assertions.) For example: “has two ends” (or three, or four). These parts are part of recognition rules, but what about: “any one connector is shared by at most two rail units”? Or the part right after it: “. . . mutually exclusive . . .”? For the last (the mutual exclusion) we can claim it to be a fact, hence it is part of a recognition rule. For the “at most two”, the case is a bit more complicated. And then, when we formalise the whole thing, as we shall see below, and when such a formalisation is compared to that of a refutable assertion, we shall see that, formally speaking, the difference is almost invisible. Thus we must be prepared for the eventuality that the pragmatics of making a distinction between designations, definitions and refutable assertions can not be carried visibly over into formal models of the things being designated or defined, or for which assertions are expressed.

### 7.2.2 Formalisation

Which of the above alternative ways of designations are we going to formalise? Typically we formalise a *designation description* as shown in the next example. It allows for the general case of several alternatively expressible designations.

#### Example 7.4 Units:

##### Formal Presentation: Formalisation of Rail Units, I

We choose to introduce just one type (i.e., sort), **U**, for rail units, and we choose to let (recognition rule-oriented) observer functions and suitable axioms help separate rail units into a partition of its more specialised rail units.

Let **W** denote a concept of undirected “ways” through a unit (Fig. 7.2).

#### type

**U**, **C**, **W**

#### value

**obs\_Cs**: **U** → **C-set**

```

obs_Ws: U→W-set
is_linear,is_switch,is_simpl_cross,is_switch_cross: U→Bool
axiom
[1] ∀ u:U, ∃ c,c',c'',c''':C •
  card{c,c',c'',c'''}=4 ∧ obs_Cs(u)⊆{c,c',c'',c'''} ⇒
  is_linear(u)⇒obs_Cs(u)⊆{c,c'}∧card obs_Ws(u)=1 ∨
  is_switch(u)⇒obs_Cs(u)={c,c',c''}∧card obs_Ws(u)=2 ∨
  is_simpl_cross(u)⇒obs_Cs(u)={c,c',c'',c'''}∧card obs_Ws(u)=2 ∨
  is_switch_cross(u)⇒obs_Cs(u)={c,c',c'',c'''}∧card obs_Ws(u)=4,

[2] ∀ c:C • card{ u | u:U • c ∈ obs_Cs(u) } ≤ 2,

[3] let lus={|u:U•is_linear(u)|}, sus={|u:U•is_switch(u)|},
  cus={|u:U•is_simpl_cross(u)|}, scus={|u:U•is_switch_cross(u)|}
  in lus ∩ sus = lus ∩ scus = lus ∩ cus = {} ∧
  sus ∩ cus = sus ∩ scus = cus ∩ scus = {} end

```

U denotes the possibly infinite set of all *designations* satisfying the *rail\_unit recognition rule*. C stands for the possibly infinite set of all *designations* satisfying the *connector recognition rule*. W denotes the concept of way (or link). The predicates *is\_linear*, *is\_switch*, *is\_simpl\_cross* and *is\_switch\_cross* further constrain the *rail\_unit recognition rule*, as indicated by the **axiom**.

Observe how the formalisation fits with the narration. ■

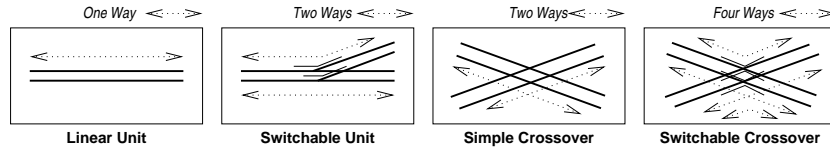


Fig. 7.2. Undirected ways through units

### 7.2.3 Observer Functions and Identification

The *observer functions*, eg. *obs\_Cs* and *obs\_Ls*, are the closest counterpart to the undefined terms of *recognition rules* that we have. These *observer functions* cannot be defined. They are postulated. They “arise” as the result of *identification*. Given an actual *universe of discourse* — to which the above *designations* are said to apply — one can now “define” these *observer functions* by “walking out and into” the *universe of discourse* and by providing the *identification*. To get a better grasp of possible relationships between recognition rules and observer functions let us give a further example.

**Example 7.5** *Rail Net: Lines and Stations:* A railway net consists of [one or more] lines and [two or more] stations. A line is a linear sequence of one or more linear rail units. A station is any composition (connection) of rail units. [A line connects exactly two distinct stations.] The above (excluding the bracketed parts) may be claimed to be a definition, and we shall soon turn to a treatment of definitions. But we claim that each of the (unbracketed parts of the) above notions can be designated. The *italic text* above may not look like *recognition rules*, but they are!

One can indeed “walk out, into” the railway system domain and, with a sweeping hand, express: *That “thing” there is a linear rail unit. That one, next to it, is likewise. That particular sequence of those two linear units I just pointed to forms (part of) a line. This “thing” here is a rail unit of a station. It is a crossover. Here is the connector that separates a line from a station. No rail unit is both a rail unit of a line and of a station, or of two otherwise distinct lines or stations.*

#### Formal Presentation: Formalisation of Rail Units, II

```

type
  N, L, S, U, C, L
value
  obs_Ls: N → L-set, obs_Ss: N → S-set,
  obs_Us: (N|L|S) → U-set
axiom
  ∀ n:N, l,l':L, s,s':S •
    card obs_Ls(n) ≥ 1 ∧ card obs_Ss(n) ≥ 2 ∧
    {l,l'} ⊆ obs_Ls(n) ∧ {s,s'} ⊆ obs_Ss(n) ⇒
    l≠l' ⇒ obs_Us(l) ∩ obs_Us(l') = {} ∧
    s≠s' ⇒ obs_Us(s) ∩ obs_Us(s') = {} ∧
    obs_Us(l) ∩ obs_Us(s) = {} ∧
    exactly_two_distinct_stations(l,obs_Ss(n))

```

We leave out the definition of `exactly_two_distinct_stations`. ■

### 7.2.4 Mathematical and Computing Entities

From a formal point of view, i.e., from the points of view of mathematics and computer and computing science, which are the kinds of designations? Which kinds of mathematical entities are they? Or, in the jargon of computing: Which types of computer and computing science entities are they?

#### Mathematical Entities

When viewed mathematically, are the designations scalars, like numbers (integers, reals (or complex), rationals, transcendental), or truth values, etc.? Or

are they composite, like sets, Cartesians, lists, maps or functions (or algebras, etc.)? And if the latter, then which are their component elements?

### Computing Entities

When viewed computer and computing sciencewise, are the designations values (of, for example, the above-mentioned mathematical kind (where models are like algebras))? Or are they types of these, i.e., types as we know them in computer science: simple lattices, Scott domains [133, 316–324] or otherwise? Or are they events — whatever they are? If lists, then what do these lists model: behaviours (of processes, in terms of traces of actions and/or events, etc.), or other? In any case, when we model computing entities (values, types, events, (process) behaviours, semantic algebras), then we model them in terms of the above-mentioned mathematical entities.

### Awareness

We do not intend to give a full answer to the question: which kind (type) of formal entities can designations be modelled by? We only advise that the practicing, professional software engineer be reasonably well-versed in these matters: modelling designations formally in terms of mathematical entities — perhaps couched in the computing jargon of, for example, types and values, events and (process) behaviours, semantic algebras or other.

### Some Guidelines

In the following we shall try to stay clear, if at all scientifically possible, of the deeper problems of conceptual modelling, viz.: [112, 340], as currently studied in computer and computing science, and in various philosophical discourses, viz.: [237]. But we must — since it cannot be avoided (if we are to cover any ground at all) — postulate some possibilities of concept modelling (since that is what it, in essence, is all about). We do so in order to identify some designation (etc.) principles and techniques.

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Some *designations* are specific, “one of a kind” things (“that rail unit there”), or they are specific events, or specific behaviours (etc.). Or *designations* are types or models (i.e., algebras) over these. Some designations are those of components of contexts and states. Contexts are entities whose properties — whose attributes, whose values — remain static over time. States are such whose values change with time — they are dynamic. Examples of contexts are: *road* and *rail nets* (when viewed topologically and over time periods that are not too large), similarly for airline *timetables*. Examples of states are: *road*, rail and air traffic, hence *trucks*, trains and aircraft, and the *hubs* (where they



meet and passengers embark and disembark, or where *freight receipt*, *transfer* and *delivery* can take place).

A *bill of lading* is somewhere “in between”: The *route* along which a *freight item* should be conveyed is, we assume, static, but the *bill of lading* is (probably) marked (“updated”) to show the (believed) current (or last recorded) position of the *freight item* for which it is the *bill of lading*.

But, in any case, all these examples can be modelled as typed, usually composite values, fixed or variable. These values we consider inert: They do not change by themselves. Some external action has to change them.

Designations thus could, alternatively, be the, or a, specific action (active phenomenon, *function*, operation, task, procedure) of *inscribing a freight item for conveyance with a logistics firm, loading it onto the truck, unloading it*, etc. In this case we refer to the designations as function values. Or designations could be the events of *truck departure* from, resp. *arrival at a hub*. Or designations could be part or the entire (process) behaviour *from inscription to unloading* described in terms of the specific sequence of inert and dynamic phenomena. In this case we refer to the designations as behaviours (a certain kind of function values).

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The above explication presents just one kind of conceptual modelling approach. It is “slanted” towards CSP and RAISE [118, 301].

**Principles.** *Conceptual Frameworks:* When setting out on a first description, identify which conceptual modelling framework you intend to work within. ■

**Techniques.** *Framework Model:* Among the conceptual modelling frameworks that can be believably offered are:

- B [4], VDM [35, 36, 104], Z [162, 341, 343, 377]
- RAISE [117, 118]
- CafeOBJ [85, 110, 111], CASL [28, 62, 251]

We list three groups. The first are solely model-oriented, the last is solely property-oriented (in the algebraic semantics style). RAISE basically offers both styles. If none of these is fully applicable, then be careful, very careful in choosing some “integrated formal specification” approach which takes a “little from this specification language”, a “little from that other language”, and so forth. If the underlying semantic models are compatible, then fine, or else there will be problems in securing an integrated semantics [16]. ■

## A “Large” Example

We illustrate the implications of the above by giving a formal definition in which we then identify various “designatable” quantities.

**Example 7.6** *Transport:*

## Formal Presentation: Formalisation of Transport

The example expands on the *italic text explications* given between the pair of ●●●s above.

**type**

Fre, BoL, Trk, Hub

**value**

m,n: Nat

obs\_BoL: Fre → BoL

**axiom**

m>0 ∧ n>0

**type**

HIdx = { | 1..m | }, TIdx = { | 1..n | }

**channel**

{ ht[h,t]:(Fre×BoL) • h:HIdx, t:TIdx }

**value**

truck: t:TIdx → Trk → **Unit** → **in,out** {ht[h,t]|h:HIdx} **Unit**

truck(t)(tr) ≡ (...; **let** h = ... **in** truck(t)(unload(tr)(h,t)) **end**)∥(...)

hub: h:HIdx → Hub → **in,out** {ht[h,t]|t:TIdx} **Unit**

hub(h)(hu) ≡ ∥ {hub(j)(add(ht[h,t]?)(hu))|t:TIdx} ∥ (...)

load: Fre×BoL → t:TIdx × h:HIdx → **in,out** ht[h,t] Trk

unload: Trk → (h:HIdx × t:TIdx) → **out** {ht[i,t]|i:HIdx} Trk

unload(tr)(h,t) ≡

**let** (f,b):(Fre×BoL)•ii(tr,f,b,h) **in** ht[h,t]!(f,b);rem(f,b)(tr) **end**

ii: tr:Trk×f:Fre×b:BoL×h:HIdx → **Bool** /\* b of f of tr mentions h \*/

rem: Fre×BoL → Trk → Trk

add: Fre×BoL → Hub → Hub

Fre, BoM, Trk and Hub name sets of freight items, bills of material, trucks and hubs. TIdx and HIdx name sets of index sets uniquely identifying trucks and hubs. The logistics system consists of *n* trucks and *m* hubs (here truck yards). The index sets enable us to model that any trucks can dock at any hub. The function names **truck** and **hub** model truck and hub behaviours. Channels **ht[h,t]** model interaction between hubs (**h**) and trucks (**t**). The functions **load** and **unload** effect the loading and unloading actions; **is\_in**, **remove** and **add** name auxiliary functions. The **truck** and **hub** processes are cyclic (never ending). The **ht[h,t]!(f,b)** and **ht[h,t]?** clauses specify synchronisation and communication events (in CSP [168,301,311]). The event concept is thus expressed in terms of channels and synchronised output/input and communication. ■

We remind the reader that the above shows just an example of how designations can be modelled — in the RAISE/CSP paradigm.

## Formal Modelling

It is not the intention of the present section to suggest neither principles nor techniques for exactly how to formally model designations. Such principles and techniques were and are the subject of earlier volumes and later chapters. Instead it is the aim of the current section to point out the following principle:

**Principle.** *Choice of Description Style:* In considering which things to designate and to describe, let the informal description style be governed by the chosen, most suitable formal description style. ■

That is, the describer, i.e., the software engineer, is well served in considering whether that which is to be described informally (rough-sketched, terminologised and narrated) can be given a formal model. If not, then perhaps what one is trying to describe informally is not the right thing to describe! Certainly, if it can be formalised, it can most likely be described well informally. And thus it may be something possibly worth describing.

### 7.2.5 Discussion: Designations

**Principle.** *Type Versus Value (Instantiation) Modelling:* In building up designations — within the conceptual paradigm (i.e., framework), say, of RAISE — one must decide whether one is trying to designate the **type** of all those things that are similar to a few designations, i.e., satisfy their common recognition rule — as for **Fre**, **BoL**, **Trk** and **Hub** above — or one is trying to define only a specific **value** (one particular thing) — as for **hub**, **truck**, **load**, **unload**, **ii**, **rem** or **add** above. ■

This distinction is not, we find, made sufficiently clear in [189]. This is probably because that book, possibly wisely so, and definitely explicitly so, avoids committing itself to any specific formal specification paradigm.

In addition to modelling designations as *types*, *channels* and *values*, one could, in the chosen conceptual framework of RAISE, think of other designation models, for example, *schemes*, *classes*, *objects* and *variables*. Other conceptual frameworks would favour similar or quite dissimilar choices — and the general *designation* principles of [189] apply to any such paradigm.

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We have, so far, only covered one of the three aspects of *descriptions: designations*. Two more remain to be covered: *definitions* and *refutable assertions*. Yet we have been able to demonstrate, in the formal example immediately above, that just by dealing with *designations* one can come a long way.

One may rightfully argue (i) whether all of what that last, albeit only sketched formal example showed was indeed a model only of *designatable* entities, and (ii) whether, as we shall soon see, introducing some *definitions*

might result in a different, perhaps more easily read description. One may also argue (iii) whether some of the axioms are, in fact, *refutable assertions*. We will return to these issues in the closing 'Discussion' sections of this chapter.

### 7.3 Explicit Definitions

First, designations represent one form of definition. In any description all terms that are special to the *universe of discourse* being described must be *defined* either through *designations* or by *explicit definitions*.

#### 7.3.1 Definitions: "The Narrow Bridge"

The example above contained only designatable quantities, or so we claimed. But that may not always be possible, or convenient. Sometimes it is easier to *define* a **term** by giving it a *definition*, but (notably) using already *defined* (including *designated*) **terms**.

**Example 7.7** *Rail Lines: A railway line is a linear, acyclic, sequence of linear rail units, that is, a sequence where adjacent elements of the sequence are rail units that share exactly one connector.*

#### Formal Presentation: Rail Lines

```

type
  U, C
  L1' = U-set
  L1 = { | us:L1' • wf_L1(us) | }
  L2' = U*
  L2 = { | us:L2' • wf_L2(us) | }
value
  obs_Cs: U → C-set
  wf_L1: U-set → Bool
  wf_L1(us) ≡ ∃ ul:U* • len ul = card us ∧ elems ul = us ∧ wf_L2(ul)
  wf_L2: U* → Bool
  wf_L2(ul) ≡
    ∀ u:U • u ∈ elems ul ⇒ is_linear(u) ∧
    ∀ i:Nat • {i,i+1} ⊆ inds ul ⇒
      card(obs_Cs(ul[i]) ∩ obs_Cs(ul[i+1]))=1 ∧
      len ul > 1 ⇒ obs_Cs(hd ul) ∩ obs_Cs(ul[len ul]) = {}

```

Here the terms *linear rail unit* and *connector* are already designated terms; and the concept of a *line* is defined in terms of those designations. Recall an earlier axiom which stated that for any connector there are at most two units sharing that connector. That axiom is assumed above. ■

**Techniques.** *Definitions* must be expressed in terms of designated and/or defined terms, and ultimately definitions must be based on designations. ■

**Principles.** Although it may seem utterly obvious we urge the developer to *develop alternative definitions*. ■

Here we developed alternatives L1 and L2. Although one may claim that lines could be designations, since, in a sense, they are tangible, we have here defined the concept of line. We follow the advice of Michael Jackson [189] when we raise it to a principle:

**Principle.** *The Narrow Bridge:* Seek as few designations as possible: Enough to define all the other possibly designatable as well as all the desirable abstract, nontangible concepts. ■

### 7.3.2 Definition of Abstract, Intangible Concepts

We shall next define some abstract, intangible concepts.

**Example 7.8** *Path:* A rail unit defines a concept of path. A path (through a rail unit) represents the ability for a train to move along that path, through the rail unit.<sup>2</sup>

(Notice that we neither, at present, designate nor define what we mean by train or move.) We define (i.e., we model) a path as a pair of connectors. (Based on the definition of path we then proceed to define further intangible, i.e., not physically manifest, concepts.) With any rail unit we can associate a state: a possibly empty set of paths. And, finally, we can associate with any rail unit its state space: the set of all the states that a rail unit may “be in” over its lifetime as a rail unit.

#### Formal Presentation: Rail Unit States

```

type
  U, C, P = C×C
  Σ = P-set, Ω = Σ-set
axiom
  ∀ (c,c'):P • c≠c'
value
  obs_Σ: U → Σ, obs_Ω: U → Ω

```

In Sect. 7.4 we shall look at possible relations between designations and definitions of the railway system. ■

<sup>2</sup> When we earlier spoke of a concept of link, that could, in a sense, be seen as an abstract concept, much like a pair of opposite direction paths.

### 7.3.3 How Much, How Little to Define?

The question is now: how much to *define*? For the *universes of discourse* of *requirements* and *software design* there is a *utility* concept: We know what we aim at, so we *define* at least that; but why more than that? For the *universes of discourse* of *application domains* we do not, as a matter of principle, know what we are aiming at, so here we go for more: as long as some interesting ideas are being *defined*, then why not? It is only, we believe, in “roaming around” and experimenting with *definitions*, and, later, with *refutable assertions*, that we may discover the *most interesting*, *most relevant*, *most fitting* domain concepts. It is in this exploration, and based on *definitions* that we can expect to build theories of specific domains.

**Principle.** *Exploring Theory Bases:* In constructing *domain models*, i.e., *descriptions* of *domains*, *designate* according to the “*narrow bridge*” principle, and then *define* as many abstract concepts as long as their *definition* (and those of attendant *refutable assertions* help) “reveals” further concepts. ■

A last example may illustrate the previous point.

**Example 7.9** *Routes:* (i) A *route* is a sequence of connected *rail units*. (ii) A *route* is an *open route* if the *states* of all *rail units* of the *route* are such that there are *paths* in those *states* that also *connect* (in one direction or another). (iii) Given a *unit* within a *station* define as *reachable* from the *station unit* all those *lines* that are incident upon and (thus) emanate from that *station*, and for which there is a *route* between the *unit* and the *line*. (iv) Given any two *rail units*, an ‘*origin*’ and a possible ‘*destination*’, of a *railway network* define a more general notion of *reachability*, one that is satisfied if there are *open routes* from the ‘*origin*’ to the ‘*destination*’.

And so forth. The present example illustrates the definition of a number of (viz.: four) concepts without a priori making sure that these concepts will serve a useful purpose. ■

### 7.3.4 Discussion: Definitions

Choosing an appropriate balance between a reasonably small set of designations and an appropriate (large, “rich”) set of definitions is an art!<sup>3</sup> No definite advice can be given, other than perhaps that given above, which, when followed, helps ensure an “appropriate balance”.

We have now covered two of the informal and formal description principles and their formalisation techniques. From the examples so far we can already now conclude that one cannot see from the formulas whether a description formalises a designation or a definition! That is, the *pragmatics* of distinguishing between designation and definitions is lost in the formulas. Thus it

<sup>3</sup> Of course, the hedge “appropriate balance” justifies the property “is an art”.

is important with some *syntax* to start the informal and formal texts in order to alert the reader to which is what!

## 7.4 Refutable Assertions

**Characterisation.** By a *refutable assertion* we mean a claim that may be shown to be wrong. ■

### 7.4.1 Designation and Definition Assertions

In Example 7.5, where we described railway nets of lines and stations, we put in some bracketed sentence parts. Those parts did not designate manifest things. They expressed suitable, desirable or other relations of designations (and, later, of definitions). We shall refer to such constraints as *refutable assertions*.

**Example 7.10** *Rail Net Assertion:* That a line connects exactly two distinct stations is one such example of a refutable assertion. It implies that there must be at least two stations in a railway net — another refutable assertion — and at least one line — yet another refutable assertion. We asserted elsewhere that every connector is shared by at most two (otherwise distinct) rail units; that linear rail units have two connectors and define one link; that switchable rail units have three connectors (and define two links); etc. These sentence parts can all be considered refutable assertions. ■

**Techniques.** *Refutable assertions* must be expressed in terms of either designated or defined terms, or both kinds of terms. ■

Designations, in a sense, are facts. And facts cannot be refuted. Definitions are definitions, and as such must be accepted. But definitions may seem general and may thus need to be characterised (“tied down”) through some form of constraining assertions.

**Example 7.11** *Unit States:* We continue the rail unit path, state and state space example given earlier. *The paths of a rail unit, i.e., the paths of any of its states, must mention only connectors of that unit.*

#### Formal Presentation: Unit States

##### axiom

$$\begin{aligned} &\forall u:U, c:C, \sigma:\Sigma \bullet \\ &\quad \text{obs\_}\Sigma(u) \in \text{obs\_}\Omega(u) \wedge \\ &\quad \sigma \in \text{obs\_}\Omega(u) \Rightarrow \forall (c, c'):C \times C \bullet (c, c') \in \sigma \Rightarrow \{c, c'\} \subseteq \text{obs\_}Cs(u) \end{aligned}$$

The above assumes an earlier axiom expressing the fact that the connectors of a path are always distinct. ■

This is a refutable assertion: One might think of cases where there could be connectors of a path in some unit's state that were not connectors of that unit, or could one? We basically will not know till someone one day shows us an understanding of a railway net that "favours" such an interpretation!

#### 7.4.2 Analysis

What was claimed, above, as examples of refutable assertions, certainly were assertions. Whether they will be refuted remains to be seen, but they are potentially refutable. Take the example of *any connector* being "constrained" *to be shared by at most two rail units*. If one could show, in the future, that there were indeed connectors that were shared by three (or more) rail units, then we have indeed refuted the assertion. But is it likely?

**Example 7.12** *Stations and Lines: A station is connected to at least one line.* ■

Refuting this assertion would amount to allowing railway nets with completely isolated stations. Is that likely?

**Example 7.13** *Stations and Open Routes: For any station, there is at least one potentially open route from, and at least one such route to, some other station of the net.* ■

If this assertion is refuted then there will be at least one station from, or into which one cannot "travel"! Is that likely?

#### 7.4.3 "Dangling" Assertions

It is much too easy to write down texts that "weave a web" of pseudodesignations, pseudodefinitions and pseudoassertions.

**Example 7.14** *Freight Transport Bill of Lading: Suppose the following text was all we had: A bill of lading lists the transport hubs where the referenced freight items first loaded onto a conveyor are being transferred from one conveyor to a next, and are finally unloaded. A transport according to a bill of lading does not visit a hub more than at most once. Then what were we to mean?*



Suppose further that we went straight ahead and formalised the above *italic text*:

Formal Presentation: Freight Transport Bill of Lading

```

type
  Hub, HIdx, Fre, Con
  BoL' = (F × (HIdx × HIdx* × HIdx))
  BoL = { | b:BoL' • wf_BoL(b) | }
  Transport = Fre × (Con × Hub)*
value
  wf_BoL: BoL' → Bool
  wf_BoL(f,(o,hl,d)) ≡ card({o,d} ∪ elems hl)=2+len bol

  obs_Fn: Fre → FIdx
  load,unload: Fre × BoL → Con → Con
  transfer: Fre × BoL → (Con × Con) → (Con × Con)
  M: BoL → Transport-inset

```

What is the problem? The problem is that we have taken the liberty of mentioning such possibly designatable phenomena as **Hub**, **Hub index**, **Freight**, **Freight index**, **Conveyor**, **load**, **unload** and **transfer**, and given signature to the function values without having attempted the slightest bit of designations! We have likewise defined a **Meaning** function which defines the meaning of a bill of lading as a possibly infinite set of defined *Transports*.

The assertion *does not visit a hub more than at most once* is thus pretty meaningless. ■

**Techniques.** *Dangling Assertions:* Make sure all terms of an assertion are either designated, defined or otherwise bound in the assertion. ■

**Example 7.15** *Transport Subtypes:* With respect to the previous example we should thus also have properly subtyped the *Transport* type, etc. ■

#### 7.4.4 Discussion: Refutable Assertions

Assertions are like axioms. They are self-evident truths — until refuted by observations made in the referenced universe of discourse. Assertions are thus not facts, nor are they theorems. Theorems are predicate expressions that can be proven to follow from designations, definitions and (the axioms of) refutable assertions. Theorems (lemmas, propositions) together with designations, definitions, the axioms of assertions and other theorems (lemmas, propositions) form a theory of the described universe of discourse.

## 7.5 Discussion: Description Principles

We have related the description principle concepts of designations, definitions and refutable assertions to (albeit sketchy examples of) formal specifications. We have, most recently above, seen how formal specifications, cf. the latest formalisation example above, without proper, necessarily informal descriptions that adhere to Jackson's principles, can too easily "drug" us into believing that the model identification principle (cf. Sect. 4.4) can be skipped.

We ourselves readily admit to having, far too often for anybody's good, fallen into that trap!

## 7.6 Bibliographical Notes

There is basically just one, major, overriding reference that "over- and fore-shadows" the present chapter, and that is [189].

## 7.7 Exercises

### 7.7.1 A Preamble

We refer to Sect. 1.7.1 for the list of 15 running domain (requirements and software design) examples; and we refer to the introductory remarks of Sect. 1.7.2 concerning the use of the term "selected topic".

### 7.7.2 The Exercises

**Exercise 7.1** *Designations and Recognition Rules.* For the fixed topic, selected by you, list some five phenomena that can be designated. State suitable recognition rules.

**Exercise 7.2** *Explicit Definitions.* For the fixed topic, selected by you, select some two or three designations that you instead decide to define explicitly in terms of other designations. Then provide these explicit definitions.

**Exercise 7.3** *Refutable Assertions.* For the fixed topic, selected by you, try come up with at least one nontrivial possibly refutable assertion. (This may sound difficult, but be open-minded and allow extreme scepticism.)