

A remark on predicate extensions of intuitionistic logic

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Hallmarks of Constructive Reasoning: Disjunction and Existence Properties

Definition

A logic \mathbf{L} is said to have the **disjunction property** (DP), if for every $A \vee B$:
 $\mathbf{L} \vdash A \vee B \Rightarrow \mathbf{L} \vdash A$ or $\mathbf{L} \vdash B$.

A logic \mathbf{L} is said to have the **existence property** (EP), if for every $\exists xA(x)$:
 $\mathbf{L} \vdash \exists xA(x) \Rightarrow$ there exists a v such that $\mathbf{L} \vdash A(v)$.

- DP and EP are regarded as distinguishing features and characteristics of constructivity of **Int**.

$$\text{“} \quad \quad \text{”}$$
$$\exists xA(x) = \bigvee_{u \in U} A(u)$$

Predicate extensions of intuitionistic logic

We consider **P**redicate **E**xtensions of **I**ntuitionistic logic (PEI's).

Int: Intuitionistic predicate logic,

Definition

- (1) The **propositional part** of **L** is the set $\{A : A \text{ is a propositional formula and provable in } \mathbf{L}\}$.
- (2) **L** is said to be a PEI, if its propositional part coincides with that of **Int**.

By the definition, PEI's must resemble **Int** in some logical characters; and PEI's must NOT resemble **Int** in other logical characters.

Int is the logic of **constructivity**.

We deal with two properties related to constructivity of **Int** ;

- Disjunction property (DP),
- Existence property (EP).

Relations in PEI's?

		DP	
		Yes	No
EP	Yes	Int etc.	Int + <i>CD</i> + <i>GJ</i> \vee <i>Z</i> etc. S. 2013(?)
	No	Int + <i>F</i> etc. Nakamura ('83), Minari ('83-'86)	Int + <i>F</i> + <i>Lin</i> [*] etc.

CD : $\forall x(p(x) \vee q) \rightarrow \forall x p(x) \vee q$,

GJ : $\bigwedge_{i=0}^2 ((q_i \rightarrow \bigvee_{j \neq i} q_j) \rightarrow \bigvee_{j \neq i} q_j) \rightarrow \bigvee_{i=0}^2 q_i$,

Z : $\exists x p(x) \rightarrow \forall x p(x)$.

F : $\exists x(p(x) \rightarrow \forall y p(y))$,

Lin^{*} : $\forall x \forall y ((p(x) \rightarrow p(y)) \vee (p(y) \rightarrow p(x)))$

How Many PEI's?

		DP	
		Yes	No
EP	Yes	2^ω	2^ω
	No	2^ω	2^ω

Theorem

- (1) There exist 2^ω PEI's having DP but lacking EP.
- (2) There exist 2^ω PEI's lacking DP and EP.

How?

By constructing a concrete sequence of axiom schemata

$$X_i \quad (i < \omega)$$

such that;

1. for $I \subseteq \omega$, $\mathbf{Int} + \{X_i ; i \in I\}$ is a PEI lacking EP,
2. for each $i < \omega$, $\mathbf{Int} + \{X_j ; j \neq i\} \not\models X_i$,
 - ▶ \implies For $I, J \subseteq \omega$,
 $I = J$ if and only if $\mathbf{Int} + \{X_k ; k \in I\} = \mathbf{Int} + \{X_k ; k \in J\}$.
3. for $I \subseteq \omega$, $\mathbf{Int} + \{X_i ; i \in I\}$ enjoys DP.

No.1.

\mathbf{A} : a finite & strongly compact Heyting algebra; the 2nd greatest = $\star_{\mathbf{A}}$.
Take a finite set $\{p_a ; a \in \mathbf{A}\}$ of propositional variables.

$$\delta(\mathbf{A}) = \{p_{a \cup_{\mathbf{A}} b} \equiv (p_a \vee p_b), p_{a \cap_{\mathbf{A}} b} \equiv (p_a \wedge p_b) ; a, b \in \mathbf{A}\} \\ \cup \{p_{a \rightarrow_{\mathbf{A}} b} \equiv (p_a \rightarrow p_b), p_{a \rightarrow_{\mathbf{A}} 0_{\mathbf{A}}} \equiv (\neg p_a) ; a, b \in \mathbf{A}\},$$

$$J(\mathbf{A}) : \left(\bigwedge \delta(\mathbf{A}) \right) \rightarrow p_{\star_{\mathbf{A}}},$$

$$QJ(\mathbf{A}) : \exists v \{ \left(\bigwedge \delta(\mathbf{A}) \right) [p_{\star_{\mathbf{A}}} := F] \rightarrow (P(v) \rightarrow \forall y P(y)) \}.$$

$$(F : \exists x (P(x) \rightarrow \forall y P(y)).)$$

Lemma

If S is a set of finite and strongly compact Heyting algebras having at least three elements, then $\mathbf{Int} + \{QJ(\mathbf{A}) : \mathbf{A} \in S\}$ is a PEI lacking EP.

No.2.

Fact (cf. Wroński '73)

\mathbf{B} : a Heyting algebra,

$E(\mathbf{B})$: the **propositional** logic determined by \mathbf{B} .

$E(\mathbf{B}) \not\equiv J(\mathbf{A})$ if and only if \mathbf{A} is embeddable into a quotient algebra of \mathbf{B} .

By making use of the **algebraic Kripke sheaf semantics** (S. '99), we have:

Lemma

\mathbf{A}, \mathbf{B} : finite and strongly compact Heyting algebras in which
there exist exactly three elements having no incomparable elements.

$L(\mathcal{K}(\mathbf{B})) \not\equiv QJ(\mathbf{A})$ if and only if \mathbf{A} is embeddable into a quotient alg. of \mathbf{B} .

$\mathcal{K}(\mathbf{B})$: a special algebraic Kripke sheaf determined by \mathbf{B} ,

$L(\mathcal{K}(\mathbf{B}))$: the **super-intuitionistic predicate** logic determined by $\mathcal{K}(\mathbf{B})$.

No.2. (Cont'd)

Suppose we have a sequence $\{\mathbf{A}_i\}_{i < \omega}$ of Heyting algebras such that

- each \mathbf{A}_i is a finite and strongly compact Ha in which there exist exactly three elements having no incomparable elements,
- for $i, j < \omega$ ($i \neq j$), \mathbf{A}_i is not embeddable into any quotient alg. of \mathbf{A}_j .

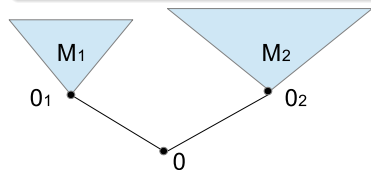
Lemma

For each $i < \omega$, $\mathbf{Int} + \{QJ(\mathbf{A}_j) ; j \neq i\} \not\cong QJ(\mathbf{A}_i)$.

No.3.

Definition

$(\mathbf{M}_i, D_i, \models_i)$ ($i = 1, 2$); Kripke models with the least elements 0_1 and 0_2 .
The **pointed join model** of them:



$$D^\uparrow(a) = \begin{cases} D_1(0_1) \times D_2(0_2) & \text{if } a = 0, \\ D_1(a) \times D_2(0_2) & \text{if } a \in \mathbf{M}_1, \\ D_1(0_1) \times D_2(a) & \text{if } a \in \mathbf{M}_2. \end{cases}$$

Lemma

If A_i is not true in $(\mathbf{M}_i, D_i, \models_i)$ ($i = 1, 2$),
then $A_1 \vee A_2$ is not true in the pointed join model of them.

No.3. (Cont'd)

Definition

A formula A is **axiomatically true** in a Kripke model, if every instance of A is true in it.

Lemma

If the axiomatic truth of A is preserved under pointed-join-model construction, then $\mathbf{Int} + A$ has DP.

Lemma

A: a finite and strongly compact Heyting algebra \mathbf{A} with $|\mathbf{A}| \geq 3$.
The axiomatic truth of $QJ(\mathbf{A})$ is preserved under pointed-join-model construction. Thus, $\mathbf{Int} + QJ(\mathbf{A})$ has DP.

Completing the proof

By modifying Wroński('73)'s technique,
we constructed a concrete sequence $\{\mathbf{A}_i\}_{i < \omega}$ of Heyting algebras such that

- each \mathbf{A}_i is a finite and strongly compact Ha in which there exist exactly three elements having no incomparable elements,
- for $i, j < \omega$ ($i \neq j$), \mathbf{A}_i is not embeddable into any quotient alg. of \mathbf{A}_j ,

Therefore, we have a concrete sequence of axiom schemata

$$QJ(\mathbf{A}_i) \quad (i < \omega)$$

such that;

1. for $I \subseteq \omega$, $\mathbf{Int} + \{QJ(\mathbf{A}_i); i \in I\}$ is a PEI lacking EP,
2. for each $i < \omega$, $\mathbf{Int} + \{QJ(\mathbf{A}_j); j \neq i\} \not\vdash QJ(\mathbf{A}_i)$,
 - ▶ \implies For $I, J \subseteq \omega$,
 $I = J$ if and only if $\mathbf{Int} + \{QJ(\mathbf{A}_k); k \in I\} = \mathbf{Int} + \{QJ(\mathbf{A}_k); k \in J\}$.
3. for $I \subseteq \omega$, $\mathbf{Int} + \{QJ(\mathbf{A}_i); i \in I\}$ enjoys DP.

Concluding Remarks

$$?? \exists x A(x) = \bigvee_{u \in U} A(u) ??$$

- 1 We constructed a **recursively enumerable** set of **concrete** axiom schemata. By adding these schemata to **Int**, we obtained a continuum of PEI's each of which has DP but lacks EP.
- 2 We have four continua of PEI's:
 - ▶ “with EP and DP,”
 - ▶ “without EP and DP,”
 - ▶ “with DP but without EP,” and
 - ▶ “with EP but without DP.”

Other than the last one, Three of them can be obtained by constructing a recursively enumerable set of concrete axioms.

But for the last continuum: “with EP but without DP,” we do not have such a set of axiom schemata, as yet.

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